Formal Methods and Functional Programming

Modelling

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The slides in this section are partly based on the course *Automata-based System Analysis* by Felix Klaedtke

Example 1: Protocol Verification

- Protocol for file access with primitives `open`, `close`, and `write`

- Task: verify that a program obeys the following (informal) rules:

  A. All opened files must be closed eventually
  B. An opened file must be closed before the next open and vice versa
  C. All files must be opened before executing a write operation

  There is only one file, which is initially closed

- Problem is typical for verification of protocols
  - Locking (acquire, access, release)
  - Authentication (authenticate, access)

Example 1: Encoding in IMP

- File is represented by variable `f`
  - Write is encoded by assignment to `f`
  - Variable `o` counts how often file was opened/closed

- Encoding of primitives:
  - `open`: `o:=o` $+ 1$
  - `close`: `o:=o` $- 1$
  - `write`: `f:=`e

- Informal rules:

  A. After `o` has been set to one, it must eventually be re-set to zero
  B. In all execution states, `o` is zero or one
  C. When `f` is being assigned to, `o` must be greater than zero

  Variable `o` is initially zero

Example 1: Specification in NS and Hoare Logic

A. For a terminating program $\mathtt{s}$, $\mathtt{o}$ must be zero in the terminal state

$$\boxed{\langle \mathtt{s}, \sigma \rangle \to \sigma' \text{ and } \sigma(\mathtt{o}) = 0 \text{ then } \sigma'(\mathtt{o}) = 0}$$

$$\boxed{\{\ \mathtt{o} = 0\ \}\ \mathtt{s}\ \{\ \mathtt{o} = 0\ \}}$$

Property cannot be expressed for non-terminating programs

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$$\{ \; \mathtt{o} = 0 \; \} \; \mathtt{s} \; \{ \; \mathtt{o} = 0 \; \}$$

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- Natural semantics and Hoare logic can express properties of initial and terminal states, but not of intermediate states

Example 1: Specification in SOS (A)

- A: After `o` has been set to one, it must eventually be re-set to zero
  - For a terminating program `s`

$$\boxed{\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \sigma' \text{ and } \sigma(\mathtt{o}) = 0 \text{ then } \sigma'(\mathtt{o}) = 0}$$

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$$\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \sigma' \text{ and } \sigma(\mathtt{o}) = 0 \text{ then } \sigma'(\mathtt{o}) = 0$$

  - For a deterministic, non-terminating program s

$$\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \langle s', \sigma' \rangle \text{ and } \sigma(\mathtt{o}) = 0 \text{ and } \sigma'(\mathtt{o}) = 1 \text{ then there exist } s'', \sigma'' \text{ such that } \langle s', \sigma' \rangle \rightarrow_1^* \langle s'', \sigma'' \rangle \text{ and } \sigma''(\mathtt{o}) = 0$$

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  - For a non-deterministic, non-terminating program s

$$wc : \mathsf{Stm} \times \mathsf{State} \times \mathbb{N} \rightarrow \mathsf{Bool}$$

$$\begin{aligned} wc(s, \sigma, n) \Leftrightarrow\ & \sigma(o) = 0 \vee \\ & (\text{for all } s', \sigma' : \text{if } \langle s, \sigma \rangle \rightarrow_1 \langle s', \sigma' \rangle \text{ then there exists} \\ & m \in \mathbb{N} \text{ such that } m < n \text{ and } wc(s', \sigma', m)) \end{aligned}$$

$$\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \langle s', \sigma' \rangle \text{ and } \sigma(\mathtt{o}) = 0 \text{ and } \sigma'(\mathtt{o}) = 1 \text{ then there exists } n \in \mathbb{N} \text{ such that } wc(s', \sigma', n)$$

Example 1: Specification in SOS (B and C)

- B: In all execution states, o is zero or one

$$\boxed{\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \langle s', \sigma' \rangle \text{ and } \sigma(\mathtt{o}) = 0 \text{ then } \sigma'(\mathtt{o}) = 0 \text{ or } \sigma'(\mathtt{o}) = 1}$$

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- C: When `f` is being assigned to, `o` must be greater than zero
  - SOS cannot express properties of occurrences of statements
  - Work-around: Enrich state and make sure `write` changes the state

Example 1: Verification

A. For a terminating program $\texttt{s}$

$$\boxed{\langle \texttt{s}, \sigma \rangle \rightarrow_1^* \sigma' \text{ and } \sigma(\texttt{o}) = 0 \text{ then } \sigma'(\texttt{o}) = 0}$$

- Proof needs to consider all possible derivations of $\langle \texttt{s}, \sigma \rangle \rightarrow_1^* \sigma'$ to find all possible terminal states
- Problematic in the presence of non-determinism or parallelism

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- Proof about infinite derivation sequence requires invariant, which cannot be found automatically

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- Proof about infinite derivation sequence requires invariant, which cannot be found automatically

