

Assignment 4

Exercise 1

Consider the following Alloy model of a counter:

```
module util/integer

sig Counter {
  n: Int
}

pred inc[c, c': Counter] {
  c'.n = c.n.add[Int[1]]
}
```

The predicate `inc` models the increment operation of the counter.

1. Use the abstract machine pattern that you have seen in the lecture to define the traces of the counter; see slides 148-149. For this task you must use the Alloy library `util/ordering`. Initially, the value of n must be 0. Define the initial state of the counter using the predicate `init[c: Counter]`.
2. Encode the following two invariants using assertions:
inv1: For any counter, $n \geq 0$.
inv2: The value of a counter never decreases.

Note that **inv1** is a state invariant. You must check whether it holds for all states. The second invariant **inv2** is a temporal invariant. You must check whether for any two states c and c' where c is a predecessor of c' , $c.n \leq c'.n$ holds. You can use the predicate `lt[c, c']` to check whether c is a predecessor of c' .

Exercise 2

Recall the two `ImageFile` class implementations from the lecture. You are given the Alloy model for the `ImageFile` implementation with eager initialization; see the file `imagefile_eager.als`. Consider the following two invariants:

inv1: `getImage()` returns a non-null value.

inv2: `getImage()` returns the same value for an entire trace.

You have two tasks:

1. Define the initial state and the traces for the `ImageFile` model. Model the two invariants as assertions and check whether they hold.
2. Create an Alloy model for the alternative implementation of `ImageFile`, which uses lazy initialization:

```
class ImageFile {
    String file;
    Image image;
    ImageFile(String f) {
        file = f;
    }
    Image getImage() {
        if(image == null) {
            // load the image
        }
        return image;
    }
}
```

Observe several traces to ensure that your model has the intended behavior. Check whether the invariants that you defined in (1) hold for your model.

Hint: You must model `getImage()` as a predicate as it may change the state. Note that we modeled `getImage()` in the eager version as a function as it does not change the state.

Exercise 3

Consider the following Alloy model:

```

sig Node {
  next: Node
}
assert demo {
  all n: Node | some m: Node | m.next = n
}

```

An assertion a is *valid* iff it evaluates to true in all instances that satisfy the models' constraints c . Note that for the given Alloy model, the constraint c is satisfied by instances where each node has exactly one next node. Recall that to check an assertion, Alloy checks whether there exists a counter-example; namely, an instance where the constraints are satisfied and the assertion is violated ($c \wedge \neg a$).

1. Encode the constraint and the assertion that correspond to checking `check demo for 1` into a boolean formula.

Hint: Treat the universal and existential quantifiers as conjunction and disjunction, respectively.

Check if the boolean formula has a satisfying assignment. If the boolean formula is satisfiable, then give a counter-example where the assertion is violated.

2. Repeat step (1) for `check demo for 2`.

3. Add a field `prev: Node` to the signature of `Node`.

Add the fact `all n: Node | n.next.prev = n` to the Alloy model.

Repeat steps (1) and (2).

Does there exist a scope, possibly larger than 2, where the assertion is violated?