

# Concepts of Object-Oriented Programming

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Chair of Programming Methodology

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# 7. Ownership Types

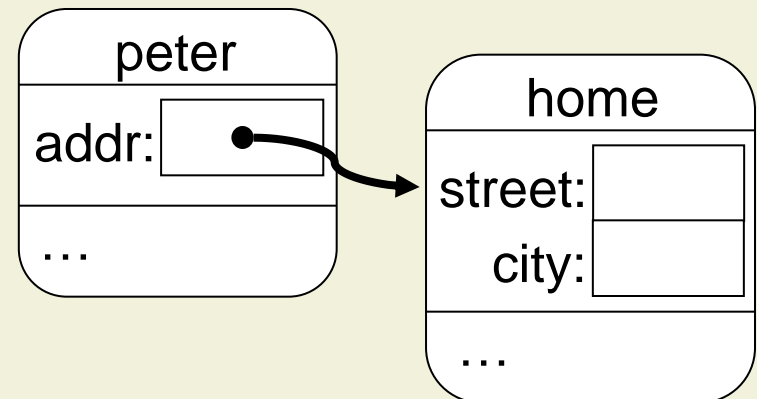
## 7.1 Readonly Types

## 7.2 Topological Types

# Object Structures Revisited

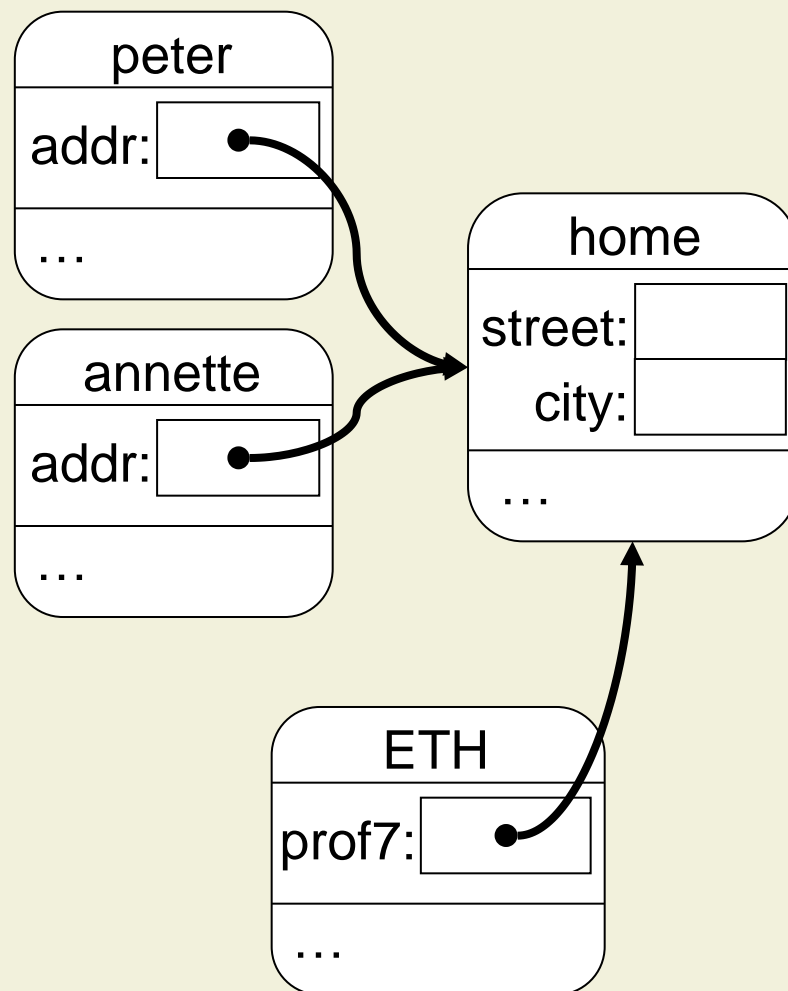
```
class Address ... {  
  private String street;  
  private String city;  
  
  public String getStreet( ) { ... }  
  public void setStreet( String s )  
    { ... }  
  
  public String getCity( ){ ... }  
  public void setCity( String s )  
    { ... }  
  ...  
}
```

```
class Person {  
  private Address addr;  
  public Address getAddr( )  
    { return addr.clone( ); }  
  public void setAddr( Address a )  
    { addr = a.clone( ); }  
  ...  
}
```



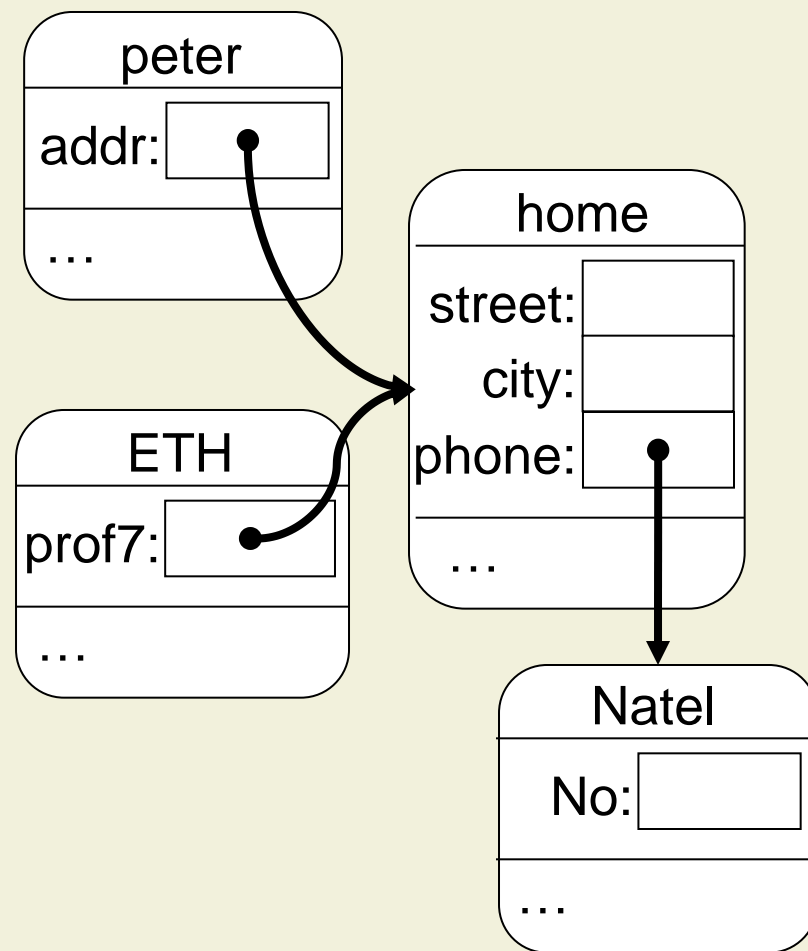
# Drawbacks of Alias Prevention

- Aliases are helpful to **share side-effects**
- Cloning objects is not efficient
- In many cases, it suffices to **restrict access** to shared objects
- Common situation: grant **read access** only



# Requirements for Readonly Access

- Mutable objects
  - Some clients can mutate the object, but others cannot
  - Access restrictions apply to references, not whole objects
- Prevent field updates
- Prevent calls of mutating methods
- Transitivity
  - Access restrictions extend to references to sub-objects



# Readonly Access via Supertypes

```
interface ReadonlyAddress {  
    public String getStreet( );  
    public String getCity( );  
}
```

```
class Address  
    implements ReadonlyAddress ... {  
    ... /* as before */ }
```

```
class Person {  
    private Address addr;  
    public ReadonlyAddress  
        getAddr( )  
        { return addr; }  
    public void setAddr( Address a )  
        { addr = a.clone( ); }  
    ... }
```

- Clients use only the methods in the interface
  - Object remains mutable
  - No field updates
  - No mutating method in the interface

# Limitations of Supertype Solution

- Reused classes might not implement a readonly interface
  - See discussion of structural subtyping
- Interfaces do not support arrays, fields, and non-public methods
- Transitivity has to be encoded explicitly
  - Requires sub-objects to implement readonly interface

```
class Address
    implements ReadonlyAddress ... {
    ...
    private PhoneNo phone;
    public PhoneNo getPhone( )
    { return phone; } }
```

```
interface ReadonlyAddress {
    ...
    public ReadonlyPhoneNo getPhone( );
}
```