B. In all execution states, o is zero or one

$$\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \langle s', \sigma' \rangle \text{ and } \sigma(\mathtt{o}) = 0 \text{ then } \sigma'(\mathtt{o}) = 0 \text{ or } \sigma'(\mathtt{o}) = 1$$

- Prove needs to consider all possible derivations

Example 2: Verification of Parallel Programs

- A (simplified) Java program

```
class Cell {
  int x = 0;

  static void main(...) {
    Cell c = new Cell();
    Thread t1 = new Even(c);
    Thread t2 = new Even(c);
    t1.start(); t2.start();
    t1.join();  t2.join();
    System.out.println(c.x);
  }
}
```

```
class Even extends Thread {
  Cell c;

  Even(Cell c) {
    this.c = c;
  }

  void run() {
    c.x = c.x + 1;
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  }
}
```

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```

```
3
```

Example 2: Encoding in IMP

- The following program `s` represents the core of the Java program
  - `x` represents shared variable `c.x`
  - `y` and `z` represent thread-local state

```
  (y := x;  y := y + 1;  x := y;
   y := x;  y := y + 1;  x := y)
par
  (z := x;  z := z + 1;  x := z;
   z := x;  z := z + 1;  x := z)
```

- Desired property:
  If `x` is zero in the initial state then `x` is even in the terminal state
  - NS and Hoare logic cannot handle parallelism
  - SOS specification:

$$\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \sigma' \text{ and } \sigma(\mathtt{x}) = 0 \text{ then } \sigma'(\mathtt{x}) \ \mathbf{mod}\ 2 = 0$$

Example 2: Verification

- In this case, spotting the counterexample is easy, but how to attempt a formal proof?
- Induction does not work because there is no suitable induction hypothesis
  - Observation also holds for corrected example
- Proof strategy: enumerate all possible derivations of $\langle \mathtt{s}, \sigma \rangle \rightarrow_1^* \sigma'$ and inspect terminal state σ'
  - Number of derivations grows exponentially in number of executed statements
  - Here, $\frac{12!}{6! \times 6!} = 924$ possible derivations!
  - Manual enumeration not feasible, especially for programs with loops

Examples: Observations

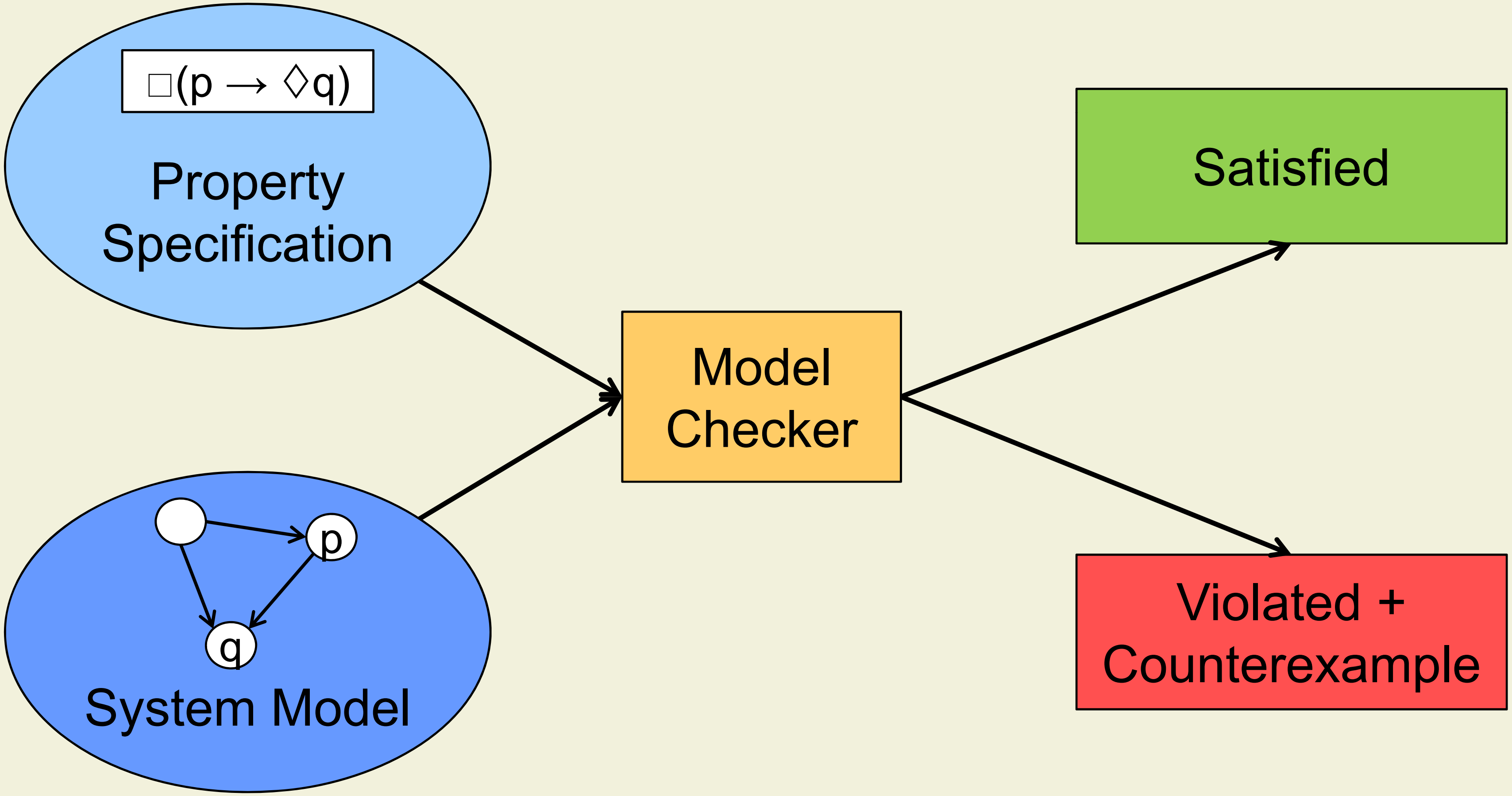
- Specification challenge
  - How to specify properties of sequences of states concisely
- Verification challenges
  - Concurrent systems: How to prove properties of all possible program executions
  - Reactive systems: How to prove automatically properties of infinite derivation sequences

