

Exercise 4

Behavioral Subtyping and Inheritance

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Task 1

Consider the following Java code:

```
interface I {};  
  
class C {};  
  
public class Main {  
    public static void main(String[] argv) {  
        C c = new C();  
        I i = (I) c;  
    }  
}
```

Try to compile it. If it compiles, try to execute it. What happens? Why? Do you expect to see the same behavior if `I` were a class, instead of an interface?

solution

- If `I` is an *interface*: the compiler allows the code to go through although it cannot prove that `c` implements `I`. The reason is that there might be a subclass `D` of `C` such that `D` implements `I` and `c` might be an object of `D`. Here Java opts for the flexibility of dynamic type checking. When the code executes a runtime exception is thrown, because `c` does not implement `I` and this is caught by the runtime check.
- If `I` were a *class*: the code does not compile. Java does not support multiple inheritance, so it is not possible to have a subclass `D` of `C`, which also extends `I`.

Task 2

Consider the following Java classes:

```
class Number {  
    int n;  
  
    /// requires true  
    /// ensures n == p  
    void set(int p) {  
        n = p;  
    }  
}  
  
class UndoNaturalNumber extends Number {  
    int undo;
```

```

/// requires 0 < q
/// ensures n == q && undo == old(n)
void set(int q) {
    undo = n;
    n = q;
}

```

Is UndoNaturalNumber a behavioral subtype of Number? Consider that we are using specification inheritance. What are the effective pre/post-conditions of the method UndoNaturalNumber.set according to the rules from Slides 67 and 71?

— solution —

UndoNaturalNumber is not a behavioral subtype of Number, because it has a stronger precondition for the method set.

The effective precondition is: $\text{true} \parallel (0 < q)$, which is equivalent to true.

The effective postcondition is: $(\text{old}(\text{true}) \implies n == p) \ \&\& \ (\text{old}(0 < q) \implies n == q \ \&\& \ \text{undo} == \text{old}(n))$. Note that p and q refer to the same parameter, so the effective postcondition actually is: $(\text{old}(\text{true}) \implies n == q) \ \&\& \ (\text{old}(0 < q) \implies n == q \ \&\& \ \text{undo} == \text{old}(n))$. Since for parameters we always have $\text{old}(q) == q$, the effective postcondition is equivalent to $(n == q) \ \&\& \ (0 < q \implies \text{undo} == \text{old}(n))$.

Task 3

From a previous exam.

Assume the following types in Java:

```

enum Shift {DayShift, NightShift, SpecialShift}

interface PostalWorker {
    boolean sick();

    ///ensures sick()
    void catchDisease();

    ///requires when == SpecialShift || when == DayShift
    ///requires !sick()
    int work(Shift when);
}

interface Bartender {
    boolean sick();

    ///ensures sick()
    void catchDisease();

    ///requires when == SpecialShift || when == NightShift
    ///requires !sick()
    int work(Shift when);
}

```

The work() method can be called to request the corresponding person to work the specified shift. The value returned by work() is the average hourly wage that was earned during the working shift including tips.

A) Now we introduce another interface:

```

interface HardWorker extends PostalWorker, Bartender {
    ///requires true
    int work(Shift when);
}

```

Assuming the improved rule for specification inheritance discussed in the course, what is the effective precondition of the `work()` method of the `HardWorker` interface?

— solution —

```

///requires
    (!sick() && (when == SpecialShift || when == DayShift))
|| (!sick() && (when == SpecialShift || when == NightShift))
|| true

```

which is equivalent to

```

///requires true

```

B) Now we add postconditions to all `work()` methods. Everything else remains as before.

```

interface PostalWorker {
    ...
    ///ensures result ≥ 15 && result ≤ 25
    int work(Shift when);
}

interface Bartender {
    ...
    ///ensures result ≥ 20 && result ≤ 30
    int work(Shift when);
}

interface HardWorker extends PostalWorker, Bartender {
    ...
    ///ensures result ≥ 25 && result ≤ 50
    int work(Shift when);
}

```

Assuming the improved rules for specification inheritance, what is the effective postcondition of the `work()` method of `HardWorker`?

— solution —

```

///ensures
    ( old(!sick() && (when == SpecialShift || when == DayShift))
    ⇒ (result ≥ 15 && result ≤ 25) )
&& ( old(!sick() && (when == SpecialShift || when == NightShift))
    ⇒ (result ≥ 20 && result ≤ 30) )
&& ( old(true)
    ⇒ (result ≥ 25 && result ≤ 50) )

```

which is equivalent to

```

///ensures
    ( old(!sick() && when != NightShift)
    ⇒ result == 25 )
&& ( old(!sick() && when == NightShift)
    ⇒ (result ≥ 25 && result ≤ 30) )
&& ( old(sick())
    ⇒ (result ≥ 25 && result ≤ 50) )

```

C) Consider the following code:

```
///requires worker != null
///requires !worker.sick()
int foo(HardWorker worker) {
    return worker.work(Shift.SpecialShift);
}
```

What is the range of possible return values of the method `foo()`?

— solution —

Only 25 is a possible return value.

