

## Module II: Methods to Compare Structured Biomedical Data

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## Structured biomedical data

## What makes it "structured"?



Image: MRI
Zhu et al. [2018]


Single-cell time series Liu et al. [2017]


Gene co-expression network Mueller et al. [2017]

## Data transformation

## A review



■ ML toolboxes expect data matrix as samples $\times$ features

- We must transform the (original) structured data into a vectorial representation

Classification
Root of the problem

$$
\begin{array}{llllllllllllllll}
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 & 2 \\
3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 \\
4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 \\
5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 \\
6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 \\
7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 \\
8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 & 8 \\
9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9 & 9
\end{array}
$$

Josef Steppan. https://commons.wikimedia.org/w/index.php?curid=64810040

- Idea: to classify we need a measure of similarity between objects
- How we measure similarity has direct impact on classification

Classification
Root of the problem

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\begin{array}{lllllllllllllllll}
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3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 & 3 \\
4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 & 4 \\
5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 & 5 \\
6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 & 6 \\
7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 & 7 \\
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\end{array}
$$

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■ Idea: to classify we need a measure of similarity between objects

- How we measure similarity has direct impact on classification


# Part I. Kernels 

What is a kernel?<br>Their properties<br>Kernelizing an algorithm

## What is a kernel?

## Intuition



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- Map two objects $\mathbf{x}$ and $\mathbf{x}^{\prime}$ onto Hilbert space $\mathcal{H}$ via mapping $\Phi$
- Compute similarity between $\mathbf{x}$ and $\mathbf{x}^{\prime}$ as inner product $\left\langle\Phi(\mathbf{x}), \Phi\left(\mathbf{x}^{\prime}\right)\right\rangle_{\mathcal{H}}$

■ Kernel trick: Compute inner product in $\mathcal{H}$ as kernel in input space

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\Phi(\mathbf{x}), \Phi\left(\mathbf{x}^{\prime}\right)\right\rangle_{\mathcal{H}}
$$

## What is a kernel?

## Example: XOR problem Gätner [2003]

- Dataset $\mathcal{D}$ with data points in $\mathbb{R}^{2}$ :

$$
(+1,+1),(-1,-1),(+1,-1),(-1,+1)
$$

- Define the mapping
$\phi:\left(x_{1}, x_{2}\right) \mapsto\left(x_{1}^{2}, \sqrt{2} x_{1} x_{2}, x_{2}^{2}\right)$
- Separation is possible with plane orthogonal to $(0,1,0)$
- Explicit mapping
$\Phi(+1,+1)=(+1,+\sqrt{2},+1)$
$\Phi(-1,-1)=(+1,+\sqrt{2},+1)$
$\Phi(+1,-1)=(+1,-\sqrt{2},+1)$
$\Phi(-1,+1)=(+1,-\sqrt{2},+1)$


## What is a kernel?

## Example: XOR problem Gärner [2003]

- Dataset $\mathcal{D}$ with data points in $\mathbb{R}^{2}$ :

$$
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$$

- Separation is possible with plane orthogonal to $(0,1,0)$
- Explicit mapping

$$
\begin{aligned}
& \Phi(+1,+1)=(+1,+\sqrt{2},+1) \\
& \Phi(-1,-1)=(+1,+\sqrt{2},+1) \\
& \Phi(+1,-1)=(+1,-\sqrt{2},+1) \\
& \Phi(-1,+1)=(+1,-\sqrt{2},+1)
\end{aligned}
$$

## EIHzürich

## Inner product

## Definition

■ Let $\mathbf{x}$ and $\mathbf{x}^{\prime} \in \mathbb{R}^{d}$
■ Inner product:

$$
\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle=\mathbf{x}^{\top} \mathbf{x}^{\prime}=\sum_{i=1}^{d} x_{i} x_{i}^{\prime}
$$

- Geometric interpretation:

$$
\begin{aligned}
& k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle \\
& k\left(\mathbf{x}, \mathbf{x}^{\prime \prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime \prime}\right\rangle\left\|\mathbf{x}^{\prime}\right\| \cos \theta_{1} \\
&=\|\mathbf{x}\|\left\|\mathbf{x}^{\prime \prime}\right\| \cos \theta_{2}
\end{aligned}
$$



