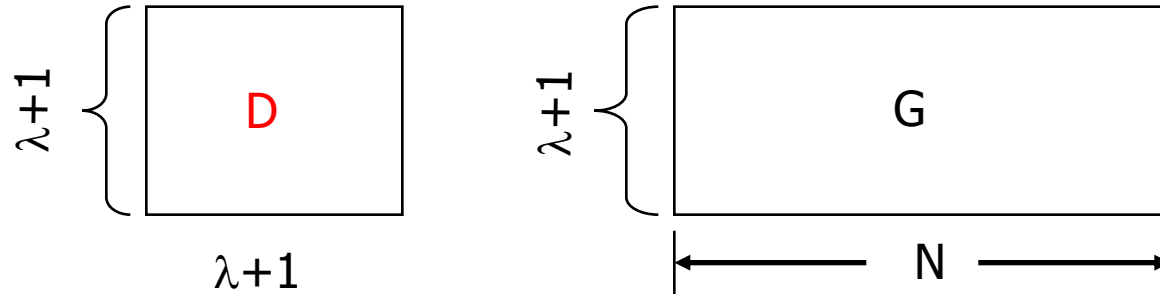


Deterministic Approaches

- Used to design the key pool and the key chains to provide better connectivity
 - Matrix Based Scheme [Blom 1985]
 - Polynomial Based Key Generation [Blundo et al. 1992]

Deterministic approaches: Blom's Scheme [B]

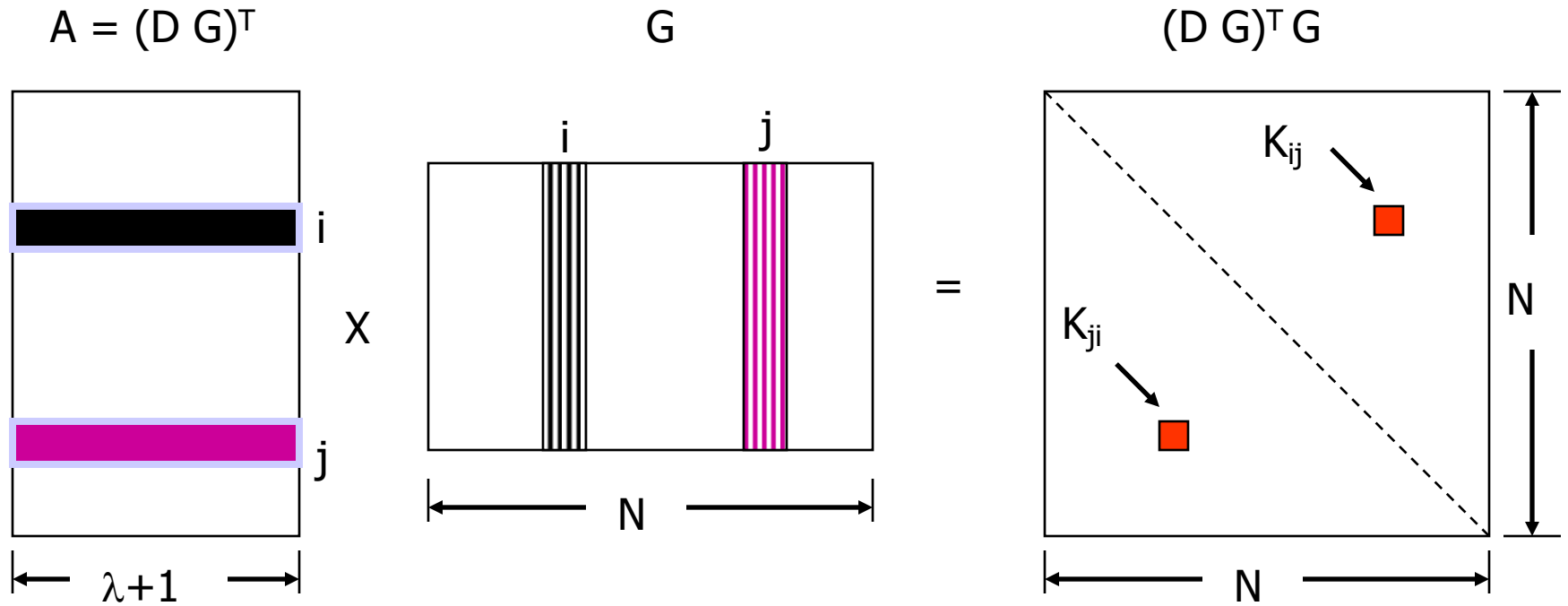
- Public matrix G
- Private matrix D (symmetric).



$$\text{Let } \mathbf{A} = (\mathbf{D} \ \mathbf{G})^T$$

$$\mathbf{A} \ \mathbf{G} = (\mathbf{D} \ \mathbf{G})^T \ \mathbf{G} = \mathbf{G}^T \ \mathbf{D}^T \ \mathbf{G} = \mathbf{G}^T \ \mathbf{D} \ \mathbf{G} = (\mathbf{A} \ \mathbf{G})^T$$