# Supertype Solution is not Safe

- No checks that methods in readonly interface are **actually side-effect free**
- **Readwrite aliases** can occur, e.g., through capturing
- Clients can use **casts** to get full access

```
class Person {  
    private Address addr;  
    public ReadonlyAddress getAddr( )  
        { return addr; }  
    public void setAddr( Address a )  
        { addr = a.clone( ); }  
    ...  
}
```

```
void m( Person p ) {  
    ReadonlyAddress ra = p.getAddr( );  
    Address a = (Address) ra;  
    a.setCity( "Hagen" );  
}
```



# Readonly Access in Eiffel

- Better support for fields
  - Readonly supertype can contain getters
  - Field updates only on “this” object
- Command-query separation
  - Distinction between mutating and inspector methods
  - But **queries** are **not checked to be side-effect free**
- Other problems as before
  - Reused classes, transitivity, arrays, aliasing, downcasts

# Readonly Access in C++: const Pointers

```
class Address {  
    string city;  
public:  
    string getCity( void )  
        { return city; }  
    void setCity( string s )  
        { city = s; }  
};
```

C++

```
class Person {  
    Address* addr;  
public:  
    const Address* getAddr( )  
        { return addr; }  
    void setAddr( Address a )  
        { /* clone */ }  
};
```

C++

- C++ supports readonly pointers
  - No field updates
  - No mutator calls

```
void m( Person* p ) {  
    const Address* a = p->getAddr( );  
    a->setCity( "Hagen" );  
    cout << a->getCity( );  
}
```

Compile-time  
errors

# Readonly Access in C++: const Functions

```
class Address {  
    string city;  
public:  
    string getCity( void ) const  
        { return city; }  
    void setCity( string s )  
        { city = s; }  
};
```

C++

```
class Person {  
    Address* addr;  
public:  
    const Address* getAddr( )  
        { return addr; }  
    void setAddr( Address a )  
        { /* clone */ }  
};
```

C++

- const functions must not modify their receiver object

```
void m( Person* p ) {  
    const Address* a = p->getAddr( );  
    a->setCity( "Hagen" );  
    cout << a->getCity( );  
}
```

Call of const  
function allowed

Compile-time  
error

# It wouldn't be C++ ...

```
class Address {  
    string city;  
public:  
    string getCity( void ) const  
        { return city; }  
    void setCity( string s ) const {  
        Address* me = ( Address* ) this;  
        me->city = s;  
    } };
```

C++

```
class Person {  
    Address* addr;  
public:  
    const Address* getAddr( )  
        { return addr; }  
    void setAddr( Address a )  
        { /* clone */ }  
};
```

C++

- const-ness can be cast away
  - No run-time check

```
void m( Person* p ) {  
    const Address* a = p->getAddr( );  
    a->setCity( "Hagen" );  
}
```

Call of const  
function allowed

# It wouldn't be C++ ... (cont'd)

```
class Address {  
    string city;  
public:  
    string getCity( void ) const  
        { return city; }  
    void setCity( string s )  
        { city = s; }  
};
```

C++

```
class Person {  
    Address* addr;  
public:  
    const Address* getAddr( )  
        { return addr; }  
    void setAddr( Address a )  
        { /* clone */ }  
};
```

C++

- const-ness can be cast away
  - No run-time check

```
void m( Person* p ) {  
    const Address* a = p->getAddr( );  
    Address* ma = ( Address* ) a;  
    ma->setCity( "Hagen" );  
}
```

C++

# Readonly Access in C++: Transitivity

```
class Phone {  
public:  
    int number;  
};
```

C++

```
class Address {  
    string city;  
    Phone* phone;  
public:  
    Phone* getPhone( void ) const  
        { return phone; }  
    ...  
};
```

C++

```
void m( Person* p ) {  
    const Address* a = p->getAddr( );  
    Phone* p = a->getPhone( );  
    p->number = 2331...;  
}
```

C++

- **const** pointers are not transitive
- **const**-ness of sub-objects has to be indicated explicitly

# Transitivity (cont'd)

```
class Address {  
    string city;  
    Phone* phone;  
public:  
    const Phone* getPhone( void ) const {  
        phone->number = 2331 ...;  
        return phone;  
    }  
    ...  
};
```

const functions may  
modify objects other  
than the receiver

C++

# Readonly Access in C++: Discussion

## Pros

- const pointers provide readonly pointers to **mutable objects**
  - Prevent field updates
  - Prevent calls of non-const functions
- Work for **library classes**
- Support for arrays, fields, and non-public methods

## Cons

- const-ness is **not transitive**
- const pointers are **unsafe**
  - Explicit casts
- **Readwrite aliases** can occur



# Pure Methods

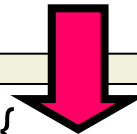
- Tag side-effect free methods as **pure**
- Pure methods
  - Must not contain field update
  - Must not invoke non-pure methods
  - Must not create objects
  - Can be overridden only by pure methods

```
class Address {  
    private String street;  
    private String city;  
    public pure String getStreet( )  
        { ... }  
    public void setStreet( String s )  
        { ... }  
    public pure String getCity( )  
        { ... }  
    public void setCity( String s )  
        { ... }  
    ...  
}
```

# Types

- Each class or interface  $T$  introduces two types
  - Denoted by  $T$  in programs
- **Readwrite type**  $rw\ T$ 
  - Denoted by  **$T$**  in programs
- **Readonly type**  $ro\ T$ 
  - Denoted by **readonly  $T$**  in programs

```
class Person {  
    private Address addr;  
    public ReadonlyAddress  
        getAddr( ) { return addr; }  
    public void setAddr( Address a )  
        { addr = a.clone( ); }  
    ... }  
}
```



```
class Person {  
    private Address addr;  
    public readonly Address  
        getAddr( ) { ... }  
    ...  
}
```

# Subtype Relation

- **Subtyping** among readwrite and readonly types is defined as in Java
  - S extends or implements T  $\Rightarrow$   $rw\ S <: rw\ T$
  - S extends or implements T  $\Rightarrow$   $ro\ S <: ro\ T$
- **Readwrite types** are subtypes of corresponding readonly types
  - $rw\ T <: ro\ T$

```
class T { ... }
```

```
class S extends T { ... }
```

```
S rwS = ...
```

```
T rwT = ...
```

```
readonly S roS = ...
```

```
readonly T roT = ...
```

```
rwT = rwS;
```

```
roT = roS;
```

```
roT = rwT;
```

```
rwT = roT;
```

# Type Rules: Transitive Readonly

```
class Address {  
    ...  
    private int[ ] phone;  
    public int[ ] getPhone( ) { ... }  
}
```

```
class Person {  
    private Address addr;  
    public readonly Address  
        getAddr( ) { return addr; }  
    ...  
}
```

- Accessing a value **of a readonly type** or **through a readonly type** should yield a readonly value

```
Person p = ...  
readonly Address a;  
a = p.getAddr( );  
int[ ] ph = a.getPhone( );
```

# Type Rules: Transitive Readonly (cont'd)

- The type of
  - A field access
  - An array access
  - A method invocation
 is determined by the type combinator ►

►	<i>rw T</i>	<i>ro T</i>
<i>rw S</i>	<i>rw T</i>	<i>ro T</i>
<i>ro S</i>	<i>ro T</i>	<i>ro T</i>

```

Person p = ...
readonly Address a;
a = p.getAddr( );

int[ ] ph = a.getPhone( );
  
```

*ro Address*

*rw int[ ]*

*ro int[ ]*

# Type Rules: Transitive Readonly (cont'd)

- The type of
  - A field access
  - An array access
  - A method invocation
 is determined by the type combinator ►

►	<i>rw T</i>	<i>ro T</i>
<i>rw S</i>	<i>rw T</i>	<i>ro T</i>
<i>ro S</i>	<i>ro T</i>	<i>ro T</i>

Person p = ...