## Inner product

## Definition

■ Let $\mathbf{x}$ and $\mathbf{x}^{\prime} \in \mathbb{R}^{d}$

- Inner product:

$$
\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle=\mathbf{x}^{\top} \mathbf{x}^{\prime}=\sum_{i=1}^{d} x_{i} x_{i}^{\prime}
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- Geometric interpretation:

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k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle=\|\mathbf{x}\|\left\|\mathbf{x}^{\prime}\right\| \cos \theta_{1}
$$

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime \prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime \prime}\right\rangle=\|\mathbf{x}\|\left\|\mathbf{x}^{\prime \prime}\right\| \cos \theta_{2}
$$

## Inner product

## Definition

■ Let $\mathbf{x}$ and $\mathbf{x}^{\prime} \in \mathbb{R}^{d}$

- Inner product:

$$
\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle=\mathbf{x}^{\top} \mathbf{x}^{\prime}=\sum_{i=1}^{d} x_{i} x_{i}^{\prime}
$$

■ Geometric interpretation:

$$
\begin{aligned}
& k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle \\
& k\left(\mathbf{x}, \mathbf{x}^{\prime \prime}\right)=\left\langle\mathbf{x}, \mathbf{x}^{\prime \prime}\right\rangle=\|\mathbf{x}\| \mathbf{x}^{\prime} \| \cos \theta_{1} \\
& \mathbf{x}^{\prime \prime} \| \cos \theta_{2}
\end{aligned}
$$

## 패zürich

## Kernel

## Revisiting the XOR problem Gärtnere etal. [2003]

- Recall the mapping

$$
\Phi:\left(x_{1}, x_{2}\right) \mapsto\left(x_{1}^{2}, \sqrt{2} x_{1} x_{2}, x_{2}^{2}\right)
$$

- Implicit transformation with kernel:

$$
\begin{aligned}
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right) & =\left\langle\mathbf{x}, \mathbf{x}^{\prime}\right\rangle^{2} \\
& =\left\langle\left(x_{1}, x_{2}\right),\left(x_{1}^{\prime}, x_{2}^{\prime}\right)\right\rangle^{2}=\left(x_{1} x_{1}^{\prime}+x_{2} x_{2}^{\prime}\right)^{2} \\
& =\left(x_{1} x_{1}^{\prime}\right)^{2}+2 x_{1} x_{2} x_{1}^{\prime} x_{2}^{\prime}+\left(x_{2} x_{2}^{\prime}\right)^{2} \\
& =\left\langle\left(x_{1}^{2}, \sqrt{2} x_{1} x_{2}, x_{2}^{2}\right),\left(x_{1}^{\prime 2}, \sqrt{2} x_{1}^{\prime} x_{2}^{\prime}, x_{2}^{\prime 2}\right)\right\rangle \\
& =\left\langle\Phi(\mathbf{x}), \Phi\left(\mathbf{x}^{\prime}\right)\right\rangle
\end{aligned}
$$

## 패zürich

## Kernels

■ Linear kernel

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\sum_{i=1}^{d} x_{i} x_{i}^{\prime}=\mathbf{x}^{\top} \mathbf{x}^{\prime}
$$

■ Polynomial kernel (of degree $p$ )

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left(\mathbf{x}^{\top} \mathbf{x}^{\prime}+c\right)^{p} \quad \text { with } c \in \mathbb{R}, p \in \mathbb{N}^{+}
$$

- Gaussian RBF kernel

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\exp \left(-\frac{1}{2 \sigma^{2}}\left\|\mathbf{x}-\mathbf{x}^{\prime}\right\|^{2}\right) \quad \text { with } \sigma \in \mathbb{R}
$$

## Implicit vector embedding

## Kernel trick

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\Phi(\mathbf{x}), \Phi\left(\mathbf{x}^{\prime}\right)\right\rangle_{\mathcal{H}}
$$

- For all data points $\mathbf{x}$ and $\mathbf{x}^{\prime}$
- Take an algorithm $\mathcal{A}$ defined in terms of inner products between $\mathbf{x}$ and $\mathbf{x}^{\prime}$

■ Substitute the inner products with a kernel $k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)$
■ Obtain kernelized version of $\mathcal{A}$

