EVENT-DRIVEN AND DATA-DRIVEN CONTROL AND OPTIMIZATION IN CYBER-PHYSICAL SYSTEMS

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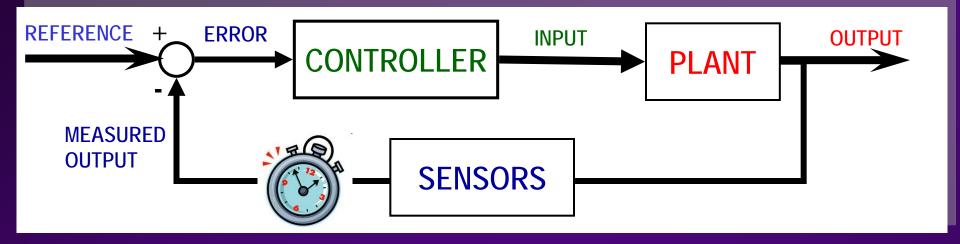
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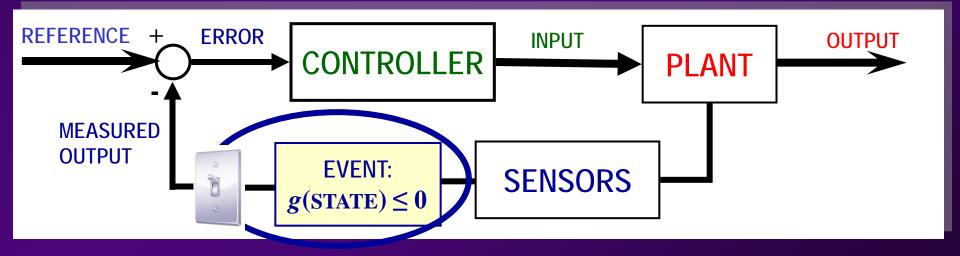
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TIME-DRIVEN v EVENT-DRIVEN CONTROL

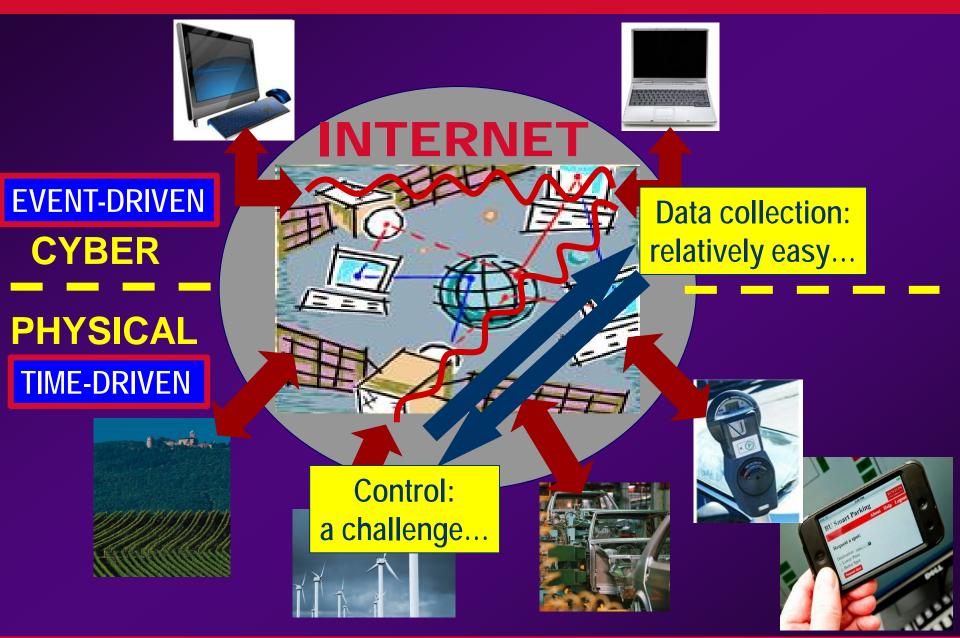


EVENT-DRIVEN CONTROL: Act only when needed (or on TIMEOUT) - not based on a clock



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CYBER-PHYSICAL SYSTEMS



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OUTLINE

Why EVENT-DRIVEN Control and Optimization ?

EVENT-DRIVEN Control in Distributed Multi-Agent Systems

A General Optimization Framework for Multi-Agent Systems

EVENT-DRIVEN + DATA-DRIVEN Control and Optimization: the IPA (Infinitesimal Perturbation Analysis) Calculus

REASONS FOR EVENT-DRIVEN MODELS, CONTROL, OPTIMIZATION

- Many systems are naturally Discrete Event Systems (DES) (e.g., Internet)
 - \rightarrow all state transitions are event-driven
- Most of the rest are Hybrid Systems (HS) \rightarrow some state transitions are event-driven
- Many systems are distributed

 → components interact asynchronously (through events)
- Time-driven sampling inherently inefficient ("open loop" sampling)

REASONS FOR *EVENT-DRIVEN* MODELS, CONTROL, OPTIMIZATION

Many systems are stochastic

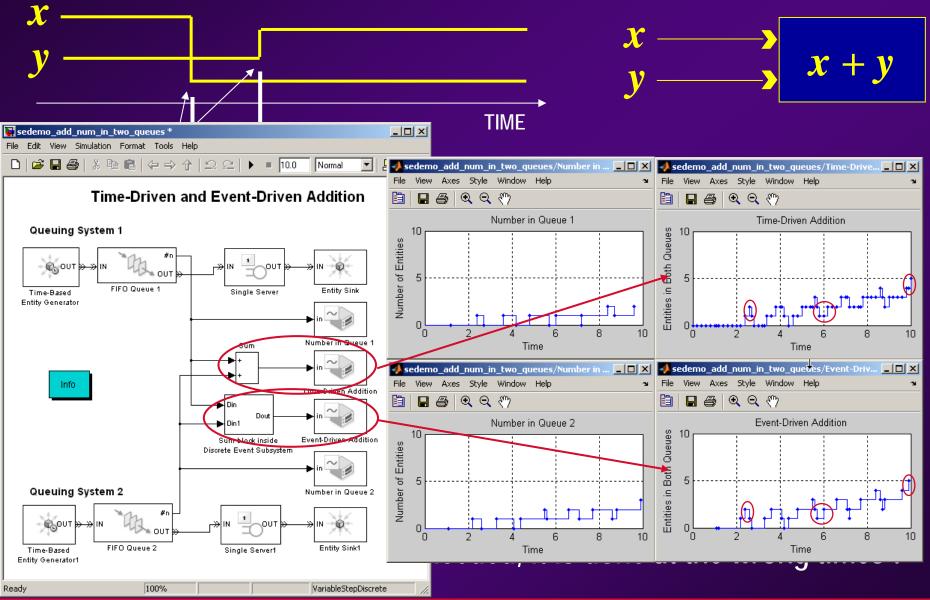
 \rightarrow actions needed in response to random events

Event-driven methods provide significant advantages in computation and estimation quality

 System performance is often more sensitive to event-driven components than to time-driven components

 Many systems are wirelessly networked → energy constrained
 → time-driven communication consumes significant energy UNNECESSARILY!