Model Checking

> Model checking is an automated technique that, given a finite-state model of a system and a formal property, systematically checks whether this property holds for (a given state in) that model. [Baier and Katoen]

- Model checkers enumerate all possible states of a system:
  - Explicit state model checking: represent state explicitly through concrete values
  - Symbolic model checking: represent state through (boolean) formulas
- We focus on explicit state model checking

Model Checking



Model Checking Process

- Modelling phase
  - Model the system under consideration using the description language of your model checker (possibly a programming language)
  - Formalize the properties to be checked

- Running phase
  - Run the model checker to check the validity of the property in the system model

- Analysis phase
  - If property is satisfied, celebrate and move on to next property
  - If property is violated, analyze counterexample
  - If out of memory, reduce model and try again

Main Purposes of Model Checking

- Model checking is mainly used to analyze system designs (as opposed to implementations)
- Typical properties to be analyzed include
  - Deadlocks
  - Reachability of undesired states
  - Protocol violations

Modelling Concurrent Systems

- Systems are modelled as finite transition systems
- We model systems as communicating sequential processes (agents)
  - Finite number of processes
  - Interleaved process execution
- Processes can communicate via:
  - Shared variables
  - Synchronous message passing
  - Asynchronous message passing

Protocol Meta Language Promala

- Input language of the Spin model checker
- Main objects are processes, channels, and variables
- C-like syntax

```
init {
  printf("Hello World!\n")
}
```

- Spin can "execute" (simulate) models
- References
  - Quick reference: `www.spinroot.com/spin/Man/Quick.html`
  - Further references: `www.spinroot.com/spin/Man/index.html`

Promela Programs

- Constant declarations

```
#define N 5
mtype = { ack, req };
```

- Structure declarations

```
typedef vector { int x; int y };
```

- Global channel declarations

```
chan buf = [2] of { int };
```

- Global variable declarations

```
byte counter;
```

- Process declarations

```
proctype myProc(int p) { ... }
```

Promela Process Declarations

- Simple form

```
proctype myProc(int p) { ... }
```

  - Body consists of a sequence of variable declarations, channel declarations, and statements
  - No arrays as parameters

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- General form

```
active [N] proctype myProc(...) provided(E) priority M { ... }
```

  - `active`: Start `N` instances of `myProc` in the initial state
  - `provided`: `E` is an enabling condition, evaluated in the inital state
  - `priority`: `M` indicates probability during random simulation ($\mathtt{M} \geq 1$)

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- Deterministic processes lead to extra check during model analysis

```
D_proctype myProc(int p) { ... }
```

Promela Types

- Primitive types

Type	Value range
`bit` or `bool`	$0 \ldots 1$
`byte`	$0 \ldots 255$
`short`	$-2^{15} \ldots 2^{15}-1$
`int`	$-2^{31} \ldots 2^{31}-1$

  - No floats and mathematical integers

- Used-defined types
  - Arrays: `int name[4]`
  - Structures
  - Type of symbolic constants: `mtype`

- Channel type: `chan`

Promela Variable and Channel Declarations

- Variable declarations

```
byte a, b = 5, c;
int d[3], e[4] = 3;
mtype msg = ack;
vector v;
```

  - Variables are initialized to zero-equivalent values

Promela Variable and Channel Declarations

- Variable declarations

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byte a, b = 5, c;
int d[3], e[4] = 3;
mtype msg = ack;
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```

  - Variables are initialized to zero-equivalent values

- Channel declarations

```
chan c1 = [2] of { mtype, bit, chan };
chan c2 = [0] of { int };
chan c3;
```

  - c1 can store up to two messages
    Messages sent via c1 consist of three parts (triples)
  - c2 models rendez-vous communication
  - c3 is uninitialized; must be assigned an initialized channel before usage