## $k$-Nearest neighbor

## Algorithm

procedure $k-\mathrm{NN}\left(\mathcal{X}_{\text {train }}, \mathcal{Y}_{\text {train }}, k, \mathbf{x}_{\text {test }}\right)$
$\triangleright \mathcal{X}_{\text {train }}, \mathcal{Y}_{\text {train }}$ : Data points, labels in training set
$\triangleright k$ : Number of neighbors
$\triangleright x_{\text {test }}$ : Data point to predict label
$\triangleright$ Neighborhood

$$
\mathcal{N} \leftarrow \emptyset
$$

for each $x \in \mathcal{X}_{\text {train }}$ do
$d \leftarrow$ get_distance $\left(\mathbf{x}, \mathbf{x}_{\text {test }}\right)$
$\mathcal{N} \leftarrow$ update_neighborhood $(\mathcal{N}, k,(\mathbf{x}, d))$
$\triangleright$ Predict the label based on $\mathcal{N}$

$$
\hat{y} \leftarrow \text { majority_vote }\left(\mathcal{N}, \mathcal{Y}_{\text {train }}\right)
$$

function get_distance ( $\mathbf{x}, \mathbf{z}$ )
$\triangleright$ Compute the Euclidean distance dist $\leftarrow 0$ for $i \leftarrow 1, d$ do dist $\leftarrow \operatorname{dist}+\left(x_{i}-z_{i}\right)^{2}$ return $\sqrt{\text { dist }}$

## Euclidean distance

## As inner product

$$
\begin{aligned}
\|\mathbf{x}-\mathbf{z}\|^{2} & =\sum_{i=1}^{d}\left(x_{i}-z_{i}\right)^{2}=\sum_{i=1}^{d}\left(x_{i}^{2}+z_{i}^{2}-2 x_{i} z_{i}\right) \\
& =\sum_{i=1}^{d} x_{i}^{2}+\sum_{i=1}^{d} z_{i}^{2}-\sum_{i=1}^{d} 2 x_{i} z_{i} \\
& =\langle\mathbf{x}, \mathbf{x}\rangle+\langle\mathbf{z}, \mathbf{z}\rangle-2\langle\mathbf{x}, \mathbf{z}\rangle
\end{aligned}
$$

## Implicit transformation with kernel

$$
\begin{aligned}
\|\mathbf{x}-\mathbf{z}\|^{2}=\|\Phi(\mathbf{x})-\Phi(\mathbf{z})\|^{2} & =\langle\Phi(\mathbf{x}), \Phi(\mathbf{x})\rangle+\langle\Phi(\mathbf{z}), \Phi(\mathbf{z})\rangle-2\langle\Phi(\mathbf{x}), \Phi(\mathbf{z})\rangle \\
& =k(\mathbf{x}, \mathbf{x})+k(\mathbf{z}, \mathbf{z})-2 k(\mathbf{x}, \mathbf{z})
\end{aligned}
$$

## Kernelized $k$-Nearest neighbor

## Possible implementation

procedure $k-\mathrm{NN}\left(\mathcal{X}_{\text {train }}, \mathcal{Y}_{\text {train }}, k, \mathbf{x}_{\text {test }}\right)$
$\triangleright \mathcal{X}_{\text {train }}, \mathcal{Y}_{\text {train }}:$ Data points, labels in training set
$\triangleright k$ : Number of neighbors
$\triangleright x_{\text {test }}$ : Data point to predict label
$\triangleright$ Neighborhood

$$
\mathcal{N} \leftarrow \emptyset
$$

for each $x \in \mathcal{X}_{\text {train }}$ do
$s \leftarrow$ get_similarity $\left(\mathbf{x}, \mathbf{x}_{\text {test }}\right)$
$\mathcal{N} \leftarrow$ update_neighborhood $(\mathcal{N}, k,(\mathbf{x}, s))$
$\triangleright$ Predict the label based on $\mathcal{N}$

$$
\hat{y} \leftarrow \text { majority_vote }\left(\mathcal{N}, \mathcal{Y}_{\text {train }}\right)
$$

```
function get_similarity(x, z)
\ Use a kernel
    return k(x, x) +k(z, z) - 2k(x, z)
```


## Kernelized $k$-Nearest neighbor

## Possible kernels

■ Linear: $k(\mathbf{x}, \mathbf{z})=\langle\mathbf{x}, \mathbf{z}\rangle$
Defaults to standard (squared) Euclidean
function get_similarity $(\mathbf{x}, \mathbf{z})$
$\triangleright$ Use a kernel

$$
\text { return } k(\mathbf{x}, \mathbf{x})+k(\mathbf{z}, \mathbf{z})-2 k(\mathbf{x}, \mathbf{z})
$$ distance