**readonly** Address a;

a = p.getAddr( );

**readonly int[ ]** ph = a.getPhone( );

*ro Address*

*rw int[ ]*

*ro int[ ]*

# Type Rules: Readonly Access

- Expressions of readonly types must not occur as receiver of
  - a **field update**
  - an **array update**
  - an **invocation** of a **non-pure method**
- Readonly types must not be **cast to readwrite types**

```
readonly Address roa;  
roa.street = "Rämistrasse";  
roa.phone[ 0 ] = 41;  
roa.setCity( "Hagen" );
```

```
readonly Address roa;  
Address a = ( Address ) roa;
```

# Discussion

- Readonly types enable **safe sharing of objects**
- Very similar to const pointers in C++, but:
  - Transitive
  - No casts to readwrite types
- All rules for pure methods and readonly types can be **checked statically by a compiler**
- Readwrite aliases can still occur, e.g., by capturing



# 7. Ownership Types

## 7.1 Readonly Types

## 7.2 Topological Types

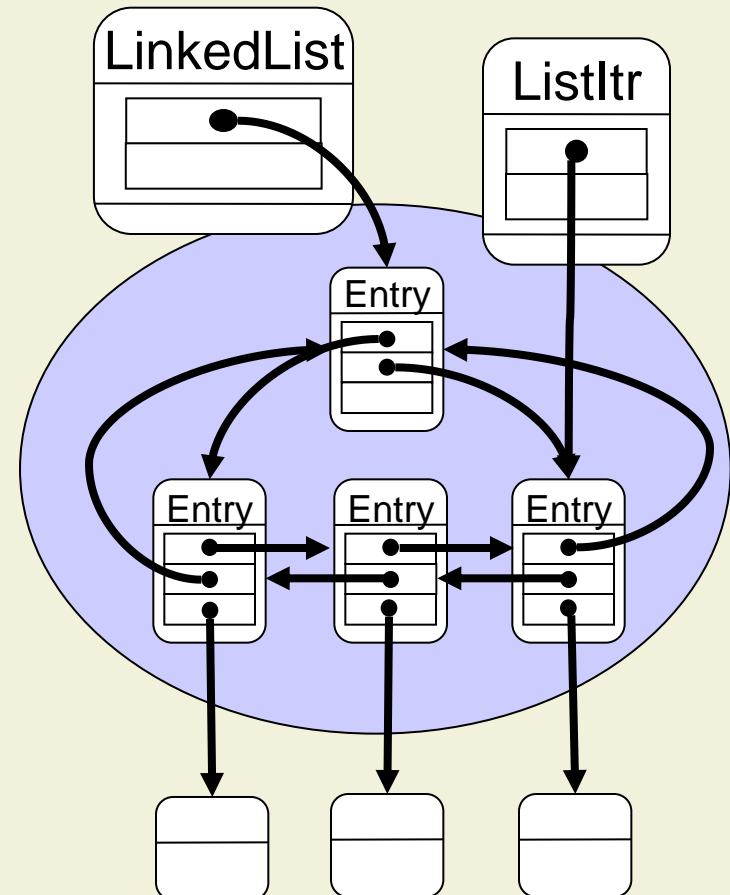
# Object Topologies

- Read-write aliases can still occur, e.g., by capturing or leaking
- We need to distinguish “internal” references from other references

```
class Person {  
    private Address addr;  
    private Company employer;  
    public readonly Address getAddr( )  
        { return addr; }  
    public void setAddr( Address a )  
        { addr = a.clone( ); }  
  
    public Company getEmployer( )  
        { return employer; }  
    public void setEmployer( Company c )  
        { employer = c; }  
  
    ...  
}
```

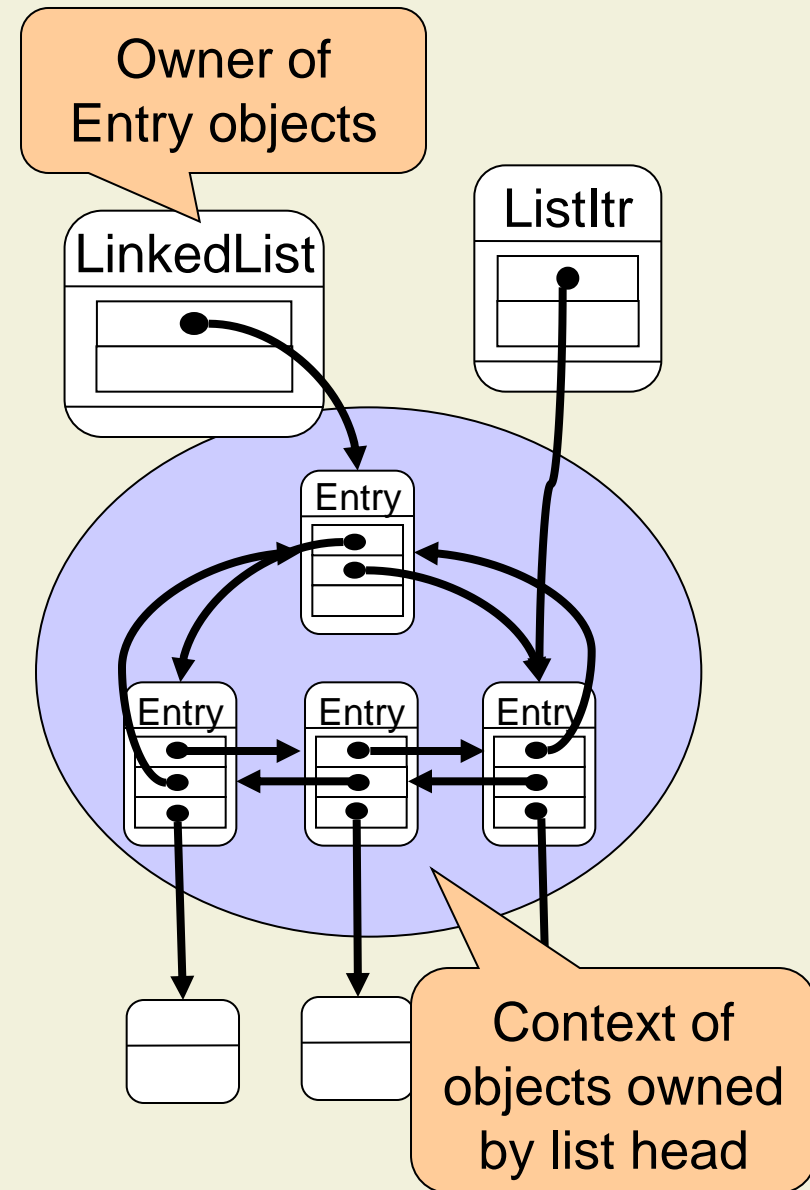
# Roles in Object Structures

- **Interface objects** that are used to access the structure
- **Internal representation** of the object structure
- **Arguments** of the object structure