Promela Variable and Channel Declarations

- Variable declarations

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byte a, b = 5, c;
int d[3], e[4] = 3;
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vector v;
```

  - Variables are initialized to zero-equivalent values

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    Messages sent via c1 consist of three parts (triples)
  - c2 models rendez-vous communication
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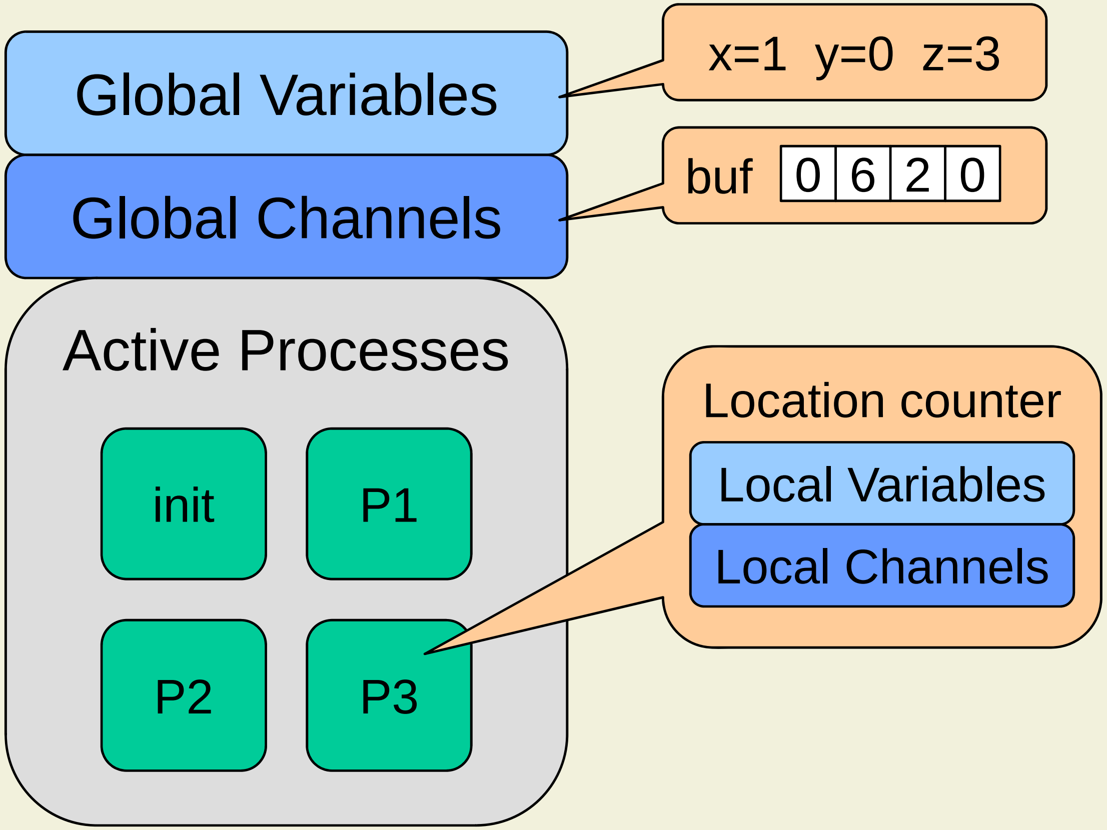
- Variable and channel declarations are local to a process or global

State Space of a Promela System



Global Variables

x=1 y=0 z=3

Global Channels

buf | 0 | 6 | 2 | 0 |

Active Processes

init

P1

P2

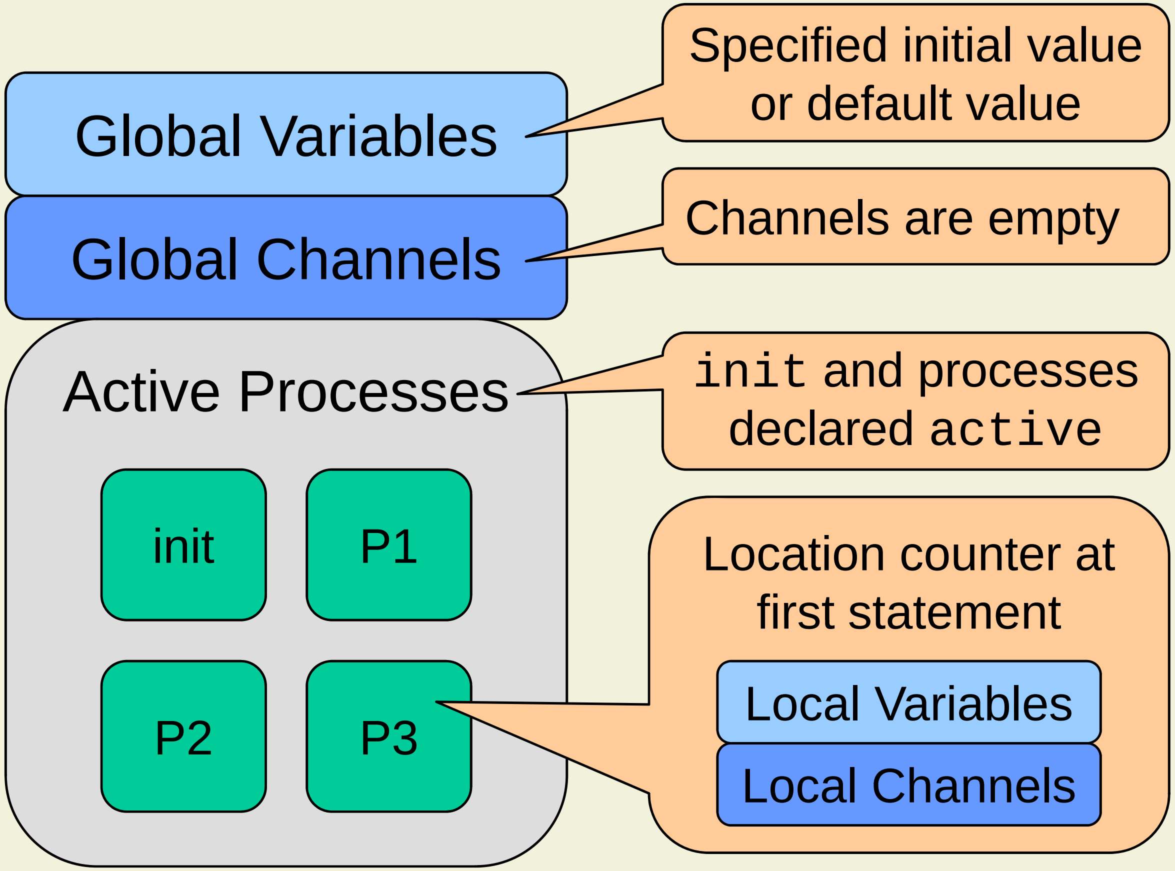
P3

Location counter

Local Variables

Local Channels

Initial State



Global Variables

Global Channels

Active Processes

init

P1

P2

P3

Specified initial value or default value

Channels are empty

`init` and processes declared `active`

Location counter at first statement

Local Variables

Local Channels

State Transitions

- A statement can be executable or blocked
  - Send is blocked if channel is full
  - `s1;s2` is blocked if `s1` is blocked
  - `timeout` is executable if all other statements are blocked

A transition is made in three steps:

- Determine all executable statements of all active processes
  - If no executable statement exists, transition system gets stuck
- Choose non-deterministically one of the executable statements
  - Non-determinism models concurrency through interleaving
- Change the state according to the chosen statement

Promela Expressions

- Variables, constants, and literals
- Structure and array accesses
- Unary and binary expressions with operators