- Polynomial: $k(\mathbf{x}, \mathbf{z})=(\langle\mathbf{x}, \mathbf{z}\rangle+1)^{3}$

Shown to perform better than original $k$-NN in liver disorder database Yuet al. [2002]

## ElHzürich

## Kernel

## Definition Hofmann et al. [2008]

$\square$ For $\mathbf{x}, \mathbf{x}^{\prime} \in \mathcal{X}$ define $k: \mathcal{X} \times \mathcal{X} \mapsto \mathbb{R}$ as

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle\Phi(\mathbf{x}), \Phi\left(\mathbf{x}^{\prime}\right)\right\rangle
$$

where $\Phi$ maps into an inner product space $\mathcal{H}$


■ Given a dataset $\left(\mathbf{x}_{1}, y_{1}\right), \ldots,\left(\mathbf{x}_{n}, y_{n}\right) \in \mathcal{X} \times \mathcal{Y}$ the Gram matrix $\mathbf{K}$ is defined as

$$
\mathbf{K}=\left[\begin{array}{ccc}
k\left(\mathbf{x}_{1}, \mathbf{x}_{1}\right) & \ldots & k\left(\mathbf{x}_{1}, \mathbf{x}_{n}\right) \\
k\left(\mathbf{x}_{2}, \mathbf{x}_{1}\right) & \ldots & k\left(\mathbf{x}_{2}, \mathbf{x}_{n}\right) \\
\vdots & \ddots & \vdots \\
k\left(\mathbf{x}_{n}, \mathbf{x}_{1}\right) & \ldots & k\left(\mathbf{x}_{n}, \mathbf{x}_{n}\right)
\end{array}\right]
$$

## Properties of kernels

## Closure Shawe-Taylor and Cristianini [2004]

Assume $k_{1}$ and $k_{2}$ are kernels over $\mathcal{X} \times \mathcal{X}$
$\mathcal{X} \in \mathbb{R}^{d}$
$\Phi: \mathcal{X} \mapsto \mathbb{R}^{m}$

■ $k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=k_{1}\left(\mathbf{x}, \mathbf{x}^{\prime}\right)+k_{2}\left(\mathbf{x}, \mathbf{x}^{\prime}\right)$
■ $k\left(x, x^{\prime}\right)=a k_{1}\left(x, x^{\prime}\right)$, with $a \in \mathbb{R}^{+}$
$\square k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=k_{1}\left(\mathbf{x}, \mathbf{x}^{\prime}\right) k_{2}\left(\mathbf{x}, \mathbf{x}^{\prime}\right)$

- $k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\mathbf{x}^{\top} \mathbf{B} \mathbf{x}^{\prime}$, with $\mathbf{B} \in \mathbb{R}^{d \times d}$ symmetric and positive semi-definite


# Part II. Kernels on structured data 

String kernels<br>Graph kernels

## Kernels on structured data

## R-convolution kernels Haussler [1999]

■ Intuition: Decompose two objects $X$ and $X^{\prime}$ into sets of substructures $S$ and $S^{\prime}$.

- The idea is to compare all pairs of the substructures of $X$ and $X^{\prime}$ :

$$
k_{R}\left(X, X^{\prime}\right)=\sum_{s \in S, s^{\prime} \in S^{\prime}} k_{\text {base }}\left(s, s^{\prime}\right)
$$

■ For example, a substructure could be the elements of a set, the nodes of a graph or the substrings of a string.

- $k_{\text {base }}$ is an arbitrary vectorial kernel, very often even the Dirac delta kernel.


## Motivation

## Comparison of strings

| Protein binds | No binding |
| :---: | :---: |
| ACTGGCA | GCATTGCTG |
| TTTCGAA | AGTGATC |
| GTAGGAA | CGCATT |
| CCTGGTACA | CCGGTAC |

- Assume we have training data with two sets of labeled DNA short sequences
- Question: To which sequence(s) is a new one most similar?