TIME-DRIVEN (SYNCHRONOUS) v EVENT-DRIVEN (ASYNCHRONOUS) COMPUTATION



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SELECTED REFERENCES - EVENT-DRIVEN CONTROL, COMMUNICATION, ESTIMATION, OPTIMIZATION

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P. Wan and M. D. Lemmon, "Event triggered distributed optimization in sensor networks," *Proc. of 8th ACM/IEEE Intl. Conf. on Information Processing in Sensor Networks*, 2009.
Zhong, M., and Cassandras, C.G., "Asynchronous Distributed Optimization with Event-Driven Communication", *IEEE Trans. on Automatic Control*, AC-55, 12, pp. 2735-2750, 2010.

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EVENT-DRIVEN DISTRIBUTED OPTIMIZATION

DISTRIBUTED COOPERATIVE OPTIMIZATION

N system components (processors, agents, vehicles, nodes), one common objective:

$$\min_{s_1,\ldots,s_N} H(s_1,\ldots,s_N)$$

s.t. constraints on each s_i

$$\min_{s_1} H(s_1,\ldots,s_N)$$

s.t. constraints on
$$s_1$$

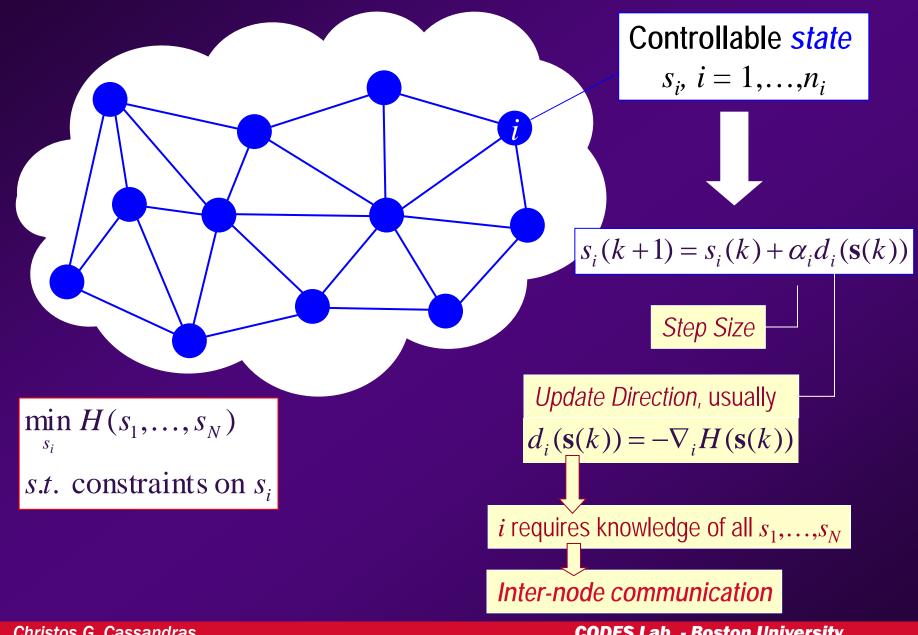


$$\min_{s_N} H(s_1, \dots, s_N)$$

s.t. constraints on s_N

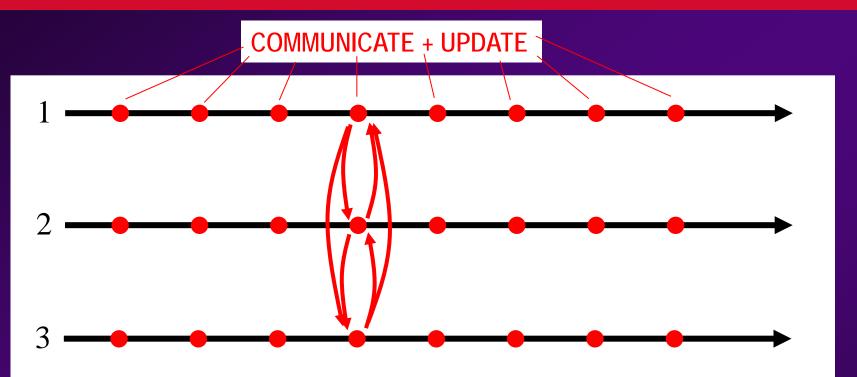
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DISTRIBUTED COOPERATIVE OPTIMIZATION



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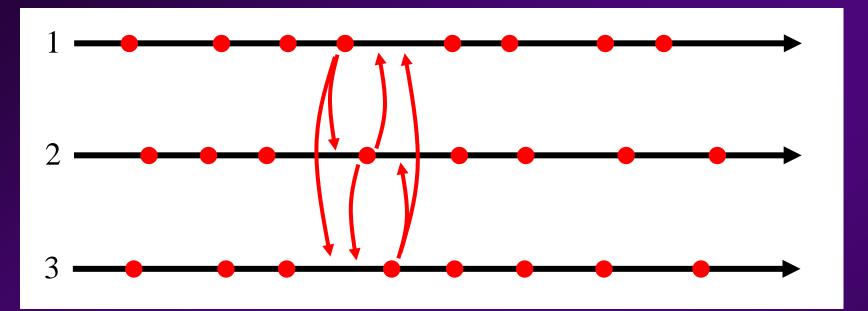
SYNCHRONIZED (TIME-DRIVEN) COOPERATION



Drawbacks:

- Excessive communication (critical in wireless settings!)
- Faster nodes have to wait for slower ones
- Clock synchronization infeasible
- Bandwidth limitations
- Security risks