```
+     -     *     /      %     >
>=    <     <=    ==     !=    !
&     ||    &&    |      ~     >>
<<    ^     ++    --
```

- Function applications

```
len()  empty()  nempty()   nfull()    full()
run    eval()   enabled()  pcvalue()
```

- Conditional expressions: `(E1 -> E2 : E3)`

Promela Statements

- `skip`
  - Does not change the state
  - Always executable
- `timeout`
  - Does not change the state
  - Executable if all other statements in the system are blocked
- `assert(E)`
  - Aborts execution if expression `E` evaluates to zero; otherwise equivalent to `skip`
  - Always executable
- Assignment
  - `x = E` assigns the value of `E` to variable `x`
  - `a[n] = E` assigns the value of `E` to array element `a[n]`
  - Always executable

Promela Statements (cont'd)

- Sequential composition
  - `s1;s2` is executable if `s1` is executable

- Expression statement
  - Evaluates expression `E`
  - Executable if `E` evaluates to value different from zero
  - `E` must not change state (no side effects)
  - Common case: function applications
  - Examples:

```
printf("Hello World!\n");
run myProcess;
x > 0;
```

Motivation: Verification of Parallel Programs

- A (simplified) Java program

```
class Cell {
  int x = 0;

  static void main(...) {
    Cell c = new Cell();
    Thread t1 = new Even(c);
    Thread t2 = new Even(c);
    t1.start(); t2.start();
    t1.join();  t2.join();
    System.out.println(c.x);
  }
}
```

```
class Even extends Thread {
  Cell c;

  Even(Cell c) {
    this.c = c;
  }

  void run() {
    c.x = c.x + 1;
    c.x = c.x + 1;
  }
}
```

```
3
```

Example: Modelling Even.run

```
class Even extends Thread {
  Cell c;

  Even(Cell c) {
    this.c = c;
  }

  void run() {
    c.x = c.x + 1;
    c.x = c.x + 1;
  }
}
```

```
int x;

proctype EvenRun() {
  x = x + 1;
  x = x + 1;
}
```

```
int x;

proctype EvenRun() {
  int y = x;
  y = y + 1;
  x = y;
  y = x;
  y = y + 1;
  x = y;
}
```

Example: Modelling Cell.main