## Motivation

## Comparison of strings

|  | No binding |
| :---: | :---: |
| ACTGGCA | GCATTGCTG |
| GTTCGAA | AGTGATC |
| CCTGGTACA | CCGGTAC |
|  | TCGGCATT |

■ Assume we have training data with two sets of labeled DNA short sequences

- Question: To which sequence(s) is a new one most similar?


## Spectrum kernel (k-mers)

## Approach Lestie et al. [2002]

- Let $\mathcal{A}$ be the underlying alphabet $\rightarrow \mathcal{A}^{k}$ all strings of length $k$
- Index the feature space by $k$-length strings in $\mathcal{A}^{k}$

■ For each $s \in \mathcal{A}^{k}$, count separately the occurrences in $\mathbf{x}$ and $\mathbf{x}^{\prime}$ in $f_{\mathbf{x}}$ and $f_{\mathbf{x}^{\prime}}$, respectively

- The kernel between two strings $\mathbf{x}$ and $\mathbf{x}^{\prime}$ is defined as:

$$
k\left(\mathbf{x}, \mathbf{x}^{\prime}\right)=\left\langle f_{\mathbf{x}}, f_{\mathbf{x}^{\prime}}\right\rangle
$$

x
$\mathbf{x}^{\prime}$
ACTGGCA GCATTGCTG

$$
\begin{aligned}
& \text { ACT } \\
& \text { CTG } \\
& \text { TGG } \\
& \text { GGC } \\
& \text { GCA }
\end{aligned}
$$

GCA
CAT
ATT
TTG TGC
 GCT

CTG

## Networks

In all walks of life


PPI network
https://commons.wikimedia.org


Chemical compound
https://commons.wikimedia.org


Protein structure
Vishwanathan et al. [2010]

## Graphs

## Some terminology

- A graph $G=(V, E)$ is a set $V$ of vertices and a set $E$ of edges
- Our focus will be on undirected graphs, i.e. edge $e_{i j}=\left(v_{i}, v_{j}\right)$ is unordered
- A graph is labeled if $v_{i} \in V$ has labels in $\mathcal{L}_{V}$ and/or $e_{i j} \in E$ has labels in $\mathcal{L}_{E}$
- Two vertices $v_{i}$ and $v_{j}$ are adjacent if there is an edge ( $v_{i}, v_{j}$ ) connecting them

- A walk $\mathcal{W}$ is a sequence of vertices $\mathcal{W}=v_{1}, v_{2}, \ldots, v_{k}$, where $v_{i}$ is adjacent to $v_{i+1}$

■ A path $\mathcal{P}$ is a walk where $v_{i} \neq v_{j}, \forall i \neq j$

## Graph representation

## Adjacency matrix



|  | a | $b$ | c | $d$ | $e$ |  |  | $g$ | $h$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 0 | 1 | 0 | 0 | 0 |  | 0 | 0 |
| $b$ | 0 | 0 | 1 | 0 | 0 |  |  | 0 | 0 |
|  | 1 | 1 | 0 | 1 | 0 |  |  | 0 | 0 |
|  | 0 | 0 | 1 | 0 | 1 |  |  | 1 | 0 |
|  | 0 | 0 | 0 | 1 | 0 |  |  | 0 | 1 |
|  | 0 | 0 | 1 | 0 | 0 |  |  | 1 | 0 |
|  | 0 |  |  |  |  |  |  |  | 1 |
|  | 0 | 0 | 0 | 0 | 1 |  |  |  | 0 ) |

## Motivation

Comparison of graphs


- Assume we have training data with two sets of labeled molecules
- Key question: How to measure similarity between graphs?
F. Llinares-López \& D. Roqueiro | ML Methods to Compare Structured Biomedical Data


## Motivation

Comparison of graphs


- Assume we have training data with two sets of labeled molecules

■ Key question: How to measure similarity between graphs?