ASYNCHRONOUS COOPERATION



Nodes not synchronized, delayed information used

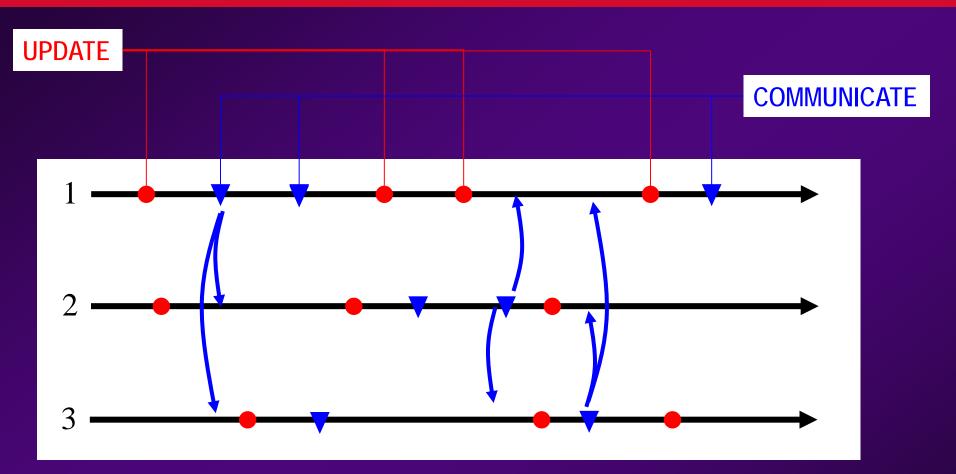
Update frequency for each node is bounded

technical conditions

 $\Rightarrow \frac{s_i(k+1) = s_i(k) + \alpha_i d_i(\mathbf{s}(k))}{\text{converges}}$

Bertsekas and Tsitsiklis, 1997

ASYNCHRONOUS (EVENT-DRIVEN) COOPERATION



UPDATE at *i* : locally determined, arbitrary (possibly periodic)
 COMMUNICATE from *i* : only when absolutely necessary

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WHEN SHOULD A NODE COMMUNICATE?

AT ANY TIME *t* :

- $x_i^j(t)$: node *i* state estimated by node *j*
- If node *i* knows how *j* estimates its state, then it can evaluate $x_i^j(t)$
- Node *i* uses
 - its own true state, $x_i(t)$
 - the estimate that j uses, $x_i^j(t)$

... and evaluates an ERROR FUNCTION $g(x_i(t), x_i^j(t))$

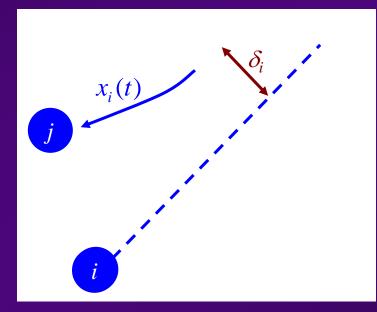
Error Function examples:
$$\left\|x_{i}(t) - x_{i}^{j}(t)\right\|_{1}, \quad \left\|x_{i}(t) - x_{i}^{j}(t)\right\|_{2}$$

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WHEN SHOULD A NODE COMMUNICATE?

Compare ERROR FUNCTION $g(x_i(t), x_i^j(t))$ to THRESHOLD δ_i

Node *i* communicates its state to node *j* only when it detects that its *true state* $x_i(t)$ deviates from *j*' *estimate of it* $x_i^j(t)$ so that $g(x_i(t), x_i^j(t)) \ge \delta_i$



⇒ *Event-Driven* Control

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CONVERGENCE

Asynchronous distributed state update process at each *i*:

$$s_i(k+1) = s_i(k) + \alpha \cdot d_i(\mathbf{s}^i(k))$$

Estimates of other nodes, evaluated by node i

$$\delta_i(k) = \begin{cases} K_{\delta} \left\| d_i(\mathbf{s}^i(k)) \right\| & \text{if } k \text{ sends update} \\ \delta_i(k-1) & \text{otherwise} \end{cases}$$

THEOREM: Under certain conditions, there exist positive constants α and K_{δ} such that

 $\lim_{k\to\infty}\nabla H(\mathbf{s}(k))=0$

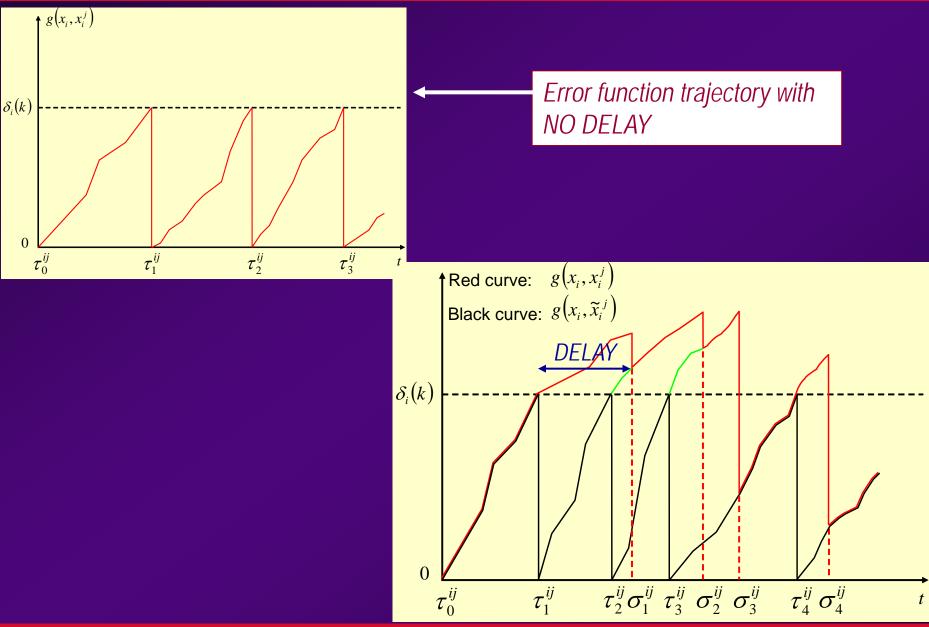
Zhong and Cassandras, IEEE TAC, 2010

INTERPRETATION:

Event-driven cooperation achievable with minimal communication requirements \Rightarrow *energy savings*

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COONVERGENCE WHEN DELAYS ARE PRESENT



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COONVERGENCE WHEN DELAYS ARE PRESENT

Add a boundedness assumption:

ASSUMPTION: There exists a non-negative integer *D* such that if a message is sent before t_{k-D} from node *i* to node *j*, it will be received before t_k .

INTERPRETATION: at most **D** state update events can occur between a node sending a message and all destination nodes receiving this message.