```
class Cell {
  int x = 0;

  static void main(...) {
    Cell c = new Cell();
    Thread t1 = new Even(c);
    Thread t2 = new Even(c);
    t1.start(); t2.start();
    t1.join();  t2.join();
    System.out.println(c.x);
  }
}
```

```
init {
  x = 0;

  run EvenRun();
  run EvenRun();

  /* wait for termination */
  _nr_pr == 1;

  printf("x: %d\n", x);
  assert x % 2 == 0;
}
```

- `_nr_pr` is a predefined global variable that yields the number of active processes
- Simulation in Spin shows the possible outcomes 2, 3, and 4 (like Java program)

Promela Statements: Selection

```
if
:: s1   /* option 1 */
:: ...
:: sn   /* option n */
fi
```

- Executable if at least one of its options it executable
- Chooses an option non-deterministically and executes it

```
if  /* Move a sprite */
:: x < maxX -> x = x + 1;
:: x > minX -> x = x - 1;
:: y < maxY -> y = y + 1;
:: y > minY -> y = y - 1;
:: color = color + 1;
fi
```

- Statement `else` is executable if no other option is executable (may occur at most in one option)

Promela Statements: Repetition

```
do
:: s1   /* option 1 */
:: ...
:: sn   /* option n */
od
```

- Executable if at least one of its options it executable
- Chooses repeatedly an option non-deterministically and executes it
- Terminates when a `break` or `goto` is executed

```
/* compute factorial of n */

int r = 1;

do
:: n > 1 -> r = r*n; n = n-1;
:: else -> break
od
```

```
/* deadlock detection */
active proctype watchDog() {
  do
  :: timeout ->
     /* reset the state */
  od
}
```

Promela Statements: Atomic

- Basic statements are executed atomically
  - No interleaving during execution of statement
  - `skip`, `timeout`, `assert`, assignment, expression statement

- `atomic { s }` executes `s` atomically
  - If any statement within `s` blocks, atomicity is lost, and other processes are then allowed to execute statements

- Example: Binary semaphores (locks)

```
bit locked; /* global */
```

```
/* lock */
locked == 0;
locked = 1;


/* critical section */
locked = 0; /* unlock */
```

```
/* lock */
atomic {
  locked == 0;
  locked = 1;
}
/* critical section */
locked = 0; /* unlock */
```

Promela Macros

- Promela does not contain procedures
- Effect can often be achieved using macros

```
inline lock() {
  atomic {
    locked == 0;
    locked = 1
  }
}
```

```
inline swap(a, b) {
  int tmp;
  tmp = a;
  a = b;
  b = tmp
}
```

- A macro just defines a replacement text for a symbolic name, possibly with parameters
  - The inline call `lock()` is replaced by the body of the definition
  - No new variable scope
  - No recursion
  - No return values
- Define macro globally before its first use

Motivation: Deadlock

- Threads are synchronized via locks
- Interleaved execution of `a.transfer(b,n)` and `b.transfer(a,m)` might deadlock
- Multi-threaded programs are extremely hard to test

```
class Account {
  int balance;

  void transfer(Account to, int amount) {
    acquire this;
    acquire to;
    this.balance -= amount;
    to.balance += amount;
    release this;
    release to;
  }
}
```

Promela Model: Account

- We need to model accounts and clients
  - General approach: omit all irrelevant details to reduce complexity
- Account
  - Balance is not relevant for potential deadlocks
  - Only model the locks of accounts

```
#define N 5

bit Account_locks[N];

inline lock(n) {
  atomic {
    Account_locks[n] == 0;
    Account_locks[n] = 1;
  }
}
```

Promela Model: Client

- Idea: model the most generic client and run several instances in parallel
    - Pick two arbitrary accounts non-deterministically
    - Lock both accounts
    - Unlock both accounts

- Choosing accounts

```
inline choose(a, l, u) {
  a = l;
  do
  :: (a < u) -> a++
  :: break
  od
}
```

```
inline chooseAccounts(f, t) {
  do
  :: (f != t) -> break
  :: (f == t) -> choose(f, 0, N-1);
                 choose(t, 0, N-1)
  od
}
```

Promela Model: Client Process

```
active [C] proctype transfer() {
  byte from, to;

  /* choose accounts non-deterministically */
  chooseAccounts(from, to);

  /* acquire locks */
  lock(from);
  lock(to);

  /* actual transfer omitted */

  /* release locks */
  Account_locks[from] = 0;
  Account_locks[to] = 0;
}
```

Alternative Account Selection

- Idea: instead of looping until two different accounts are found, restrict range for second choice

```
do
:: (f != t) -> break
:: (f == t) -> choose(f, 0, N-1);
               choose(t, 0, N-1)
od
```

```
choose(f, 0,   N-2);
choose(t, f+1, N-1)
```

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:: (f == t) -> choose(f, 0, N-1);
               choose(t, 0, N-1)
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```