## Why is it difficult to compare graphs?

## Graph isomorphism

- Two graphs $G$ and $H$ are isomorphic,
- if there exists a bijection $f: V(G) \mapsto V(H)$
$■$ such that, for every $(u, v) \in E(G) \Leftrightarrow(f(u), f(v)) \in E(H)$
- Complexity: Claimed to be solvable in quasi-polynomial time Babai [2015]



## Subgraph isomorphism

## Subgraph isomorphism

- Given two graphs $G$ and $H$
$\square$ determine whether a subgraph in $G$ is isomorphic to $H$
- determine the size of the largest common subgraph in $G$ and $H$

■ Shown to be NP-complete Garey and Johnson [1990]

Implications of NP-completeness

- Runtime may grow exponentially with the number of nodes
- For large graphs (many nodes) and for large datasets of graphs this can be a serious problem


## Subgraph isomorphism

## Subgraph isomorphism

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## Graph/subgraph isomorphism

## Special cases

■ Graph and subgraph isomorphism are proved to be solvable in linear time for:
■ Trees

- Planar graphs
- Interval graphs
- and others
- Yet, we will focus on the comparison of graphs in general


## Kernels on graphs

## Steps

- Map objects from $\mathcal{X}$ to Hilbert space $\mathcal{H}$

$$
\Phi: \mathcal{X} \mapsto \mathcal{H}
$$

■ Compute similarity between $G$ and $G^{\prime}$ in $\mathcal{H}$

$$
\begin{aligned}
k\left(G, G^{\prime}\right) & =\left\langle\Phi(G), \Phi\left(G^{\prime}\right)\right\rangle_{\mathcal{H}} \\
& =\Phi(G)^{\top} \Phi\left(G^{\prime}\right)
\end{aligned}
$$



## Kernels on graphs

## Categories of graph kernels we will discuss today

- Random walks
- Paths

■ Limited size sub-graphs

- Sub-tree patterns


## Random walk kernels



## Random walk kernel

## Characteristics

- $p_{s}\left(v_{1}\right)$ to select the starting vertex
- In ith. step
- $p_{t}\left(v_{i} \mid v_{i-1}\right)$ to decide next vertex
- $p_{e}\left(v_{i-1}\right)$ to halt

■ If no prior knowledge $\rightarrow$ uniform probability

- In general, a walk $\mathcal{W}$ rooted at $v_{1}$ is defined as


$$
\mathcal{W}_{v_{1}}=\left(v_{1}, e_{v_{1}, v_{2}}, v_{2}, e_{v_{2}, v_{3}}, \ldots, v_{k}\right)
$$

## Random walk kernel

## Main idea

■ Given a pair of graphs, perform random walks on both
■ Count the number of "matching" walks


## Random walk kernel

## Measuring similarity of walks

■ Define the similarity between the graphs as similarity between walks

$$
k\left(G, G^{\prime}\right)=\sum_{\mathcal{W} \in G} \sum_{\mathcal{W}^{\prime} \in G^{\prime}} k_{\mathcal{W}}\left(\mathcal{W}, \mathcal{W}^{\prime}\right)
$$

- Dirac delta $k_{\mathcal{W}}\left(\mathcal{W}, \mathcal{W}^{\prime}\right)=\left\{\begin{array}{l}1 \text { if } \mathcal{W}=\mathcal{W}^{\prime} \\ 0 \text { otherwise }\end{array}\right.$

> Gärtner et al. [2003]

## Alternative walk kernels Borgwardt et al. [2005]

- Define $k_{\mathcal{W}}\left(\mathcal{W}, \mathcal{W}^{\prime}\right)$ in terms of three kernels
i type kernel
ii length kernel
iii node label kernel



## Random walk kernel

## Challenges

- Computationally demanding $O\left(n^{6}\right)$

■ Tottering: revisiting a vertex immediately after leaving it
■ Halting: downweight of longer walks such that the similarity score is dominated by walks of length 1 Sugivama and Borgwardt [2015]

## Paliative measures

- Fast computation of random walk kernels Vishwanathan et al. [2006]
- Adjustment of transition probabilities to prevent tottering


## Random walk kernel

## Key insigths

- 1. Walk of length $k \rightarrow$ look at $k$ th. power of adjacency matrix

Original adjacency matrix $\mathbf{A}=\mathbf{A}^{0}$


G

## Random walk kernel

## Key insights

- 1. Walk of length $k \rightarrow$ look at $k$ th. power of adjacency matrix

Walks of length $k=2$ in $\mathbf{A}^{2}$


$\mathbf{A}^{2}=$| $a$ |
| :---: |
| $\mathbf{a}$ |
| $b$ |
| $c$ |
| $d$ |
| $\mathbf{e}$ |
| $f$ |
| $g$ |
| $h$ |\(\left(\begin{array}{llllllll}1 \& 1 \& 0 \& 1 \& 0 \& 1 \& 0 \& 0 <br>

