

Exercise 10

Type Erasure, Templates and Information Hiding

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Task 1

Consider the following Java method:

```
String concatenate(List<?> list) {
    String result = "";
    String separator = "";
    if(list instanceof List<String>) {
        result = "String:";
        separator = " ";
    }
    else if(list instanceof List<Integer>) {
        result = "Integers:";
        separator = "+";
    }
    for(Object el : list)
        result = result + separator + el.toString();
    return result;
}
```

A) This program is rejected by the Java compiler. Why?

— solution —

The Oracle and the Open JDK compilers both produce these short errors:

```
illegal generic type for instanceof
illegal generic type for instanceof
```

The Eclipse compiler tries to be more helpful:

```
Cannot perform instanceof check against parameterized type
List<String>. Use the form List<?> instead since further
generic type information will be erased at runtime
```

```
Cannot perform instanceof check against parameterized type
List<Integer>. Use the form List<?> instead since further
generic type information will be erased at runtime
```

This happens because of type erasure in Java.

B) Using the advice given by the Eclipse Java compiler (replace `List<...>` with `List<?>`), rewrite and compile the program. What are the results of executing the method passing each of the following:

- A list of strings containing only one element "word"?
- A list of Integers containing only one element `Integer(1)`?
- A list of Objects containing only one element (initialized by `new Object()`)?

— solution —

First of all, we follow the output of the compiler, and so we rewrite the method to:

```
String concatenate(List<?> list) {
    String result = "";
    String separator = "";
    if(list instanceof List<?>) {
        result = "String:";
        separator = " ";
    }
    else if(list instanceof List<?>) {
        result = "Integers:";
        separator = "+";
    }
    for(Object el : list)
        result = result + separator + el.toString();
    return result;
}
```

The Java compiler will compile this program without any warning. The output of the method is:

```
String: word
String: 1
String: java.lang.Object@3e25a5
```

C) Is this behaviour consistent with what you would expect from the initial program? If not, how can you fix it?

— solution —

No, in the original program we expected:

```
String: word
Integers:+1
java.lang.Object@3e25a5
```

We can try to fix it in the following way:

```
String concatenate(List<?> list) {
    String result = "";
    String separator = "";
    if(list.size() >= 1 )
        if(list.get(0) instanceof String) {
            result = "Strings:";
            separator = " ";
        }
        else if(list.get(0) instanceof Integer) {
```

```

        result = "Integers:";
        separator = "+";
    }
    for(Object el : list)
        result = result + separator + el.toString();
    return result;
}

```

But this requires to have at least one element in the list. Moreover, there is no guarantee that if the first element is, for example, a String, that this is not a list of Objects. Therefore, an improved solution would be to iterate over all the elements of the list and to compute their smallest common supertype.

D) What would happen if you tried to implement the different cases using method overloading instead of just one method? Why is this the case?

— solution —

If we introduce separate methods which differ only by the generic types of their arguments, we get compile-time errors such as:

```
Method concatenate(List<? extends Object>) has the same
erasure concatenate(List<E>) as another method in type C
```

This restriction is imposed to ensure that when choosing which of the overloaded method definitions to call, we always have a “best fit”. Java class files do however include generic versions of the method signatures in the class (to enable separate compilation and type-checking of generic code). For this reason, it might seem surprising that we cannot disambiguate between these different overloaded methods, since at compile-time the type information is all available. However, Java also supports raw types - versions of generic classes in which no type parameter is provided (e.g., `List` for a `List<X>` class). These are supported for backwards compatibility with pre-generics Java code. For this reason, we need to consider the possibility that a client calling our method provides an argument of raw type `List`. In this case, we would not be able to choose between our different method overloads.

E) What happens if you compile and execute the initial program in C# ? Why? (Assume that we replace the wildcard by a method type parameter `T` to make it work in C#.)

— solution —

The program is compiled and we obtain the expected results (“String: word”, “Integers:+1”, “...”), since in C# there is no type erasure and the information about generics is preserved at runtime.

Task 2

From a previous exam

Consider the following Java program, which compiles correctly and makes use of generics:

```

1 class Animal {}
2 class Mammal extends Animal {}
3 class Tiger extends Mammal {}
4

```

```

5 class Ship<T extends Animal> {
6     public T content;
7 }
8
9 class Cage<T extends Mammal> {
10    public T content;
11
12    void takeFromShip(Ship<T> other) {
13        this.content = other.content;
14        other.content = null;
15    }
16 }
17
18 class Zoo<T extends Mammal> {
19    void swapTigers(Ship<T> mammalShip, Cage<Tiger> tigerCage) {
20        Tiger tiger = tigerCage.content;
21        Cage<Tiger> tmpCage = new Cage<Tiger>();
22        tmpCage.takeFromShip((Ship<Tiger>) mammalShip);
23        mammalShip.content = (T) tiger;
24        tigerCage.content = tmpCage.content;
25    }
26 }

```

A) List all the typecasts that the virtual machine will perform at runtime, when executing the methods `takeFromShip` and `swapTigers`. For each cast, write at which line number in the original program it is performed, and what expression is cast to which type. Do not optimize away casts that are statically known to succeed.