```
choose(f, 0,   N-2);
choose(t, f+1, N-1)
```

- Alternative model is less general
  - It guarantees `from` $<$ `to`
  - So locks are acquired in order and deadlock is prevented!

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```

- Alternative model is less general
  - It guarantees `from` < `to`
  - So locks are acquired in order and deadlock is prevented!
- General strategy
  - Start with most general model
  - If model contains errors that cannot occur in real system (spurious error), revise model

Promela Channels

- `chan ch = [d] of { t1, ..., tn }` declares a channel
- Channel can buffer up to `d` messages
  - `d` > 0: buffered channel (FIFO)
  - `d` $= 0$: unbuffered channel (rendez-vous)
- Each message is a tuple whose elements have types `t1, ..., tn`
- Example

```
mtype = { req, ack, err };

chan ch = [5] of { mtype, int }
```

Send and Receive: Buffered Channels

```
chan ch = [5] of { mtype, int }
```

- `ch ! e1, ..., en` sends message
  - Type of `ei` must correspond to `ti` in channel declaration
  - Send is executable iff buffer is not full

- `ch ? a1, ..., an` receives message
  - `ai` is a variable or constant of type `ti`
  - Receive is executable iff buffer is not empty and the oldest message in the buffer matches the constants `ai`
  - Variables `ai` are assigned values of the message

```
ch ! req, 7;
ch ! ack, 1
```

```
int n;
ch ? req, n;
printf("Received: %d\n", n);
ch ? req, n;
printf("Received: %d\n", n);
```

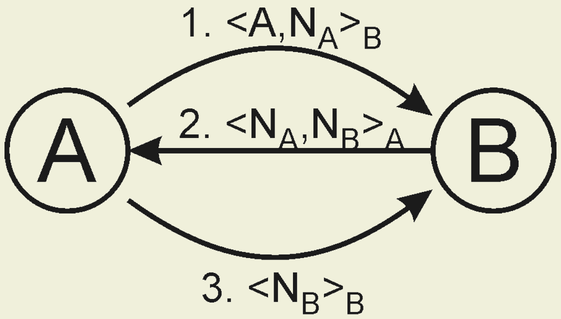
Send and Receive: Unbuffered Channels

```
chan ch = [0] of { int };
```

- `ch ! e1, ..., en` sends message
  - Send is executable if there is a receive operation that can be executed simultaneously
- `ch ? a1, ..., an` receives message
  - Receive is executable if there is a send operation that can be executed simultaneously
- Unbuffered channels model synchronous communication (rendez-vous)

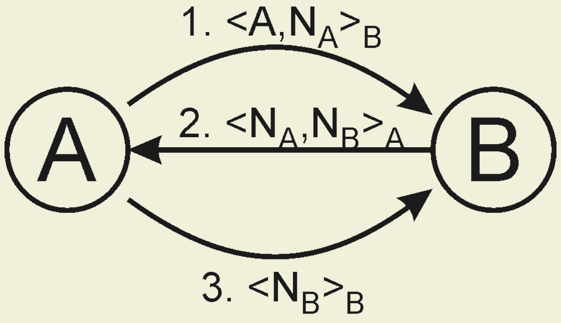
Motivation: Needham-Schroeder Protocol

- Establish a common secret over an insecure channel
  1. Alice sends random number N_A to Bob, encrypted with Bob's public key: $\langle A, N_A \rangle_B$
  2. Bob sends random number N_B to Alive, encrypted with Alice's public key: $\langle N_A, N_B \rangle_A$
  3. Alice responds with $\langle N_B \rangle_B$



- Intruders may:
  - Intercept, store, and replay messages
  - Initiate or participate in runs of the protocol
  - Decrypt messages only if encrypted with intruder's public key
- Error: intruder can pretend to be another party

Promela Model: Network



- We model the protocol for two agents plus intruder
- Agents communicate synchronously

```
chan network = [0] of {
  mtype, /* tag:                msg1, msg2, msg3       */
  mtype, /* intended receiver: agentA, agentB, agentI */
  Crypt  /* message                                    */
};
```

- We use enumeration type `mtype` for all constants
    - Spin treats `mtype` constants as symbols, not values
    - Speeds up model checking

Promela Model: Messages

- Message consists of key and up to two contents

```
typedef Crypt {
  mtype key,          /* public key used to encrypt */
        content1,     /* agent or nonce             */
        content2      /* nonce or don't care        */
};
```

- We model encryption by putting the public key into the message
  - Agent *a* will only look at message content if message key is *a*'s public key
  - No need to model private keys and encryption

- Constants for message tag, public keys, agents, and nonces