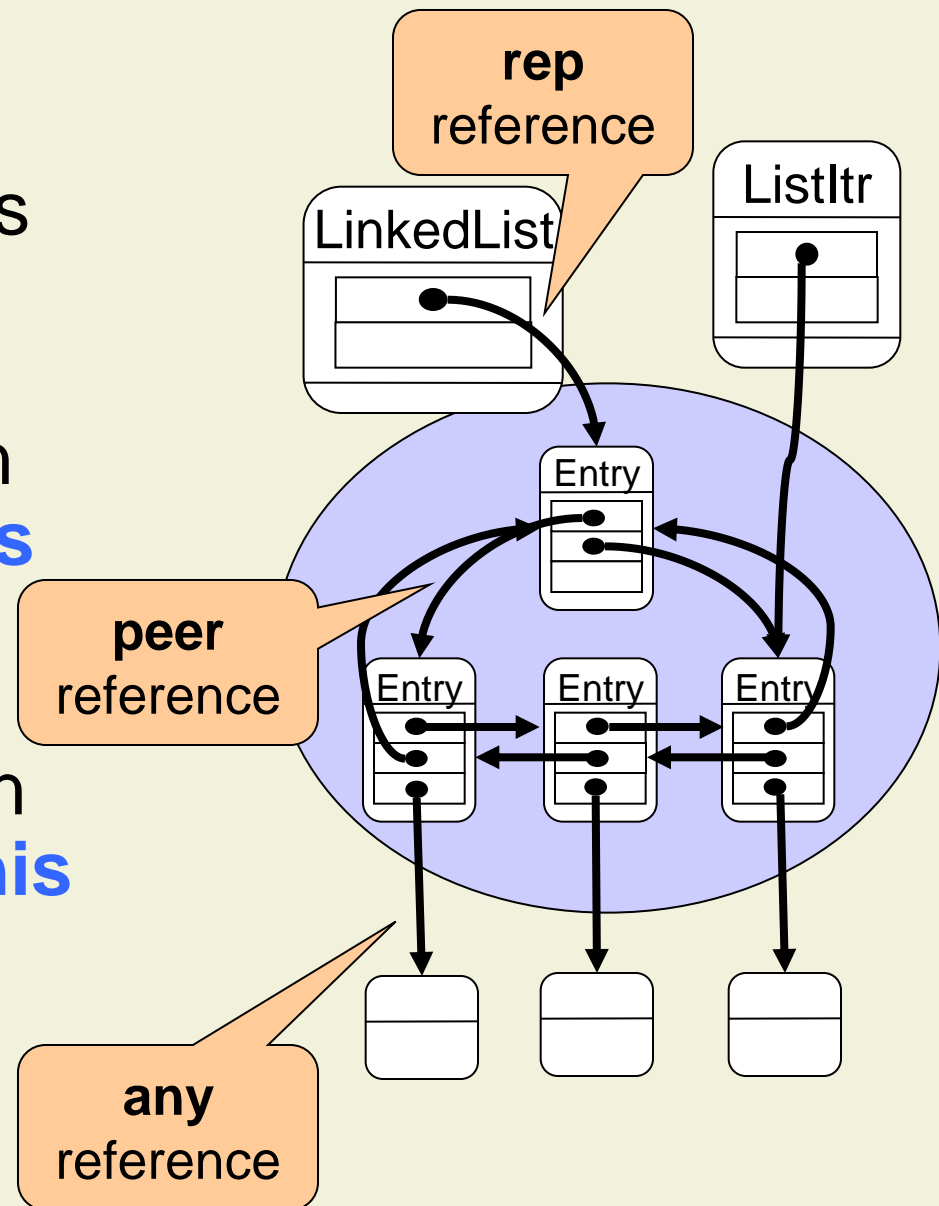
# Ownership Model

- Each object has **zero or one owner objects**
- The set of objects with the same owner is called a **context**
- The ownership relation is **acyclic**
- The heap is structured into a forest of **ownership trees**



# OwnershipTypes

- We use types to express ownership information
- **peer** types for objects in the **same context as this**
- **rep** types for representation objects in the **context owned by this**
- **any** types for argument objects **in any context**



# Example

```
class LinkedList {  
  private rep Entry header;  
  ...  
}
```

A list owns  
its nodes

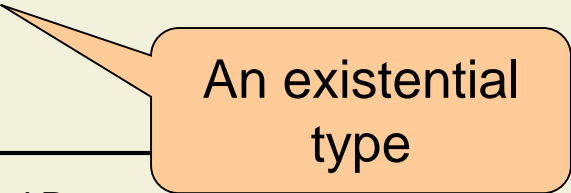
Lists store  
elements with  
arbitrary owners

```
class Entry {  
  private any Object element;  
  private peer Entry previous, next;  
  ...  
}
```

All nodes have  
the same owner

# Type Safety

- Run-time type information consists of
  - The class of each object
  - The **owner** of each object
- Type invariant: the **static ownership information** of an expression  $e$  **reflects the run-time owner** of the object  $o$  referenced by  $e$ 's value
  - If  $e$  has type **rep**  $T$  then  $o$ 's owner is **this**
  - If  $e$  has type **peer**  $T$  then  $o$ 's owner is the **owner of this**
  - If  $e$  has type **any**  $T$  then  $o$ 's owner is **arbitrary**



An existential  
type

# Subtyping and Casts

- For types with identical ownership modifier, subtyping is defined as in Java
  - $\text{rep } S <: \text{rep } T$
  - $\text{peer } S <: \text{peer } T$
  - $\text{any } S <: \text{any } T$
- rep types** and **peer types** are subtypes of corresponding **any types**
  - $\text{rep } T <: \text{any } T$
  - $\text{peer } T <: \text{any } T$

```
class T { ... }
```

```
class S extends T { ... }
```

```
peer T peerT = ...
```

```
any T anyT = ...
```

```
rep S repS = ...
```

```
rep T repT = ...
```

```
repT      = repS;
```

```
anyT      = repT;
```

```
peerT     = ( peer T ) anyT;
```

```
repT      = ( rep T ) anyT;
```

Run-time checks

```
repT      = peerT;
```

```
peerT     = repT;
```

```
repT      = anyT;
```



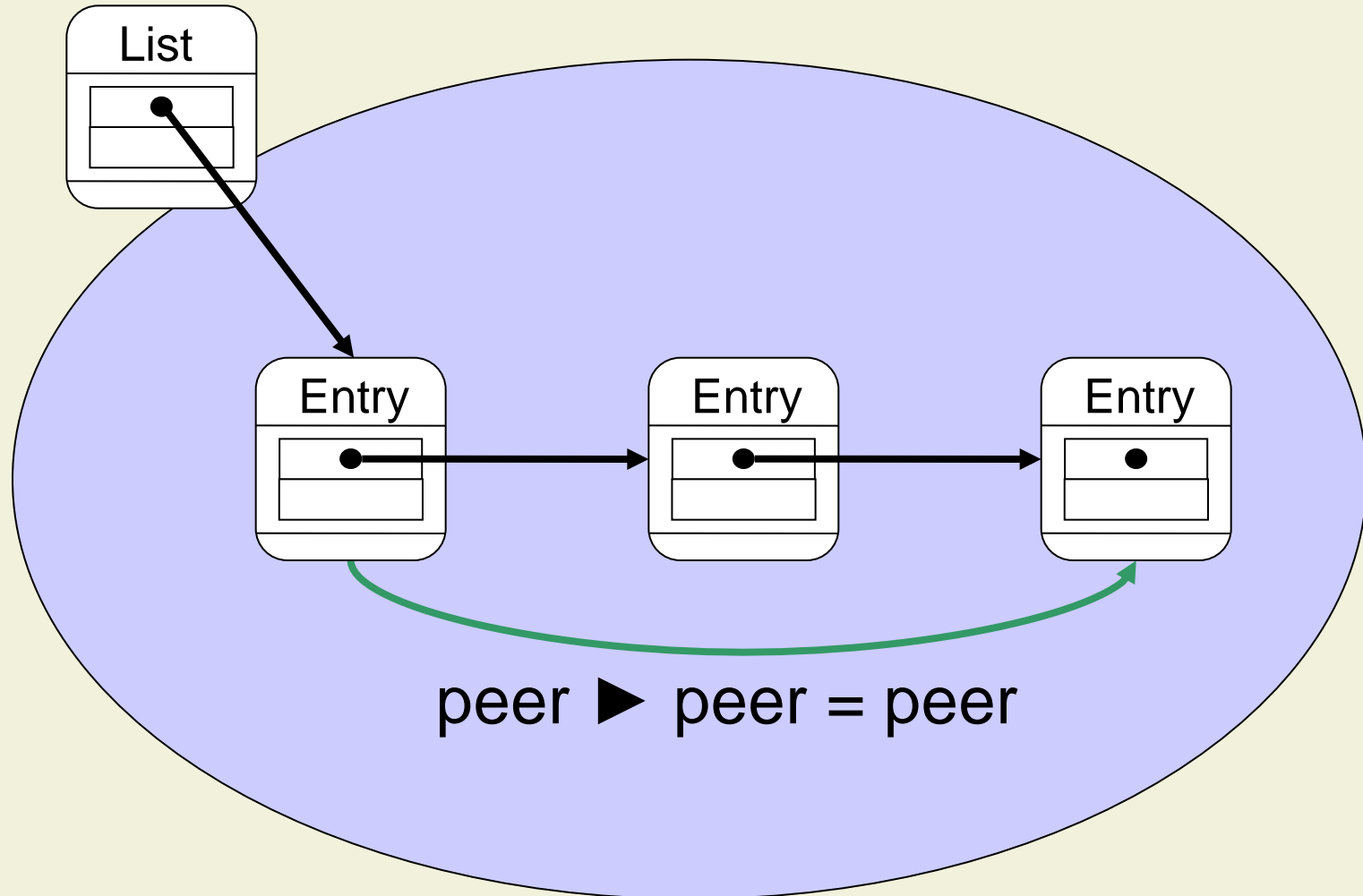
# Example (cont'd)

```
class LinkedList {  
  private rep Entry header;  
  public void add( any Object o ) {  
    rep Entry newE = new rep Entry( o, header, header.previous );  
    ...  
  }  
}
```