THEOREM: Under certain conditions, there exist positive constants α and K_{δ} such that

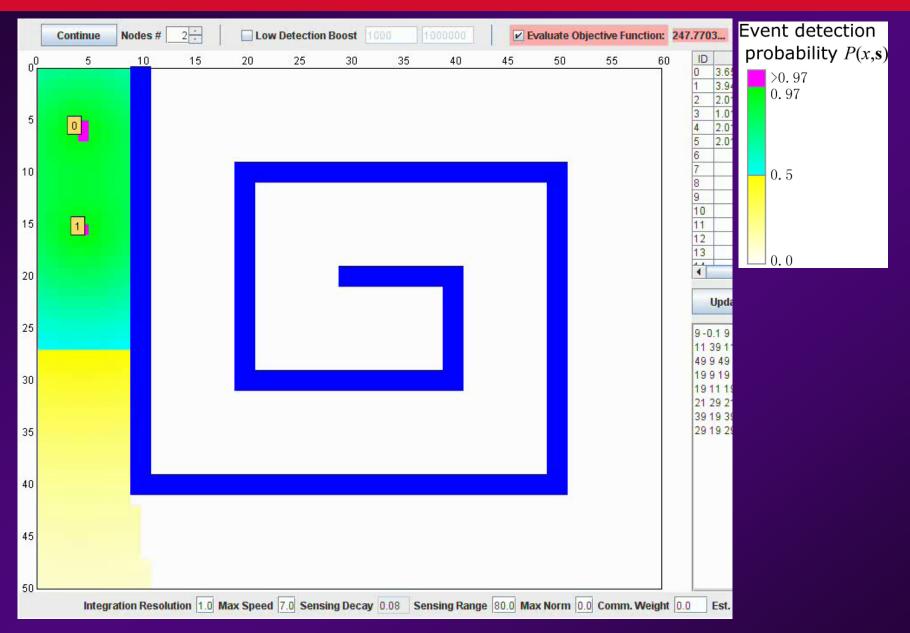
 $\lim_{k\to\infty}\nabla H(\mathbf{s}(k))=0$

NOTE: The requirements on α and K_{δ} depend on **D** and they are tighter.

Zhong and Cassandras, IEEE TAC, 2010

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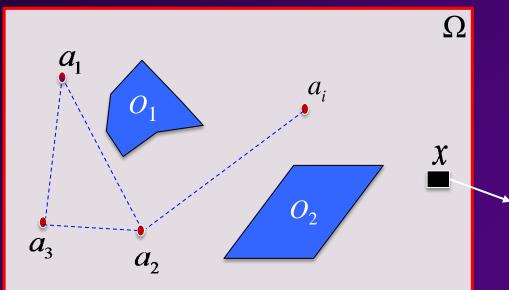
OPTIMAL COVERAGE IN A MAZE



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A GENERAL OPTIMIZATION FRAMEWORK FOR MULTI-AGENT SYSTEMS

NETWORKED MULTI-AGENT OPTIMIZATION: PROBLEM 1: PARAMETRIC OPTIMIZATION



- *s_i*: agent state, *i* = 1,..., *N s*=[*s*₁, ..., *s*_N]
- *O_j*: obstacle (constraint)
- R(x): property of point x
- P(x, s): reward function

$$\max_{\mathbf{s}} H(\mathbf{s}) = \int_{\Omega} P(x, \mathbf{s}) R(x) dx$$
$$s_i \in F \subset \Omega \quad i = 1 \cdots N$$

GOAL: Find the best state vector $s = [s_1, ..., s_N]$ so that agents achieve a maximal reward from interacting with the mission space

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NETWORKED MULTI-AGENT OPTIMIZATION: PROBLEM 2: DYNAMIC OPTIMIZATION

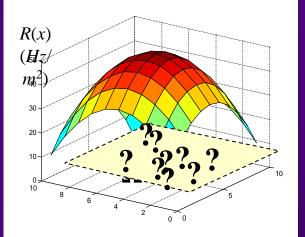
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PROBLEMS THAT FIT THIS FRAMEWORK

COVERAGE

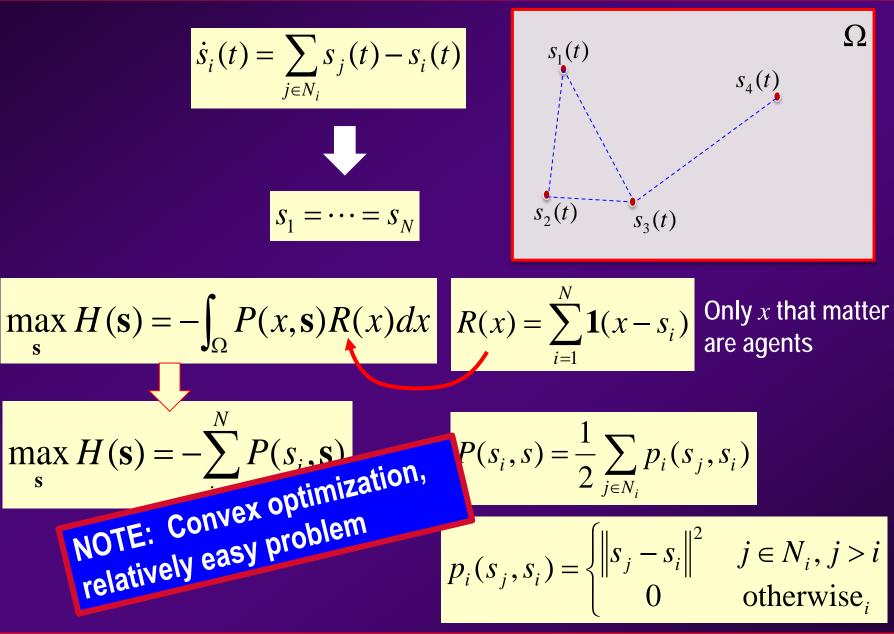
Deploy sensors to maximize "event" detection probability - unknown event locations



$$\max_{\mathbf{s}} H(\mathbf{s}) = \int_{\Omega} P(x, \mathbf{s}) R(x) dx$$

Joint event detection probability: $P(x, \mathbf{s}) = 1 - \prod_{i=1}^{N} \left[1 - p_i(x, s_i) \right]$ Event sensing probability *Event density*: Prior estimate of event occurrence frequency