```
mtype = { msg1, msg2, msg3,
          keyA, keyB, keyI,
          agentA, agentB, agentI,
          nonceA, nonceB, nonceI };
```

Promela Model: Alice

- Protocol runs are always started by Alice

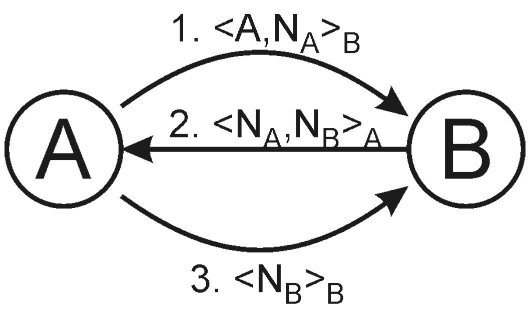
```
mtype partnerA;
bit   statusA; /* 1 = success */

active proctype Alice() {
  mtype pkey;      /* the partner's public key          */
  mtype pnonce;    /* nonce that we receive from partner */
  Crypt message;   /* our message to the partner         */
  Crypt data;      /* received message                   */

  if /* choose a partner for this run */
  :: partnerA = agentB; pkey = keyB;
  :: partnerA = agentI; pkey = keyI;
  fi;

  /* Protocol run below */
  statusA = 1; /* Success */
}
```

Promela Model: Alice's Protocol Run



```
/* Prepare and send first message */
build(message, pkey, agentA, nonceA);
network ! msg1, partnerA, message;

/* Wait for answer */
network ? (msg2, agentA, data);

/* Proceed only if the key matches keyA and the
   nonce is the one that we have sent earlier */
(data.key == keyA) && (data.content1 == nonceA);

/* Obtain partner's nonce */
pnonce = data.content2;

/* Prepare and send the last message */
build(message, pkey, pnonce, 0);
network ! msg3, partnerA, message;
```

Intruder

- Intruders may:
  - Intercept messages
  - Store one message
  - Replay messages
  - Initiate or participate in runs of the protocol
  - Decrypt messages only if encrypted with intruder's public key

- How can we model the most powerful attack using these capabilities?

- Solution: Model intruder fully non-deterministically
  - Intruder has no intellegence whatsoever
  - Model checker will explore all possible behaviors of intruder

Promela Model: Intruder

```
bool knows_nonceA, knows_nonceB;

active proctype Intruder() {
  mtype tag;          /* message tag                */
  mtype recpt;        /* recipient for our message */
  Crypt data          /* received message           */
  Crypt intercepted; /* stored message             */

  do
  :: /* Receive and learn */

  :: /* Replay or send    */
  od
}
```

Promela Model: Intruder Receives

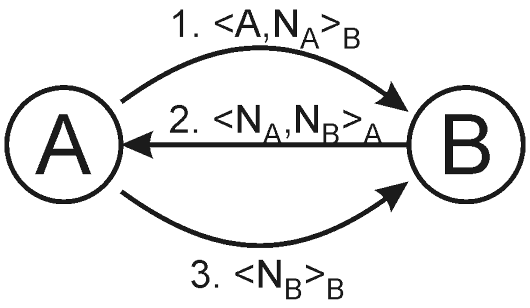
```
do
:: network ? (tag, _, data) ->
   if /* perhaps store the message */
   :: copy(data, intercepted);
   :: skip;
   fi;
   if /* record newly learnt nonces */
   :: (data.key == keyI) ->
      knows_nonceA = knows_nonceA ||
                     (data.content1 == nonceA) ||
                     (data.content2 == nonceA);
      knows_nonceB = knows_nonceB ||
                     (data.content1 == nonceB) ||
                     (data.content2 == nonceB);
   :: else -> skip;
   fi;
:: /* Replay or send    */
od
```

Promela Model: Intruder Sends

```
do
:: /* Receive and learn */
:: /* Replay or send    */
   if /* choose message type */
   :: tag = msg1;
   :: tag = msg2;
   :: tag = msg3;
   fi;
   if /* choose recipient */
   :: recpt = agentA;
   :: recpt = agentB;
   fi;
   if /* replay intercepted message or assemble it */
   :: copy(intercepted, data);
   :: /* assemble new message */
   fi;
   network ! tag, recpt, data;
od
```

Promela Model: Intruder Sends (cont'd)

```
    :: /* assemble new message */
       if
       :: data.key = keyA;
       :: data.key = keyB;
       fi;
       if
       ::                 data.content1 = agentA;
       ::                 data.content1 = agentB;
       ::                 data.content1 = agentI;
       :: knows_nonceA -> data.content1 = nonceA;
       :: knows_nonceB -> data.content1 = nonceB;
       ::                 data.content1 = nonceI;
       fi;
       if
       :: knows_nonceA -> data.content2 = nonceA;
       :: knows_nonceB -> data.content2 = nonceB;
       ::                 data.content2 = nonceI;
       fi;
```



Summary

- Models are abstractions of the real world
- Omit irrelevant details to reduce complexity
  - Example: balance in account example
- Keep model small to avoid state space explosion
  - As few process as possible
  - As little data as possible
- Non-determinism is a powerful modelling tool
  - Let model checker explore all options
- Typical sources of non-determinism are:
  - Abstraction
  - Modelling of the environment