

Challenge II: DLL to BST

Description

This challenge is to verify an in-place algorithm to convert a sorted doubly-linked list (DLL) into a balanced binary search tree (BST). The algorithm runs in linear time, and proceeds by first making a pass over the list to calculate its length, say n . It then recursively constructs a BST out of the first $n/2$ nodes, makes this the left subtree of the middle node, and then recursively constructs the right subtree using the remaining nodes.

Assume that each node has 3 fields, `data`, `prev`, and `next`, and when constructing the BST, we use the `prev` field to link a node to its left child and the `next` field to link a node to its right child.

The code for this algorithm is given in Figure 1.

Verification Tasks

1. Prove that this algorithm converts an input list into a tree.
2. Prove that the algorithm is memory-safe.
3. Prove that if the input list is sorted then the resulting tree is a BST.
4. Prove that the resulting BST is balanced.
5. Prove that the algorithm terminates.
6. (Optional) Prove the above for an iterative version of size.

```

1 // Ref is the type of nodes used for both list and tree, and has these fields:
2 field data: Int
3 field prev: Ref // Also used as left subtree pointer
4 field next: Ref // Also used as right subtree pointer
5
6 method size(head: Ref) returns (count: Int) {
7   if (head != null) {
8     count := size(head.next)
9     count := count + 1
10  } else {
11    count := 0
12  }
13 }
14
15 method dll_to_bst(head: Ref) returns (root: Ref) {
16   var n: Int
17   var right: Ref
18   n := size(head)
19   root, right := dll_to_bst_rec(head, n)
20 }
21
22 // Converts a sorted DLL into a balanced BST
23 // head: Pointer to doubly linked list
24 // n: number of nodes of list to convert to tree
25 method dll_to_bst_rec(head: Ref, n: Int) returns (root: Ref, right: Ref) {
26   if (n > 0) {
27     // Recursively construct the left subtree
28     var left: Ref
29     left, root := dll_to_bst_rec(head, n/2)
30     // [head, root) is a tree rooted at left, [root, ...] is a list
31
32     // Set pointer to left subtree
33     root.prev := left
34
35     // Recursively construct the right subtree
36     // size(right subtree) = n - size(left subtree) - 1 (for root)
37     var temp: Ref
38     temp, right := dll_to_bst_rec(root.next, n-n/2-1)
39     // [head, root) is a tree rooted at left, [root.next, right) is tree at temp
40
41     // Set pointer to right subtree
42     root.next := temp
43     // [head, right) is a tree rooted at root
44   } else {
45     root := null
46     right := head
47   }
48 }

```

Figure 1: Code for the algorithm that converts a sorted DLL into a balanced BST.