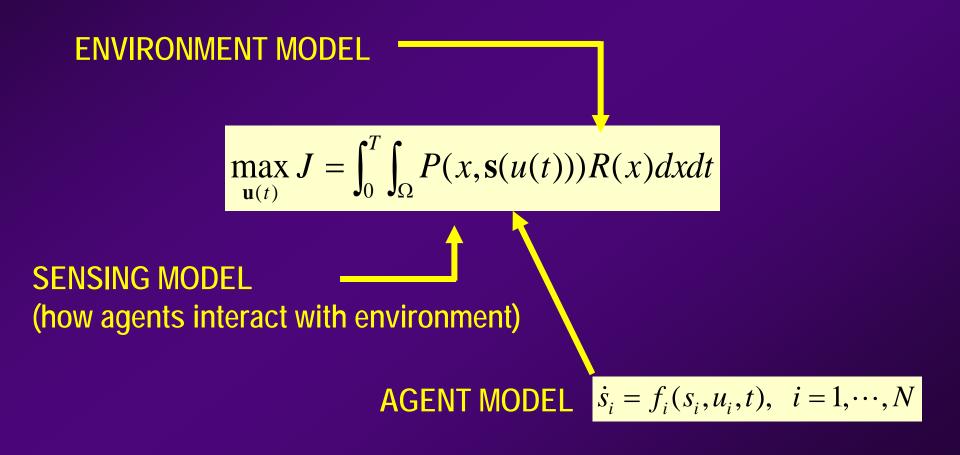
CONSENSUS



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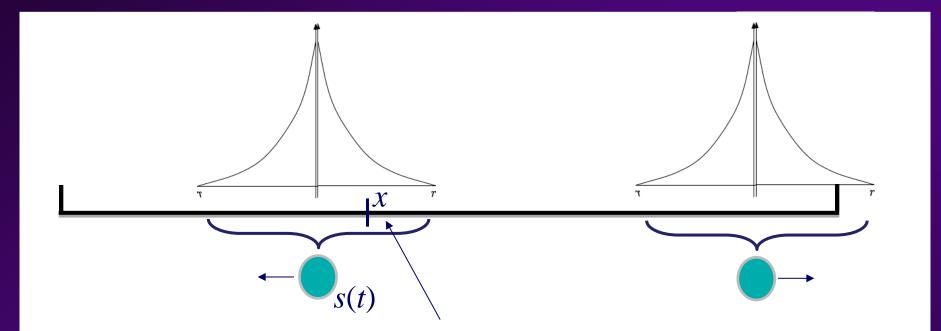
PERSISTENT MONITORING

GOAL: Find the best state trajectories $s_i(t)$, $0 \le t \le T$ so that agents achieve a maximal reward from interacting with the mission space



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PERSISTENT MONITORING



ENVIRONMENT MODEL: Associate to *x* Uncertainty Function *R*(*x*,*t*)

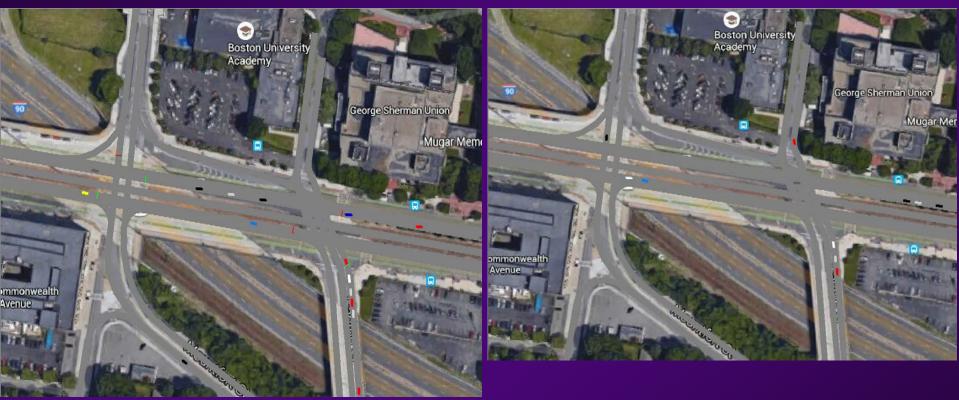
$$\dot{R}(x,t) = \begin{cases} 0 & \text{if } R(x,t) = 0, A(x) < Bp(x,s(t)) \\ A(x) - Bp(x,s(t)) & \text{otherwise} \end{cases}$$
NOTE: Could be stochastic !

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THE INTERNET OF CARS...

With traffic lights (non-Cooperative)

No traffic lights: decentralized control of CAVs (Cooperative)



One of the worst-designed double intersections ever... (BU Bridge – Commonwealth Ave, Boston)

Malikopoulos, Cassandras, Zhang et al, Automatica, 2018

Zhang et al, Proc. of IEEE, 2018

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RELATED WORK

COVERAGE AND FORMATION CONTROL

Choset 2001, Leonard and Olshevsky 2013, Tron et al 2014, Egerstedt and Hu 2001

Cortes et al 2004 Zhong and Cassandras 2010, Sun and Cassandras 2016

CONSENSUS

Jadbabaie et al, 2003, Olfati-Saber and Murray, 2004, Ren and Beard 2005 Nedich et al, 2010

SAMPLING AND TRACKING

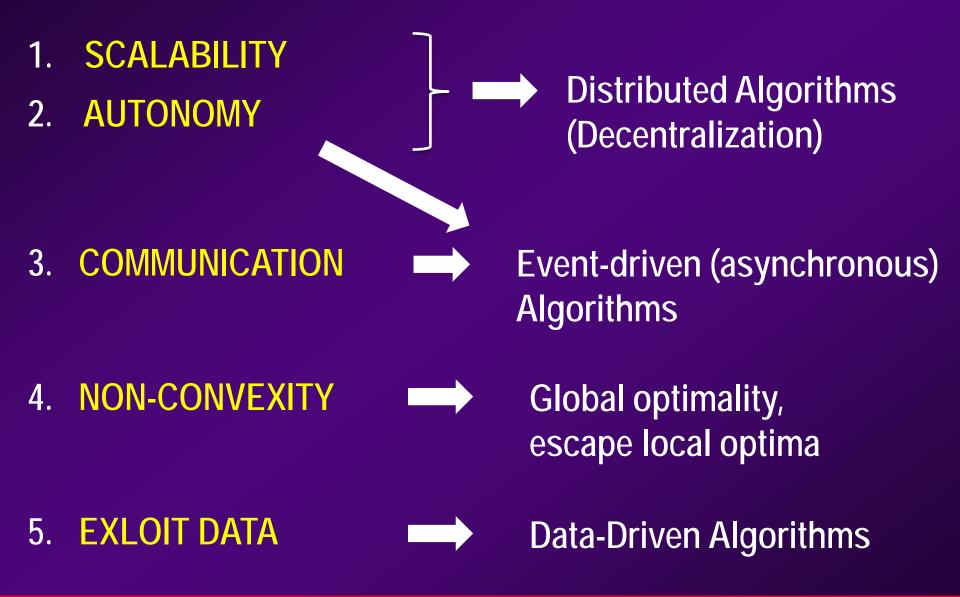
Leonard and Zhang 2010, Ashley and Andersson 2016