— solution —

- At line 13: (Mammal) `other.content`
- At line 20: (Tiger) `tigerCage.content`
- At line 22 (Ship) `mammalShip`
- At line 23: (Mammal) `tiger`

B) For each of the following two methods (from B.1 and B.2), write if they would compile without errors if added to the class `Cage`. If they do not compile, briefly explain why.

B.1

```

Cage<Tiger>[] getTigers(int number) {
    Cage<Tiger>[] cages = new Cage<Tiger>[number];
    for (int i = 0; i < number; i++) {
        cages[i] = new Cage<Tiger>();
        cages[i].content = new Tiger();
    }
    return cages;
}

```

B.2

```

int numCageFields() {
    Class cl = Cage<Animal>.class;
    return cl.getFields().length;
}

```

solution

1. Compiler error: array of generic types not allowed.
2. Compiler error: class object of generic types not available, or the type parameter does not respect the constraint of class Cage.

Task 3

A C++ template class can inherit from its template argument:

```
template <typename T>
class SomeClass : public T { ... }
```

A) Using this technique and given the following class definition

```
class Cell {
public:
    virtual void setVal(int x) { x_ = x; }
    virtual int value() { return x_; }
private:
    int x_;
}
```

write two template classes that can be used as “mixins” for the class Cell:

- Doubling - doubles the value stored in the cell.
- Counting - counts the number of times the value of the cell was read.

Do not use multiple inheritance. It should be possible to use the classes like this:

```
auto c = new Doubling<Counting<Cell>>(); // instantiation
c->setVal(5);
c->value(); // returns 10
c->numRead(); // returns 1
```

solution

```
template <typename T> class Doubling : public T {
public:
    virtual void setVal(int x) override {
        T::setVal(x * 2);
    }
}

template <typename T> class Counting : public T {
public:
    virtual int value() override {
        ++numRead_;
        return T::value();
    }
    int numRead() { return numRead_; }
private:
    int numRead_;
}
```

B) Describe how the instantiation above will look like.

— solution —

When the mix-ins are instantiated the following two classes will be generated:

```
class CountingCell : public Cell {  
public:  
    virtual int value() override {  
        ++numRead_  
        return Cell::value();  
    }  
    int numRead() { return numRead_; }  
private:  
    int numRead_  
}  
  
class DoublingCountingCell : public CountingCell {  
public:  
    virtual void setVal(int x) override {  
        CountingCell::setVal(x * 2);  
    }  
}
```

C) How does this concept of mixins in C++ differ from Scala traits?

— solution —

While this concept is similar to Scala traits there are some notable differences. In Scala it is possible to mix any number of traits in a class and use this in any location of the code that requires the same class and a subset of the traits:

```
var x = new X with A with B with C with D  
var x1: (X) = x // OK  
var x2: (X with A) = x // OK  
var x3: (X with B) = x // OK  
var x4: (X with A with D with C) = x // OK
```

Using the proposed solution in C++ however is more restrictive, as there is no way to refer to the class X with arbitrary mix-ins:

```
auto x = new D<C<B<A<X>>>>();  
X* x1 = x; // OK  
A<X>* x2 = x; // OK  
B<X>* x3 = x; // Does not compile  
C<D<A<X>>>* x4 = x; // Does not compile
```

This is particularly important for traits that introduce new methods like Counting.numRead() since any client code that uses this new behavior would have to know exactly how the trait was mixed-in.

Another problem of the C++ solution is object construction. If the base class does not have a default constructor then the mix-ins should know to call the correct constructor and provide appropriate parameters. An alternative here is for the mixin to just inherit the base class constructors: using T::T; which will allow clients of the mixin to use all constructor available in the base class. This works fine if the state of the mixin can be initialized with default values.

A further difference to Scala is that in the C++ solution it is possible to include the same “trait” more than once:

```
auto x = new Doubling<Doubling<X>>();  
x->setVal(5);
```

```
x->value(); // returns 20
```

An advantage of the C++ solution is that we do not need to declare the base class that the mix-ins extend. Thus it is possible to use them with different base classes as long as they have matching virtual methods.

Task 4

The type correctness of a C++ template class is checked only when the template is instantiated. This makes it difficult to develop templates modularly. We can try to make templates more modular by extending C++ with a new way to declare type arguments:

```
template<T s_extends SomeClass>
class TemplateClass {...}
```

Here T is the template argument and `SomeClass` is the name of a class which is an upper type bound for T . A template defined in this way may only be instantiated with a class T that is a *structural* subtype of `SomeClass`. Assume that the type checker checks such a template *definition* without having any concrete instantiation, under the assumption that T is a structural subtype of `SomeClass`.

This new feature is the only place where we introduce structural subtyping in C++, all other subtype relations in the language remain nominal as usual. Assume in general for any subtyping mode that method argument types are contravariant and method return types are covariant. Also, assume that all the methods are public and virtual.