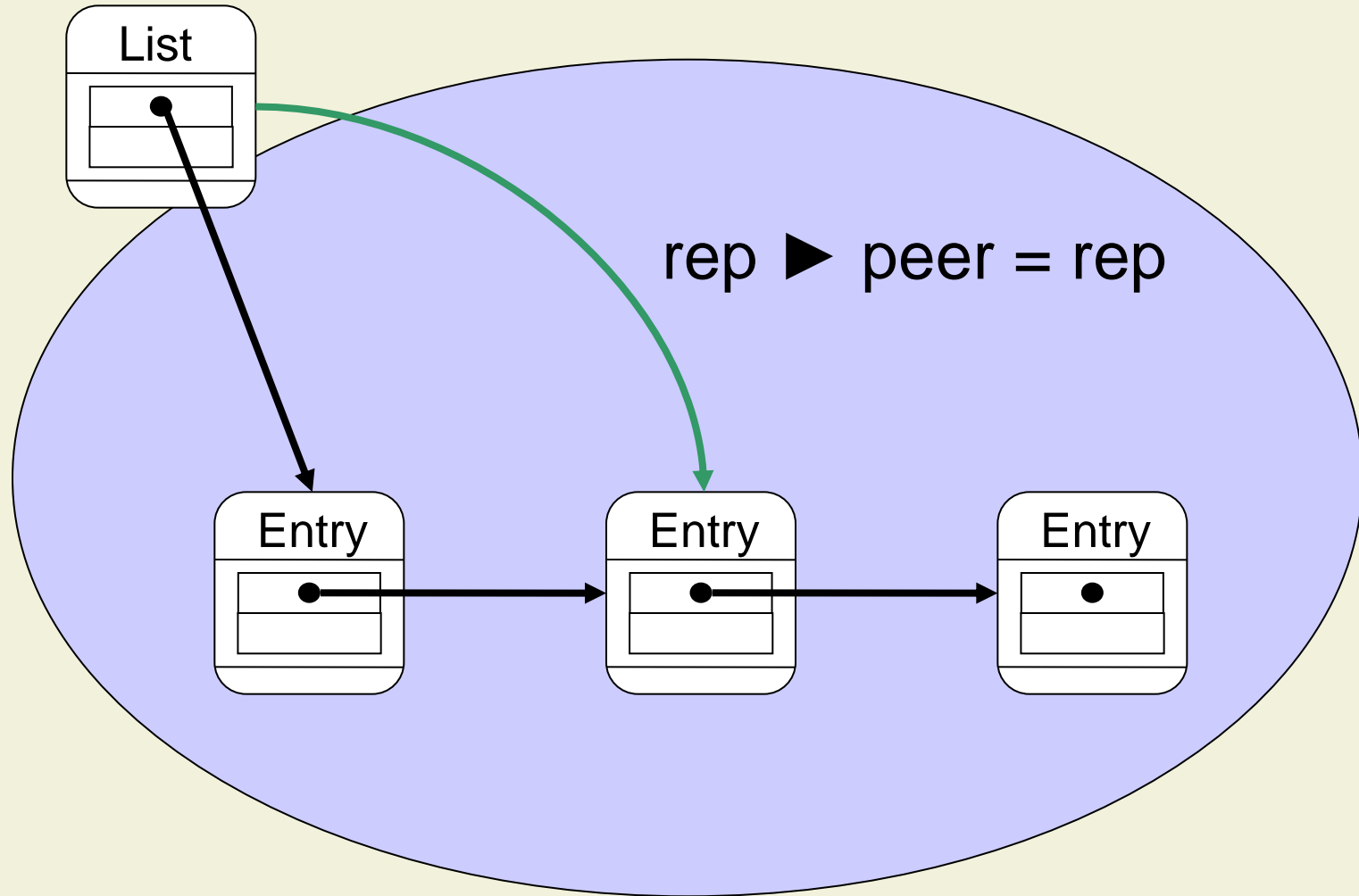
Ownership information  
is relative to **this**  
reference (viewpoint)

```
class Entry {  
  private any Object element;  
  private peer Entry previous, next;  
  public Entry( any Object o, peer Entry p, peer Entry n ) { ... }  
}
```

# Viewpoint Adaptation: Example 1



# Viewpoint Adaptation: Example 2



# Viewpoint Adaptation

►	<i>peer T</i>	<i>rep T</i>	<i>any T</i>
<i>peer S</i>	<i>peer T</i>	?	<i>any T</i>
<i>rep S</i>	<i>rep T</i>	?	<i>any T</i>
<i>any S</i>	?	?	<i>any T</i>

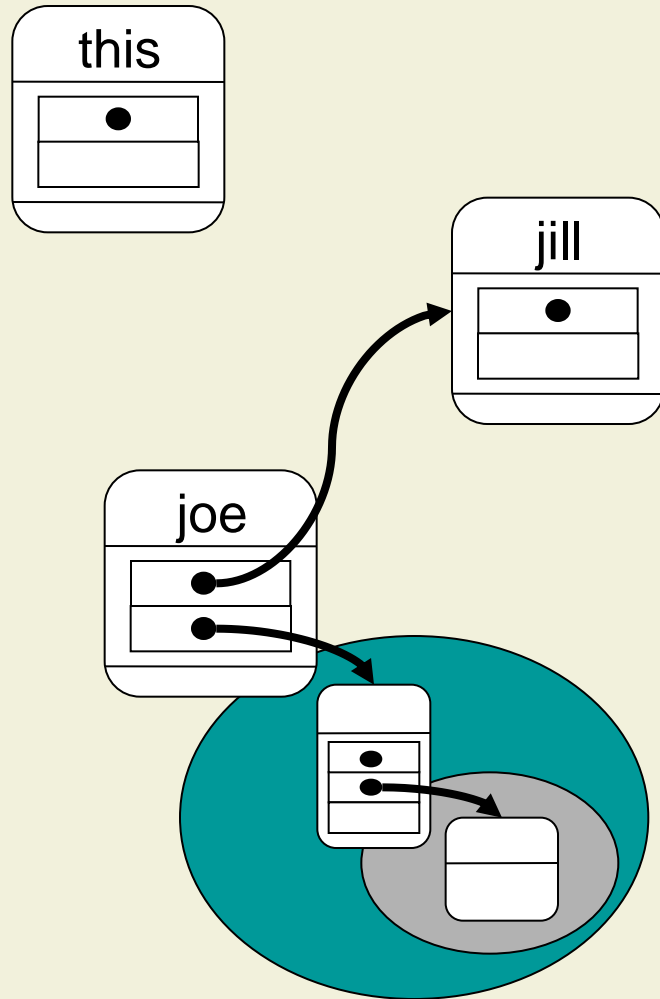
$$v = e.f;$$

$$\tau(e) \blacktriangleright \tau(f) \leq \tau(v)$$

$$e.f = v;$$

$$\tau(v) \leq \tau(e) \blacktriangleright \tau(f)$$

# Read vs. Write Access



```
class Person {  
    public rep Address addr;  
    public peer Person spouse;  
    ...  
}
```

```
peer Person joe, jill;
```

```
joe.spouse = jill;
```

```
any Address a = joe.addr;
```

```
joe.addr = new rep Address( );
```