1 \& 1 \& 0 \& 1 \& 0 \& 1 \& 0 \& 0 <br>
0 \& 0 \& 4 \& 0 \& 1 \& 0 \& 2 \& 0 <br>
1 \& 1 \& 0 \& 3 \& 0 \& 2 \& 0 \& 2 <br>
0 \& 0 \& 1 \& 0 \& 2 \& 0 \& 2 \& 0 <br>
1 \& 1 \& 0 \& 2 \& 0 \& 2 \& 0 \& 1 <br>
0 \& 0 \& 2 \& 0 \& 2 \& 0 \& 3 \& 0 <br>
0 \& 0 \& 0 \& 2 \& 0 \& 1 \& 0 \& 2\end{array}\right)\)

## Random walk kernel

Key insights Vishwanathan et al. [2010]

- 2. Use direct product graph $G_{\times}$


Given two graphs $G=(V, E)$ and $G^{\prime}=\left(V^{\prime}, E^{\prime}\right)$
$G_{\times}$contains vertex set $V_{\times}$and edge set $E_{\times}$

$$
\begin{aligned}
& V_{\times}=\left\{\left(v_{i}, v_{r}^{\prime}\right): v_{i} \in V, v_{r}^{\prime} \in V^{\prime}\right\} \\
& E_{\times}=\left\{\left(\left(v_{i}, v_{r}^{\prime}\right),\left(v_{j}, v_{s}^{\prime}\right)\right):\left(v_{i}, v_{j}\right) \in E,\left(v_{r}^{\prime}, v_{s}^{\prime}\right) \in E^{\prime}\right\}
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Intuition Random walk on $G_{x}$ is simultaneous random walk on $G$ and $G^{\prime}$


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\end{aligned}
$$

Intuition Random walk on $G_{\times}$is simultaneous random walk on $G$ and $G^{\prime}$

$$
k\left(G, G^{\prime}\right)=\sum_{i=1}^{|V|} \sum_{j=1}^{\left|V^{\prime}\right|} \sum_{k=0}^{\infty} \lambda_{k}\left[\mathbf{A}_{\times}^{k}\right]_{i j}
$$



## Path-based kernels



## EIHzürich

## Shortest path kernel

## Intuition

- Define a graph kernel based on paths
- Avoid drawbacks from random walk kernels: tottering and halting


## Considerations

- Finding all paths in a graph is NP-hard
- Finding the longest paths is NP-hard
- Finding the shortest paths can be done in $O\left(n^{3}\right)$
- Shortest paths may not be unique

- Define kernel based on shortest path distances


## ElHzürich

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## Shortest path kernel

## Steps Borgwardt and Kriegel [2005]

■ Run Floyd-Warshall algorithm to compute all-pairs shortest paths for $G$ and $G^{\prime}$ For directed graphs $\rightarrow$ edge weights $\in \mathbb{R}$ and no negative cycles
For undirected graphs $\rightarrow$ edge weights $\in \mathbb{R}^{+}$

- Define a kernel to compare all pairs of shortest path lengths from $G$ and $G^{\prime}$

$$
k\left(G, G^{\prime}\right)=\sum_{v_{i}, v_{j} \in V} \sum_{v_{r}, v_{s} \in V^{\prime}} k_{\text {length }}\left(d\left(v_{i}, v_{j}\right), d\left(v_{r}, v_{s}\right)\right)
$$

with $d\left(v_{i}, v_{j}\right)$ the length of the shortest path between $v_{i}$ and $v_{j}$

## Shortest path kernel

## Kernel for lengths Borgwardt and Kriegel [2005]

$$
k\left(G, G^{\prime}\right)=\sum_{v_{i}, v_{j} \in V} \sum_{v_{r}, v_{s} \in V^{\prime}} k_{\text {length }}\left(d\left(v_{i}, v_{j}\right), d\left(v_{r}, v_{s}\right)\right)
$$

- $k_{\text {length }}(\cdot, \cdot)$ is kernel to compare the lengths of two shortest paths. Possible implementations:
A linear kernel $k_{\text {length }}\left(d\left(v_{i}, v_{j}\right), d\left(v_{r}, v_{s}\right)\right)=d\left(v_{i}, v_{j}\right) d\left(v_{r}, v_{s}\right)$
A Dirac delta kernel $k_{\text {length }}\left(d\left(v_{i}, v_{j}\right), d\left(v_{r}, v_{s}\right)\right)=\left\{\begin{array}{l}1 \text { if } d\left(v_{i}, v_{j}\right)=d\left(v_{r}, v_{s}\right) \\ 0 \text { otherwise }\end{array}\right.$