PERSISTENT MONITORING

Smith et al, 2011, Michael et al, 2011, Lan and Schwager, 2014

Cassandras et al, 2013, Yu et al, 2017

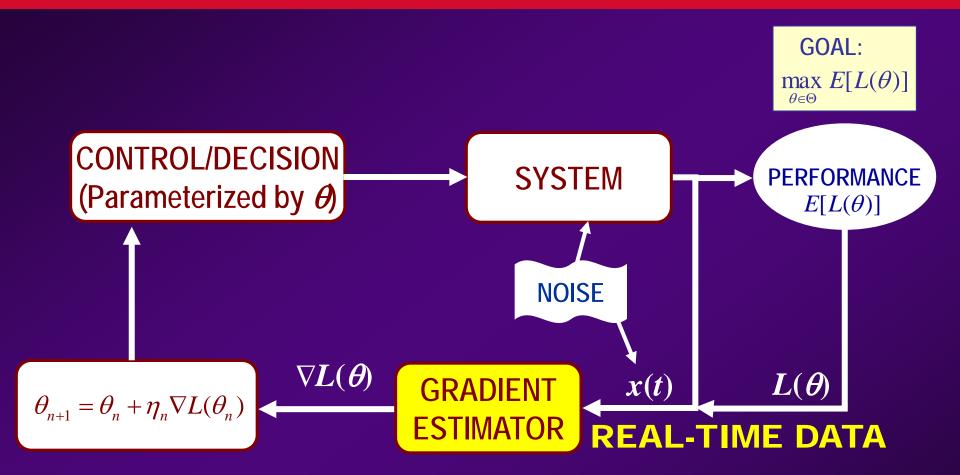
NETWORKED MULTI-AGENT OPTIMIZATION-CHALLENGES



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EVENT-DRIVEN + DATA-DRIVEN OPTIMIZATION

DATA-DRIVEN STOCHASTIC OPTIMIZATION



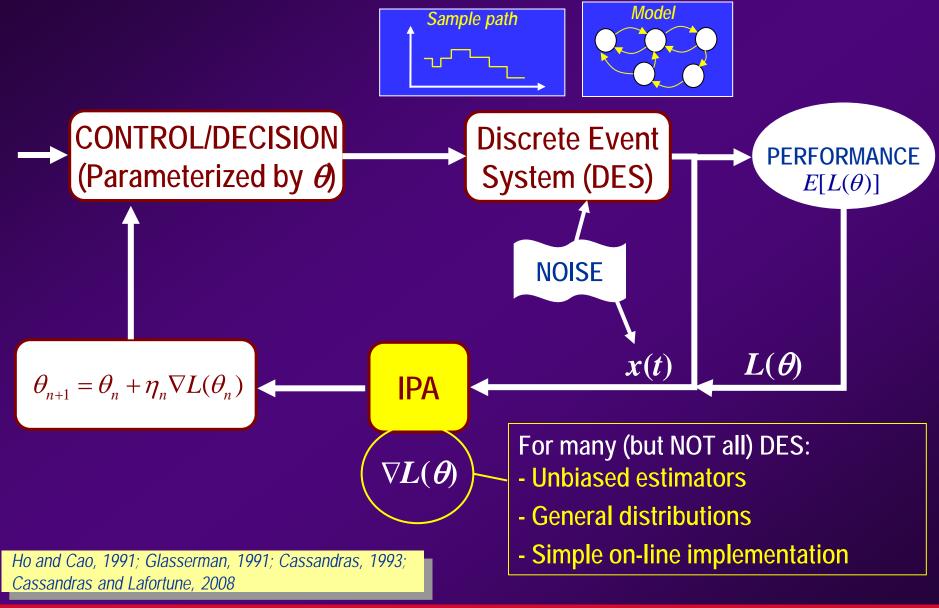
DIFFICULTIES: - $E[L(\theta)]$ NOT available in closed form

- - $\nabla L(\theta)$ not easy to evaluate
- $-\nabla L(\theta)$ may not be a good estimate of $\nabla E[L(\theta)]$

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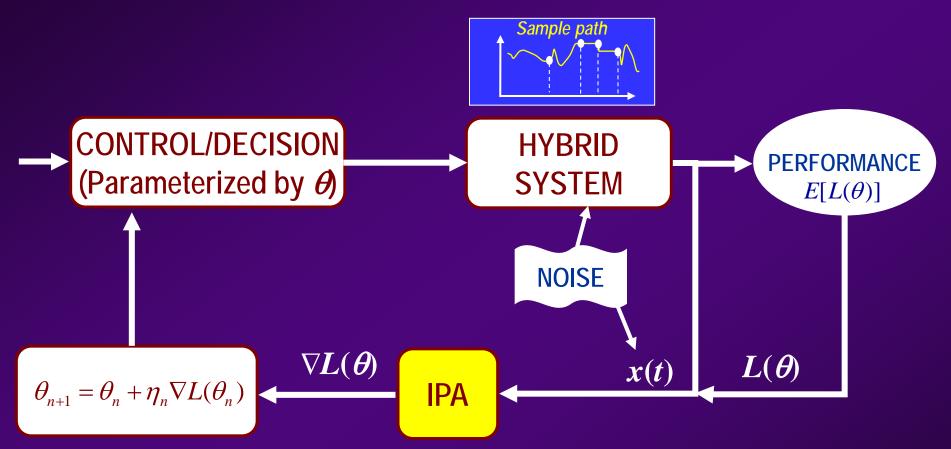
REAL-TIME STOCHASTIC OPTIMIZATION FOR DES: INFINITESIMAL PERTURBATION ANALYSIS (IPA)



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REAL-TIME STOCHASTIC OPTIMIZATION: *HYBRID SYSTEMS, CYBER-PHYSICAL SYSTEMS*



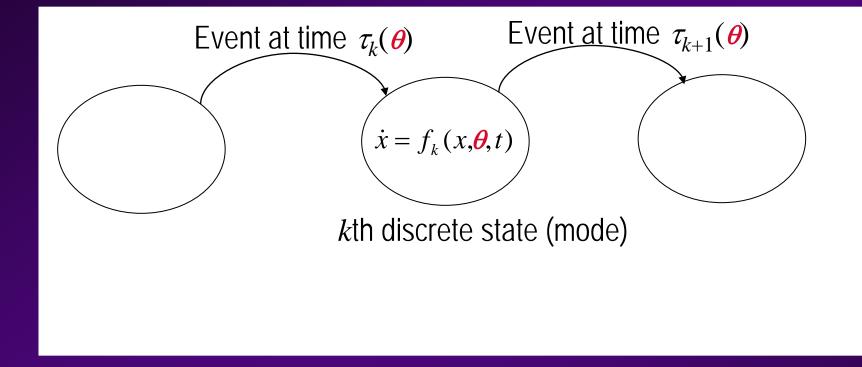
A general framework for an IPA theory in Hybrid Systems?