D) Change the body of the method `foo()` such that it calls the `work()` method of `worker` in a way that makes it possible for this call to return 50.

— solution —

```
int foo(HardWorker worker) {
    worker.catchDisease();
    return worker.work(Shift.SpecialShift);
}
```

Task 4

From a previous midterm.

Imagine extending the syntax of the Java language to support the following keywords:

- `subtypes`: used to declare that a class is a subtype of another class (without inheritance)
- `inherits`: used to declare that a class inherits from another class (without subtyping)

Now consider the following classes:

```
class A {
    public int foo (int n) { return n - 1; }
}

class B {
    public int foo (int n) { return n + 1; }
    public int bar (int n) { return foo(n) - 1; }
}

class C inherits A subtypes B {
    public int bar (int n) { return foo(n); }
}

class Main {
    public static void main(String[] args) {
        B b = new C();
        System.out.println( b.bar(3) );
    }
}
```

What would happen if we tried to compile the code and execute the method `main` from the class `Main`?

- (a) The code will be rejected by the compiler

- (b) The code will compile but the execution will fail
- (c) **CORRECT:** The code will compile and print 2
- (d) The code will compile and print 4
- (e) None of the above

Task 5

Consider two classes `Stack` and `Queue`, implementing the standard LIFO/FIFO data structures, both of which have methods with the following signatures:

```
void push(Object o);
Object pop();
bool isEmpty();
int size();
void reverse();
```

A) Despite having identical signatures, these two classes cannot be behavioral subtypes of one another. Why not?

— solution —

The intended behavior is that a `Stack` is LIFO, while a `Queue` is FIFO. Therefore, the `pop` and `push` have different behavior and so neither can be considered a behavioral subtype of the other.

B) When implementing these two classes, is there any possibility of code reuse? If so, give details.

— solution —

Depending on the internal representation, either the `pop()` or the `push()` method (but not both) could be reused, from one implementation to the other. For example, if one implements a `Queue` by pushing to the end of a linked list, and popping from the beginning, then a `Stack` could be implemented either by pushing on the beginning of the list and reusing the `pop()` method, or by reusing the `push()` method and popping from the end of the list. Furthermore, it's likely that the `isEmpty()`, `size()` and `reverse()` methods could all be reused.

C) Describe at least one way of reusing the code in one class by the other - which programming language features are needed for this to work?

— solution —

Any mechanism which allows code reuse without subtyping, e.g., private inheritance in C++ or aggregation.

Another option would be to have a “common super class” used by both implementations. This super-class, however, would either be too wide (allowing insertion/removal at both ends) or rather thin (allowing only insertion on one side). In the wide case we could use a kind of linked list, for example, that can insert/remove at the beginning and end, and use private inheritance to expose only the relevant operations to the clients of each data structure.

Task 6

Suppose that we have a database, for which we want an “automated key generation” feature. This means that each time the user inserts a new tuple, a unique key is automatically generated for the tuple by the system. An obvious way to do that is to write a counter, which increments by 1 the value that it returns each time it is called. The method that generates a new key is called `generate`.

A) Write a Java class `IncCounter` and an accompanying specification for such a counter.

— solution —

```
class IncCounter {
  /// constraint old(key) <= key
  int key;

  IncCounter () { key = 0; }

  /// ensures (key == old(key) + 1) ∧ (result == old(key))
  int generate () { return key++; }
}
```

B) Annotate the following Java class with specifications and show that it is not a behavioural subtype of `IncCounter`.

```
class DecCounter {
  int key;
  DecCounter () { key = 0; }
  int generate () { return key--; }
}
```

— solution —

The postcondition which precisely describes the behavior of the method `DecCounter.generate` is $(key == \text{old}(key) - 1) \wedge (\text{result} == \text{old}(key))$. This postcondition does not refine the postcondition of `IncCounter.generate`. The history constraint is $\text{old}(key) \geq key$ and also does not strengthen the one of `IncCounter`.

C) Write an abstract class `GenerateUniqueKey` together with a specification, such that both `IncCounter` and `DecCounter` are behavioural subtypes of `GenerateUniqueKey`. In the specification, you may use helper methods and fields.

— solution —

The abstract parent class can be declared using a helper pure method `boolean used(int)`. Informally, the helper method returns `true` if `x` has been used as a key before. Furthermore, the correctness of the class relies on the property that once a number is used, it never becomes unused again. This can be expressed with a two-state history constraint.

The definitions of the classes follow:

```
abstract class GenerateUniqueKey {
  /// constraint  $\forall x:\text{int} \mid (\text{old}(\text{used}(x)) \Rightarrow \text{used}(x))$ 
  abstract boolean used(int);

  /// ensures  $\neg \text{old}(\text{used}(\text{result})) \wedge \text{used}(\text{result})$ 
  abstract int generate ();
}
```

```
class IncCounter { // ... and similarly for DecCounter
  int key;
  IncCounter () { key = 0; }

  boolean used (int x) { return x < key; }

  /// ensures key == old(key) + 1  $\wedge$  result == old(key)
  int generate () { return key++; }
}
```