## Kernels based on limited-size subgraphs



## Graphlet kernels

## Intuition

■ Graphlet: small connected non-isomorphic induced subgraphs Prǔul [2007]
■ Similarity of graphlet distributions $\Rightarrow$ similarity between corresponding graphs

## Challenge

- Counting graphlets of size $k$ takes $O\left(n^{k}\right)$

■ Very computationally demanding


## Graphlet kernels

## Approach Shervashidze et al. [2009]

- Define $g_{G}^{k}=$ count of graphlets of size $k$ in $G$
- Normalize the counts $n_{G}^{k}=\frac{1}{\# \text { all graphlets in } G} g_{G}^{k}$

■ Define $\mathbf{f}_{G} \in \mathbb{R}^{K}$ containing the normalized frequencies
$\rightarrow$ the ith. entry in $\mathbf{f}_{G}=n_{G}^{i}$

$$
k\left(G, G^{\prime}\right)=\left\langle\mathbf{f}_{G}, \mathbf{f}_{G^{\prime}}\right\rangle
$$

## Considerations

- Perform random sampling

■ Consider graphlets of size $k \in\{3,4,5\}$


## Subtree-based kernels

## Weisfeiler-Lehman kernel

## Algorithm Shervashidze et al. [2011]

function WL_iteration $\left(i, G, G^{\prime}, \mathcal{H}(\cdot)\right)$
$\triangleright$ Step 1. Represent each node as sorted list of neighbors
for each $v \in G$ and $G^{\prime}$ do
$\mathcal{L}_{v} \leftarrow \operatorname{sort}(\mathcal{N}(v))$
$s_{v} \leftarrow$ to_string $\left(\mathcal{L}_{v}\right)$
$\triangleright$ Step 2. Compress list of neighbors into hash value
for each $s_{v}$ do
$h_{v} \leftarrow \mathcal{H}\left(s_{v}\right)$
$\triangleright$ Step 3. Relabel nodes
for each $v \in G$ and $G^{\prime}$ do
label $(v) \leftarrow h(v)$
$\triangleright$ Step 4. Compute the kernel return $k\left(G, G^{\prime}\right)$


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$\triangleright$ Step 3. Relabel nodes
for each $v \in G$ and $G^{\prime}$ do
label $(v) \leftarrow h(v)$
$\triangleright$ Step 4. Compute the kernel
return $k\left(G, G^{\prime}\right)$
$\mathrm{A}, \mathrm{D}$
$\mathrm{B}, \mathrm{C}$
$\mathrm{B}, \mathrm{CE}$
$\mathrm{B}, \mathrm{CE}$
C F

## Weisfeiler-Lehman kernel

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## ElHzürich

## Weisfeiler-Lehman kernel

## Algorithm: Iteration 1, steps 1-3 Shervashidze et al. [2011]



Iteration 0


Step 1 (neighbors)


Step 2 (compression)


Step 3 (relabeling)

## Weisfeiler-Lehman kernel

## Computing the kernel for $m$ iterations

- For any basic kernel, we have

$$
k_{W L}^{m}\left(G, G^{\prime}\right)=k\left(G_{0}, G_{0}^{\prime}\right)+k\left(G_{1}, G_{1}^{\prime}\right)+\ldots+k\left(G_{m}, G_{m}^{\prime}\right)
$$

- Assume $k(\cdot, \cdot)$ counts pairs of nodes with matching labels

$$
k\left(G, G^{\prime}\right)=\left\langle\Phi(G), \Phi\left(G^{\prime}\right)\right\rangle=\sum_{v \in V} \sum_{v^{\prime} \in V^{\prime}} \delta\left(\text { label }(v), \text { label }\left(v^{\prime}\right)\right)
$$

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$$



## Conclusions contd.

## Kernels

- Kernel methods are an efficient tool to compare structured data

■ Gram matrices are easily plugged in into ML toolboxes, e.g. SciKit learn
■ Graph kernels focus on how to compute and compare graph features efficiently

## Acknowledgements

Machine Learning and Computational Biology Lab


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