[B] Scheme



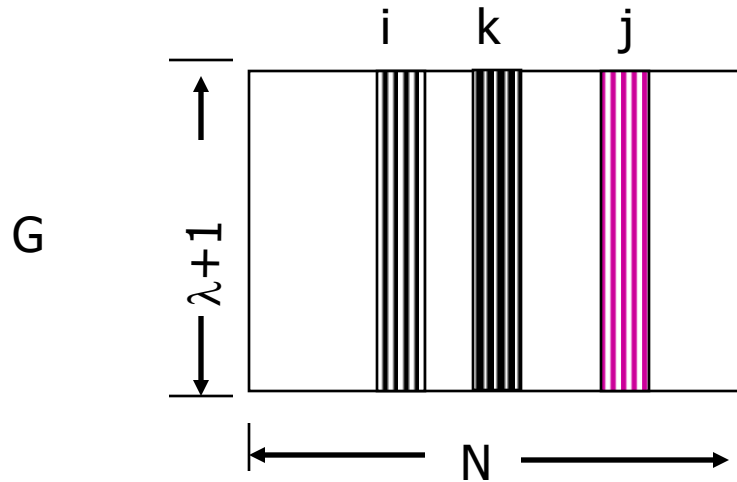
Node i carries:



Node j carries:

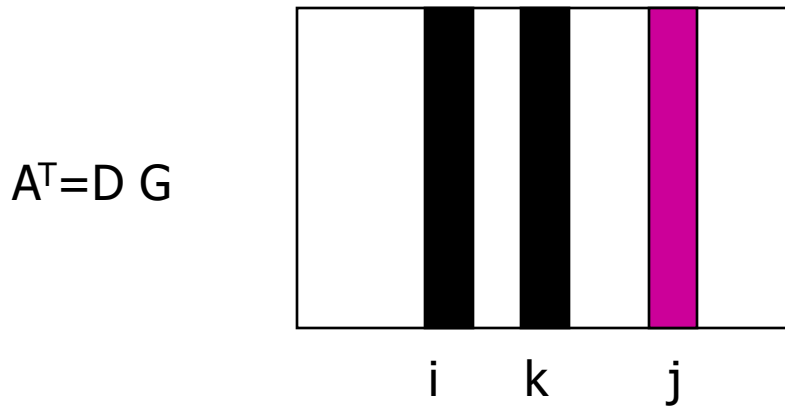


[B] λ -secure Property



Undesirable Situation:
if
 $u \cdot G(i) + v \cdot G(j) = G(k)$

then
 $u \cdot A(i) + v \cdot A(j) = A(k)$



this would allow
colluding nodes (i and j) to
impersonate other nodes (k)

[B] λ -secure Property

- **ALL** $\lambda+1$ columns in G are linear independent.
 - Different from saying that G has rank $\lambda+1$
 - **Rank:** there are $\lambda+1$ linearly independent columns
- Can tolerate compromise up to λ nodes.
 - Once $\lambda+1$ nodes are compromised, the rest can be calculated if these $\lambda+1$ columns are linear independent.
- How to find such a matrix G ?

[B] Vandermonde Matrix

$$\mathbf{G} = \begin{pmatrix} 1 & 1 & 1 & \text{-----} & 1 \\ s & s^2 & s^3 & \text{-----} & s^N \\ s^2 & (s^2)^2 & (s^3)^2 & \text{-----} & (s^N)^2 \\ & & \vdots & & \\ s^\lambda & (s^2)^\lambda & (s^3)^\lambda & \text{-----} & (s^N)^\lambda \end{pmatrix}$$

[B] Properties of Blom Scheme

- Blom's Scheme
 - Network size is N
 - Any pair of nodes can **directly** find a secret key
 - Tolerate compromise up to λ nodes
 - Need to store $\lambda+2$ keys

Key distribution schemes for sensor networks

<http://www.cs.rpi.edu/research/pdf/05-07.pdf>

Problem	Approach	Mechanism	Keying style	Papers
Pair-wise	Probabilistic	Pre-distribution	Random key-chain	C, E, F, J K, N, S
			Pair-wise key	E
	Deterministic	Pre-distribution	Pair-wise key	G, M
			Combinatorial	P, Q
			Dynamic Key Generation	Master key
		Key matrix	A	
		Polynomial	B, G	
	Hybrid	Pre-distribution	Combinatorial	P, Q
		Dynamic Key Generation	Key matrix	H, M, R
		Polynomial	I, R	
Group-wise	Deterministic	Dyn. Key Gen.	Polynomial	B, R

The papers are: A[Blom 1985], B[Blundo et al. 1992], C[Eschenauer and Gligor 2002], D[Lai et al. 2002], E[Chan et al. 2003], F[Pietro et al. 2003], G[Liu and Ning 2003c], H[Du et al. 2003], I[Liu and Ning 2003b], J[Zhu et al. 2003], K[Du et al. 2004], L[Dutertre et al. 2004], M[Lee and Stinson 2004b], N[Hwang et al. 2004], P[Camtepe and Yener 2004], Q[Lee and Stinson 2004a], R[Huang et al. 2004], S[Hwang and Kim 2004].