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THE IPA CALCULUS

STOCHASTIC HYBRID AUTOMATA



 θ : control parameter, $\theta \in \Theta$ (system design parameter,parameter of an input process,or parameter that characterizes a control policy)

System dynamics over
$$(\tau_k(\theta), \tau_{k+1}(\theta)]$$
: $\dot{x} = f_k(x, \theta, t)$

OTATION:
$$x'(t) = \frac{\partial x(\theta, t)}{\partial \theta}, \quad \tau'_k = \frac{\partial \tau_k(\theta)}{\partial \theta}$$

1. Continuity at events: $x(\tau_k^+) = x(\tau_k^-)$

Take $d/d\theta$:

Ν

$$x'(\tau_k^+) = x'(\tau_k^-) + [f_{k-1}(\tau_k^-) - f_k(\tau_k^+)]\tau'_k$$

If no continuity, use reset condition \Rightarrow

$$x'(\tau_k^+) = \frac{d\rho(q, q', x, \upsilon, \delta)}{d\theta}$$

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2. Take $d/d\theta$ of system dynamics $\dot{x} = f_k(x, \theta, t)$ over $(\tau_k(\theta), \tau_{k+1}(\theta)]$:

$$\frac{dx'(t)}{dt} = \frac{\partial f_k(t)}{\partial x} x'(t) + \frac{\partial f_k(t)}{\partial \theta}$$

Solve
$$\frac{dx'(t)}{dt} = \frac{\partial f_k(t)}{\partial x} x'(t) + \frac{\partial f_k(t)}{\partial \theta}$$
 over $(\tau_k(\theta), \tau_{k+1}(\theta)]$:

$$x'(t) = e^{\int_{\tau_k}^{t} \frac{\partial f_k(u)}{\partial x} du} \left[\int_{\tau_k}^{t} \frac{\partial f_k(v)}{\partial \theta} e^{-\int_{\tau_k}^{v} \frac{\partial f_k(u)}{\partial x} du} dv + x'(\tau_k^+) \right]$$

initial condition from 1 above

NOTE: If there are no events (pure time-driven system), IPA reduces to this equation

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3. Get τ'_k depending on the event type:

- Exogenous event: By definition, $\tau'_k = 0$
- Endogenous event: occurs when $g_k(x(\theta, \tau_k), \theta) = 0$

$$\tau'_{k} = -\left[\frac{\partial g}{\partial x}f_{k}(\tau_{k}^{-})\right]^{-1}\left(\frac{\partial g}{\partial \theta} + \frac{\partial g}{\partial x}x'(\tau_{k}^{-})\right)$$

- Induced events:

$$\tau'_{k} = -\left[\frac{\partial y_{k}(\tau_{k})}{\partial t}\right]^{-1} y'_{k}(\tau_{k}^{+})$$

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Ignoring resets and induced events:

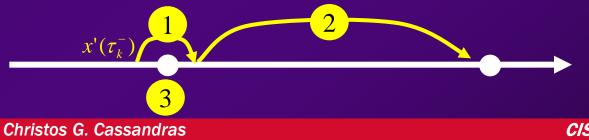
1.
$$x'(\tau_k^+) = x'(\tau_k^-) + [f_{k-1}(\tau_k^-) - f_k(\tau_k^+)] \cdot \tau'_k$$

2. $x'(\tau_k) = e^{\int_{\tau_{k-1}}^{\tau_k} \frac{\partial f_k(u)}{\partial x} du} \int_{\tau_k}^{\tau_k} \frac{\partial f_k(v)}{\partial \theta} e^{-\int_{\tau_{k-1}}^{v} \frac{\partial f_k(u)}{\partial x} du} dv + x'(\tau_{k-1}^+)$

 τ_{k-1}

$$x'(t) = \frac{\partial x(\theta, t)}{\partial \theta}$$
$$\tau'_{k} = \frac{\partial \tau_{k}(\theta)}{\partial \theta}$$

3.
$$\tau'_{k} = 0$$
 or $\tau'_{k} = -\left[\frac{\partial g}{\partial x}f_{k}(\tau_{k}^{-})\right]^{-1}\left(\frac{\partial g}{\partial \theta} + \frac{\partial g}{\partial x}x'(\tau_{k}^{-})\right)$



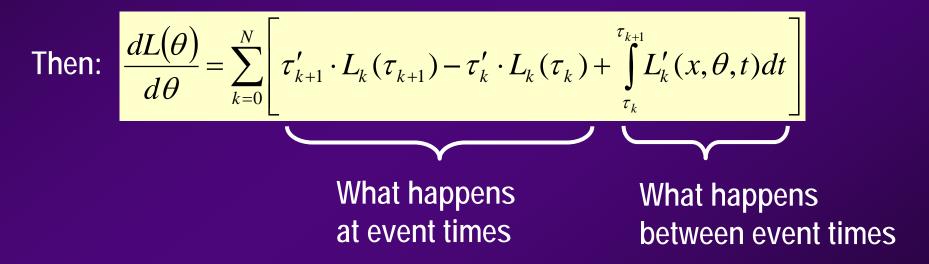
Cassandras et al, Europ. J. Control, 2010

IPA PROPERTIES

Back to performance metric:

$$L(\theta) = \sum_{k=0}^{N} \int_{\tau_k}^{\tau_{k+1}} L_k(x,\theta,t) dt$$

NOTATION:
$$L'_k(x,\theta,t) = \frac{\partial L_k(x,\theta,t)}{\partial \theta}$$



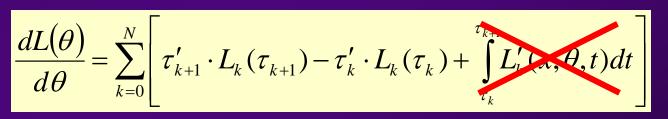
IPA PROPERTIES: *ROBUSTNESS*

THEOREM 1: If either 1,2 holds, then $dL(\theta)/d\theta$ depends only on information available at event times τ_k :