A) Provide a declaration of the `Operation` class such that the class `Compose` can be type-checked before it is instantiated.

```
template<T s_extends Operation , U s_extends Operation>
class Compose : public Operation {
    public:
        T* t;
        U* u;
        int compute(int x) {
            return t->compute(u->compute(x));
        }
}
```

— solution —

```
class Operation {
    int compute(int x);
}
```

B) We also allow template parameters to occur as type arguments in upper bounds of the same template:

```
template<T s_extends Bound<T>>
class TemplateClass{...}
```

The above limits the possibilities for T to only structural subtypes of `Bound<T>`.

Consider the classes below:

```
class A : { void foo(A* a); };
class B : public A { B* bar(); };
```

```

class C : public B    {};

template <class T>
class FOO {
    void foo(T* t){...}
};

template <T s_extends FOO<T>>
class X { ... };

template <class T>
class BAR {
    T* bar(){...}
};

template <T s_extends BAR<T>>
class Y { ... };

```

Which of the following instantiations typecheck:

X
X<C>
Y
Y<C>

Explain why each combination does or does not typecheck.

— solution —

X can be instantiated with both classes:

- B: B.foo(A*) overrides FOO.foo(B*)
- C: C.foo(A*) overrides FOO<C>.foo(C*)

Y can be instantiated only with B:

- B: B* B.bar() overrides B* BAR.bar()
- C: B* C.bar() does not override C* BAR<C>.bar(), therefore C is not a structural subtype of BAR<C>

C) As a bound we also allow the template that is being declared:

```

template <T s_extends X<T>>
class X {
    int foo(T* t) {...}
}

```

Let the class A be:

```

class A {};

```

- Write an implementation of the body of the foo method of X such that X typechecks with the bound above (T s_extends X<T>) and also typechecks if the bound is changed to T s_extends A.
- Write an implementation of the body of the foo method of X such that X typechecks with the bound above (T s_extends X<T>), but does not typecheck if the bound is changed to T s_extends A.
- Write a class B that can be used to instantiate X.

— solution —

- `int foo(T* t) { return t ? 1 : 0; }`
- `int foo(T* t) { return t->foo(t); }`
- `class B { int foo(B* b) { return 0; } }`

D) A C++ template class can inherit from its template argument:

```
template <class T>
class Mixin : public T { ... }
```

Such a template is called a mixin. We want to use the newly introduced template bound feature `<T s_extends ...>` in order to create a mixin that is guaranteed only to override existing methods but not introduce new ones. Show how this can be done.

— solution —

```
template <T s_extends Mixin<T>>
class Mixin : public T
```

Here we restrict the base class `T` to be a structural subtype of `Mixin<T>`. Thus all admissible base classes `T`, have at least the methods that `Mixin<T>` defines in its body.

If the mixin tried to add a new method, this would fail. For example:

```
template <T s_extends Mixin<T>> class Mixin : public T {
    public: int foo(int x) { return x + T::foo(x); } // overridden method
           int bar(int y) { return 2 + y; } // newly-added method
}

class A { public: int foo(int x) { return 2 * x; } }
```

The instantiation `Mixin<A>` fails, because `A` is not a structural subtype of `Mixin<A>`, since `A` does not have the `bar` method.

Task 5

Suppose that we have a language with the information hiding rules of Java, but with structural subtyping. What should the subtyping relations between the following three classes be?

```
class A { int foo(); }
class B { protected int foo(); }
class C { public int foo(); }
```

— solution —

The subtyping relations are as follows: `C <: B <: A`.

In a language with structural subtyping, methods and fields of subclasses should be more accessible than those of the superclasses. For access modifiers, this means that methods with more permissive modifiers may override methods with less permissive modifiers.

Task 6

From a previous exam

Consider the following Java program consisting of two packages:

```

1  package A;
2
3  public abstract class Person {
4      _____ int tickets = 0;
5      _____ final int maxTickets = 3;
6
7      _____ abstract void buy(int t);
8  }
9
10 public class Buyer extends Person {
11     _____ void inc(int t) {
12         if (this.tickets + t <= this.maxTickets) this.tickets += t;
13     }
14     _____ void buy(int t) { if (t >= 0) inc(t); }
15 }
16
17
18
19 package B;
20 import A.*;
21
22 public class SmartBuyer extends Buyer {
23     _____ void inc(int t) { this.tickets += t; }
24 }
25
26 public class Main {
27     public static void main(String args[]) {
28         Buyer b = new SmartBuyer();
29         b.buy(9);
30     }
31 }

```

Provide the *most restrictive* access modifiers for the fields `tickets` and `maxTickets` and the methods `inc()` and `buy()` such that the program is still accepted by the compiler.

— solution —

The field `tickets` must be protected (since we need to access it from the class `SmartBuyer` which belongs to another package). The field `maxTickets` must have a default access modifier (because we need to access it from the class `Buyer` which belongs to the same package). The method `inc()` can be declared private in both `Buyer` and `SmartBuyer`. The method `buy()` in class `Person` must have a default access modifier (because abstract methods cannot be private), while the method `buy()` in class `Buyer` must be public (because we need to access it from the class `Main` which belongs to another package and is not a subclass of `Buyer`).