# The lost Modifier

- Some ownership relations **cannot be expressed** in the type system
- Internal modifier **lost** for fixed, but unknown owner
- Reading locations with lost ownership is allowed
- Updating locations with lost ownership is unsafe

```
class Person {  
    public rep Address addr;  
    public peer Person spouse;  
    ...  
}
```

```
peer Person joe, jill;
```

```
joe.spouse = jill;
```

**lost** Address

```
any Address a = joe.addr;
```

```
joe.addr = new rep Address( );
```

**lost** Address

# The lost Modifier: Details

►	<i>peer T</i>	<i>rep T</i>	<i>any T</i>
<i>peer S</i>	<i>peer T</i>	<i>lost T</i>	<i>any T</i>
<i>rep S</i>	<i>rep T</i>	<i>lost T</i>	<i>any T</i>
<i>any S</i>	<i>lost T</i>	<i>lost T</i>	<i>any T</i>
<i>lost S</i>	<i>lost T</i>	<i>lost T</i>	<i>any T</i>

Another  
existential type

- Subtyping
  - $\text{rep } T <: \text{lost } T$
  - $\text{peer } T <: \text{lost } T$
  - $\text{lost } T <: \text{any } T$

# Type Rules: Field Access

- The field read

$$v = e.f;$$

is correctly typed if

- $e$  is correctly typed
- $\tau(e) \blacktriangleright \tau(f) <: \tau(v)$

- The field write

$$e.f = v;$$

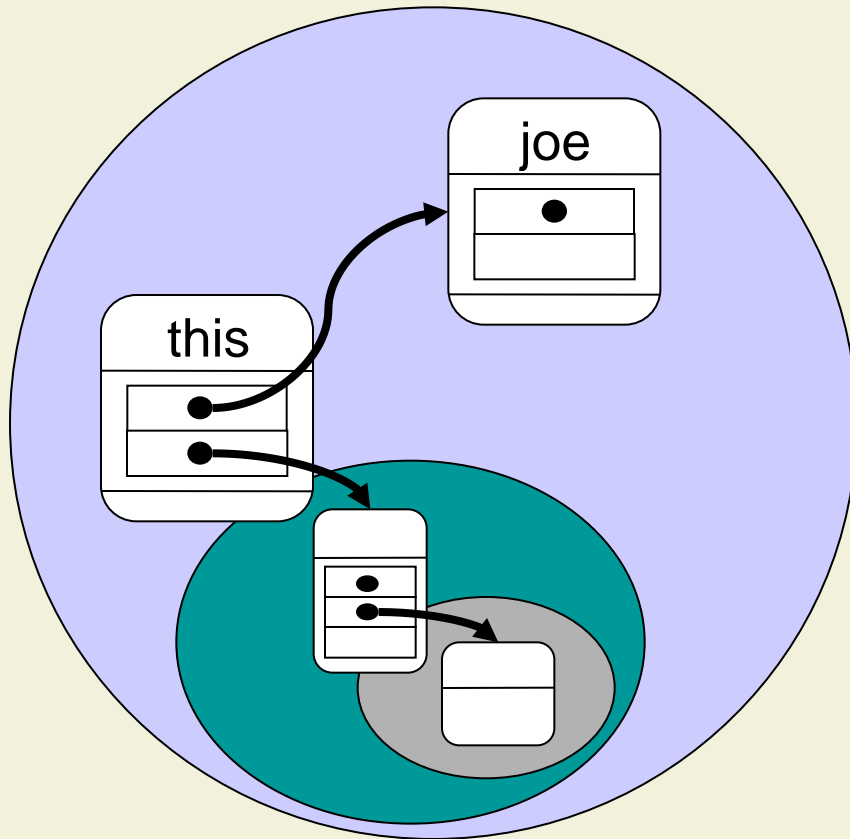
is correctly typed if

- $e$  is correctly typed
- $\tau(v) <: \tau(e) \blacktriangleright \tau(f)$
- $\tau(e) \blacktriangleright \tau(f)$  does not have **lost** modifier

- Analogous rules for method invocations
  - Argument passing is analogous to field write
  - Result passing is analogous to field read



# The self Modifier



```
class Person {  
    public rep Address addr;  
    public peer Person spouse;  
    ...  
}
```

```
peer Person joe;
```

```
joe.addr = new rep Address( );
```

```
this.addr = new rep Address( );
```

- Internal modifier **self** only for the **this** literal

# The self Modifier: Details

►	<i>peer T</i>	<i>rep T</i>	<i>any T</i>
<i>peer S</i>	<i>peer T</i>	<i>lost T</i>	<i>any T</i>
<i>rep S</i>	<i>rep T</i>	<i>lost T</i>	<i>any T</i>
<i>any S</i>	<i>lost T</i>	<i>lost T</i>	<i>any T</i>
<i>lost S</i>	<i>lost T</i>	<i>lost T</i>	<i>any T</i>
<i>self S</i>	<i>peer T</i>	<i>rep T</i>	<i>any T</i>

- Subtyping

- self T* <: *peer T*

$v = e.f;$

 $\tau(e) \blacktriangleright \tau(f) <: \tau(v)$ 

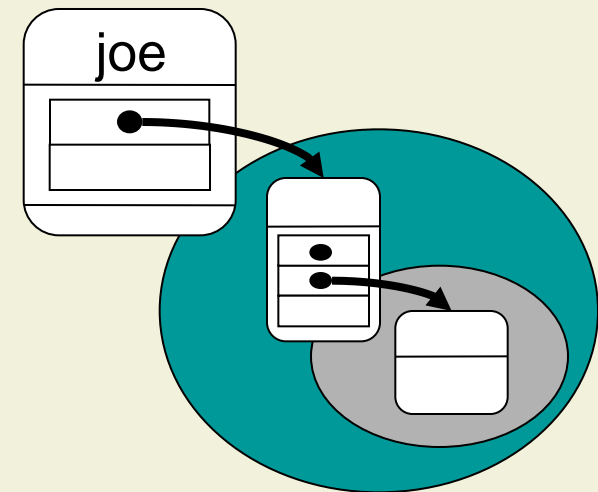
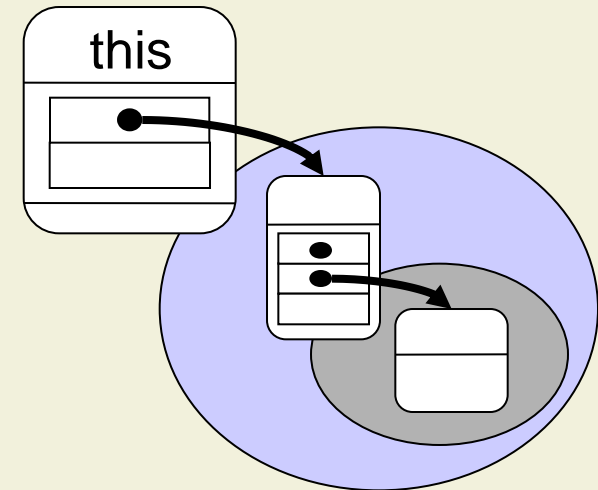
$e.f = v;$

 $\tau(v) <: \tau(e) \blacktriangleright \tau(f)$ 
 $\tau(e) \blacktriangleright \tau(f)$  does not have **lost** modifier

# Example: Sharing

```
class Person {  
  public rep Address addr;  
  ...  
}
```

- Different Person objects have different Address objects
  - No unwanted sharing



# Example: Internal vs. External Objects

```
class Person {  
  private rep Address addr;  
  
  public rep Address getAddr( ) {  
    return addr;  
  }  
  
  public void setAddr( rep Address a ) {  
    addr = a;  
  }  
  
  public void setAddr( any Address a ) {  
    addr = new rep Address( a );  
  }  
}
```

Address is part of  
Person's internal  
representations

Clients receive a  
lost-reference

Cannot be called  
by clients

Cloning  
necessary

# Internal vs. External Objects (cont'd)

```
class Person {  
  private any Company employer;  
  
  public any Company getEmployer( ) {  
    return employer;  
  }  
  
  public void setEmployer( any Company c ) {  
    employer = c;  
  }  
}
```