- 1. $L(x, \theta, t)$ is independent of t over $[\tau_k(\theta), \tau_{k+1}(\theta)]$ for all k
- 2. $L(x, \theta, t)$ is only a function of x and for all t over $[\tau_k(\theta), \tau_{k+1}(\theta)]$:

 $\frac{d}{dt}\frac{\partial L_k}{\partial x} = \frac{d}{dt}\frac{\partial f_k}{\partial x} = \frac{d}{dt}\frac{\partial f_k}{\partial \theta} = 0$

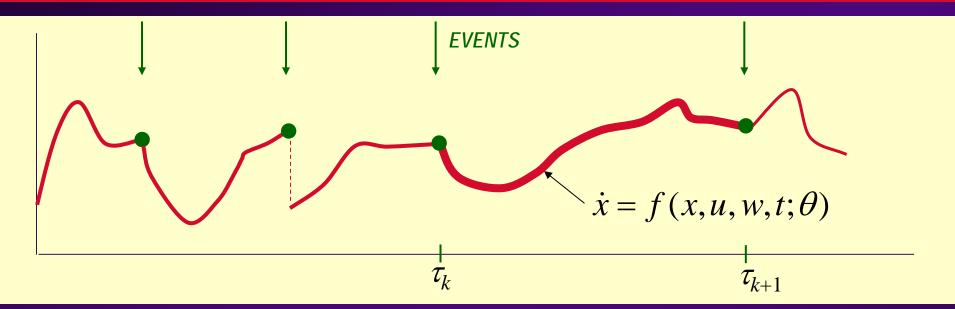
Yao and Cassandras, J. DEDS, 2011



 IMPLICATION: - Performance sensitivities can be obtained from information limited to event times, which is easily observed
 - No need to track system in between events !

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IPA PROPERTIES



Evaluating $x(t; \theta)$ requires full knowledge of w and f values (obvious)

However, $\frac{dx(t;\theta)}{d\theta}$ may be *independent* of *w* and *f* values (*NOT* obvious)

It often depends only on: - event times τ_k - possibly $f(\tau_{k+1}^-)$

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IPA estimators are EVENT-DRIVEN \Rightarrow IPA scales with the EVENT SET, not the STATE SPACE ! \Rightarrow no time discretization needed

As a complex system grows with the addition of more states, the number of EVENTS often remains unchanged or increases at a much lower rate.

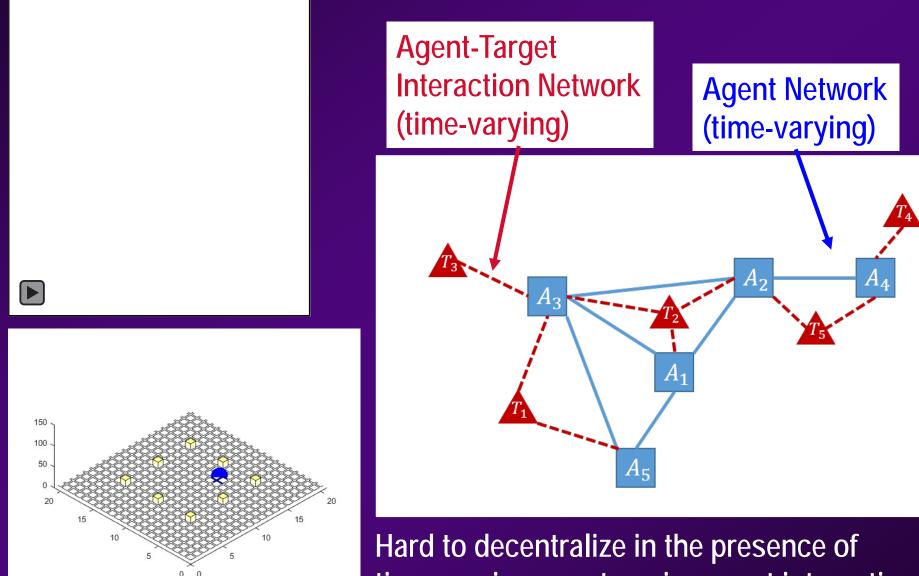
EXAMPLE: A queueing network may become very large, but the basic events used by IPA are still "arrival" and "departure" at different nodes.

DECENTRALIZING CAN BE HARD

DECENTRALIZED SOLUTION = CENTRALIZED SOLUTION (AGENTS ACTING USING ONLY LOCAL INFO.)

(no performance loss due to decentralization)

PERSISTENT MONITORING WITH KNOWN TARGETS



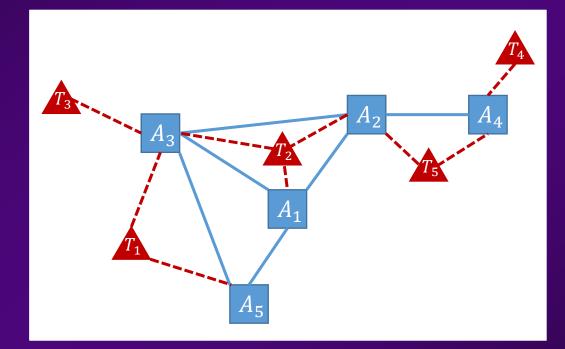
time-varying agent-environment interactions

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THREE TYPES OF NEIGHBORHOODS

The agent neighborhood of an agent (conventional) The agent neighborhood of agent j is the set $\mathcal{A}_j(t) = \{k : ||s_k(t) - s_j(t)|| \le r_c, k \ne j, k = 1, \dots, N\}.$ The target neighborhood of an agent The target neighborhood of agent j is the set $\mathcal{T}_{j}(t) = \{i : |x_{i} - s_{j}(t)| \le r_{j}, i = 1, ..., M\}.$



The agent neighborhood of an target The agent neighborhood of target *i* is the set $\mathcal{B}_i(t) = \{j : |s_j(t) - x_i| \le r_j, j = 1, ..., N\}.$

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"ALMOST DECENTRALIZATION" RESULT

- Show that optimal trajectories consist of *hybrid dynamics*: segments defined by observable *EVENTS* e.g., agent enters target sensing range, agent leaves neighborhood
- Develop EVENT-DRIVEN gradient-based algorithms using the Infinitesimal Pertubation Analysis (IPA) calculus: Each agent evaluates its IPA derivative
- Does an agent's IPA derivative depend only on LOCAL events?
 DECENTRALIZATION EVENT OBSERVABILITY

THEOREM: Each agent's IPA derivative depends only on LOCAL events except for one global event

Zhou et al, IEEE TAC 2018

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NETWORKED MULTI-AGENT OPTIMIZATION-CHALLENGES