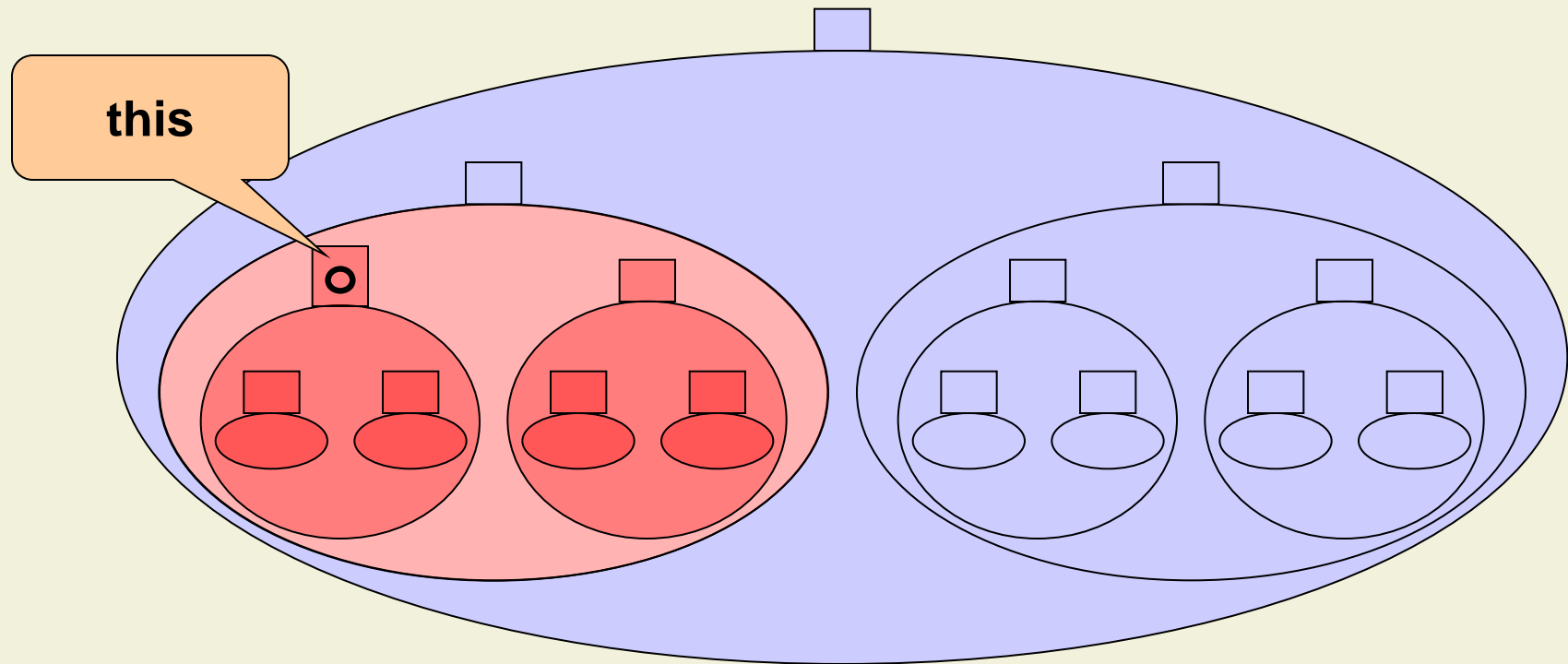
Company is shared  
between many  
Person objects

Can be called  
by clients

# Owner-as-Modifier Discipline

- Based on the topological type system we can strengthen encapsulation with extra restrictions
  - Prevent modifications of internal objects
  - Treat **any** and **lost** as readonly types
  - Treat **self**, **peer**, and **rep** as readwrite types
- Additional rules enforce owner-as-modifier
  - Field write  $e.f = v$  is valid only if  $\tau(e)$  is **self**, **peer**, or **rep**
  - Method call  $e.m(\dots)$  is valid only if  $\tau(e)$  is **self**, **peer**, or **rep**, or called method is **pure**

# Owner-as-Modifier Discipline (cont'd)



- A method may modify only objects directly or indirectly owned by the owner of the current **this** object

# Internal vs. External Objects Revisited

```
class Person {  
  private rep Address addr;  
  private any Company employer;  
  
  public rep Address getAddr( ) { return addr; }  
  public void setAddr( any Address a ) {  
    addr = new rep Address( a );  
  }  
  
  public any Company getEmployer( ) { return employer; }  
  public void setEmployer( any Company c ) { employer = c; }  
}
```

Company is shared;  
cannot be modified

Clients receive  
(transitive)  
readonly reference

Accidental capturing  
is prevented



# (Simplified) Programming Discipline

## ■ Rule 1: No Role Confusion

Different types for different roles

- Expression with one alias mode must not be assigned to variables with another mode, except to an **arg**-variable

## ■ Rule 2: No Representation Exposure

Viewpoint adaptation for **rep** types

- **rep**-mode must not occur in an object's interface
- Methods must not take or return **rep**-objects
- Fields with **rep**-mode may only be accessed on **this**

## ■ Rule 3: No Argument Dependence

- Implementations must not depend on the state of argument objects

Like with programming discipline

# Achievements

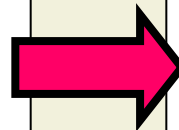
- **rep** and **any** types enable **encapsulation of whole object structures**
- Encapsulation **cannot be violated** by subclasses, via casts, etc.
- The technique **fully supports subclassing**
  - In contrast to solutions with final, private inner classes, etc.

```
class ArrayList {  
    protected rep int[ ] array;  
    private int next;  
    ...  
}
```

```
class MyList extends ArrayList {  
    public peer int[ ] leak( ) {  
        return array;  
    }  
}
```

# Exchanging Implementations

```
class ArrayList {  
  private int[ ] array;  
  private int next;  
  
  // requires ia != null  
  // ensures  $\forall i. 0 \leq i < \text{ia.length}:$   
  //           isElem( old( ia[ i ] ) )  
  public void setElems( int[ ] ia )  
    { array = ia; next = ia.length; }  
  
  ...  
}
```

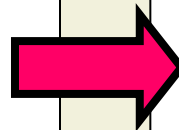


```
class ArrayList {  
  private Entry header;  
  
  // requires ia != null  
  // ensures  $\forall i. 0 \leq i < \text{ia.length}:$   
  //           isElem( old( ia[ i ] ) )  
  public void setElems( int[ ] ia )  
    { ... /* create Entry for each  
              element */ }  
  
  ...  
}
```

- Interface including contract remains unchanged

# Exchanging Implementations (cont'd)

```
class ArrayList {  
  private rep int[ ] array;  
  private int next;  
  
  // requires ia != null  
  // ensures  $\forall i. 0 \leq i < \text{ia.length}$ :  
  //           isElem( old( ia[ i ] ) )  
  public void  
  setElems( any int[ ] ia )  
  { System.arraycopy(...);  
    next = ia.length; }  
  ...  
}
```



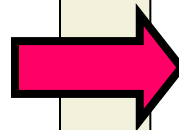
```
class ArrayList {  
  private rep Entry header;  
  
  // requires ia != null  
  // ensures  $\forall i. 0 \leq i < \text{ia.length}$ :  
  //           isElem( old( ia[ i ] ) )  
  public void  
  setElems( any int[ ] ia )  
  { ... /* create Entry for each  
        element */ }  
  ...  
}
```

Accidental capturing  
is prevented

# Exchanging Implementations (cont'd)

```
class ArrayList {  
  private rep int[ ] array;  
  private int next;  
  
  public any int[ ] getElems( )  
  { return array; }  
  ...  
}
```

Leaking is still possible



```
class ArrayList {  
  private rep Entry header;  
  
  public void any int[ ] getElems( )  
  { /* create new array */ }  
  ...  
}
```

```
peer ArrayList list = new peer ArrayList( );  
list.prepend( 0 );  
any int[ ] ia = list.getElems( );  
list.prepend( 1 );  
assert ia[ 0 ] == 1;
```

- Observable behavior is changed

# Consistency of Object Structures

- Consistency of object structures depends on fields of several objects
- Invariants are usually specified as part of the contract of those objects that represent the interface of the object structure

```
class ArrayList {  
    private int[ ] array;  
    private int next;  
  
    // invariant array != null    &&  
    //  0<=next<=array.length  &&  
    //   $\forall i. 0 \leq i < \text{next}: \text{array}[i] \geq 0$   
  
    public void add( int i )    { ... }  
    public void setElems( int[ ] ia )  
        { ... }  
  
    ...  
}
```

# Invariants for Object Structures

- The invariant of object *o* **may depend on**
  - Encapsulated fields of *o*
  - Fields of objects (transitively) owned by *o*
- Interface objects have **full control** over their rep-objects

```
class ArrayList {  
    private rep int[ ] array;  
    private int next;  
  
    // invariant array != null    &&  
    //  0<=next<=array.length  &&  
    //  ∀i.0<=i<next: array[ i ] >= 0  
  
    public void add( int i )  { ... }  
    public void setElems  
        ( any int[ ] ia )  { ... }  
  
    ...  
}
```

# Security Breach in Java 1.1.1

```
class Malicious {
```

```
  void bad( ) {
```

```
    Identity[ ] s;
```

```
    Identity trusted = java.Security...;
```

```
    s = Malicious.class.getSigners( );
```

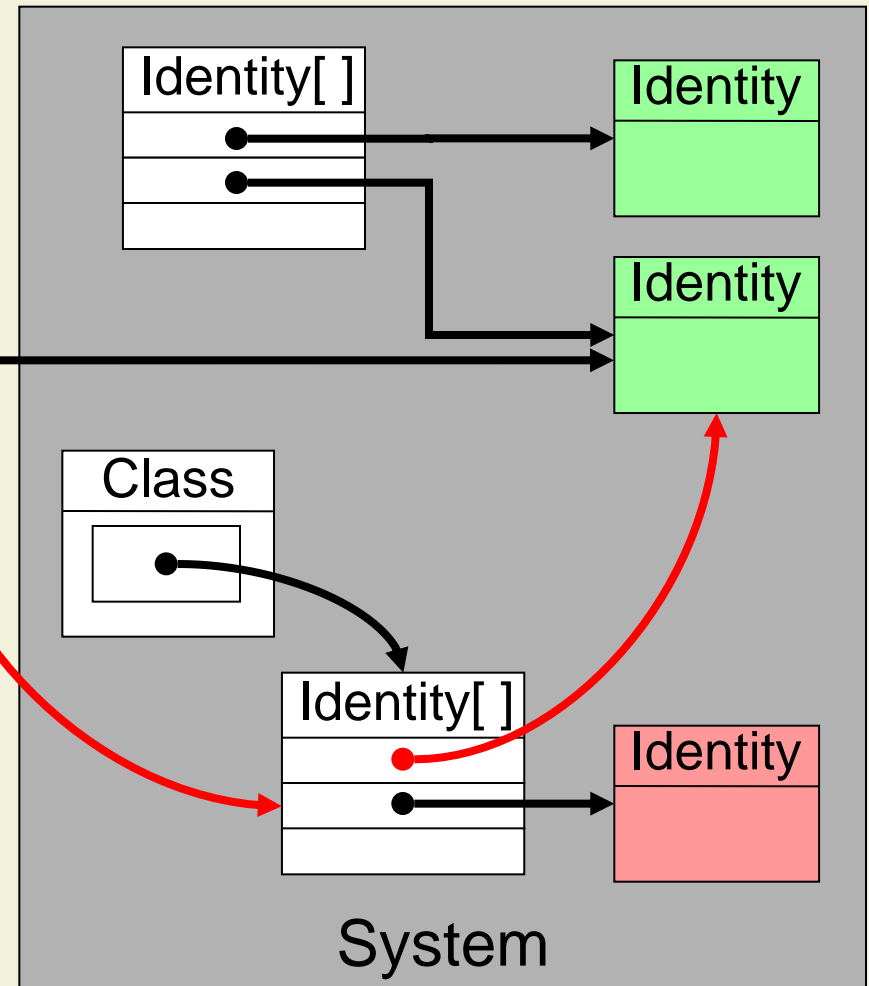
```
    s[ 0 ] = trusted;
```

```
    /* abuse privilege */
```

```
  }
```

```
}
```

Identity[ ] getSigners( )  
{ **return** signers; }





# Security Breach in Java 1.1.1 (cont'd)

```
class Malicious {
```

```
void bad( ) {
```

```
any Identity[ ] s;
```

```
Identity trusted = java.Security...;
```

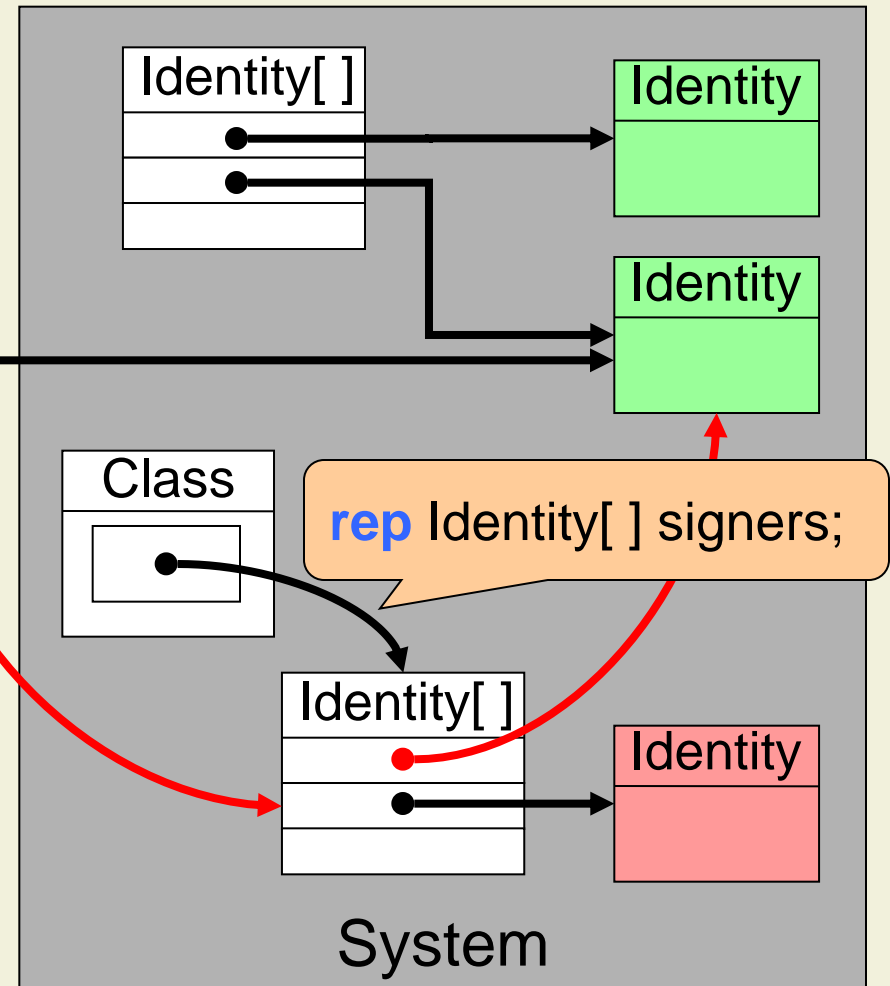
```
s = Malicious.class.getSigners( );
```

```
s[ 0 ] = trusted;
```

```
}
```

```
}
```

**rep** Identity[ ] getSigners( )  
{ **return** signers; }



# Ownership Types: Discussion

- Ownership types express **heap topologies** and enforce **encapsulation**
- Owner-as-modifier is helpful to **control side effects**
  - Maintain object invariants
  - Prevent unwanted modifications
- Other applications also need **restrictions of read access**
  - Exchange of implementations
  - Thread synchronization
- Ownership types are an area of current research

# References

- Werner Dietl and Peter Müller: *Universes: Lightweight Ownership for JML*. Journal of Object Technology, 2005
- Werner Dietl, Sophia Drossopoulou, and Peter Müller: *Separating Ownership Topology and Encapsulation with Generic Universe Types*. ACM Trans. Program. Lang. Syst., 2011