# Numerical tensor methods and their applications 

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## What is this course is about

This course is mostly on numerical methods of linear algebra in multilinear settings.

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This course is mostly on numerical methods of linear algebra in multilinear settings.

Goal: develop universal tools for working with high-dimensional problems.

## All lectures

4 lectures,

- 2 May, 08:00-10:00: Introduction: ideas, matrix results, history.
- 7 May, 08:00-10:00: Novel tensor formats (TT, HT, QTT).
- 8 May, 08:00-10:00: Advanced tensor methods (eigenproblems, linear systems).
- 14 May, 08:00-10:00: Advanced topics, recent results and open problems.


## Lecture 1

- Motivation
- Matrix background
- Canonical and Tucker formats
- Historical overview


## Motivation

## Main points

- High-dimensional problems appear in diverse applications
- Standard methods do not scale well in many dimensions


## Motivation

Solution of high-dimensional differential and integral equations on fine grids
Typical cost: $\mathcal{O}\left(N^{3}\right) \rightarrow \mathcal{O}(N)$ or even $\mathcal{O}\left(\log ^{\alpha} N\right)$.


## Motivation

Ab initio computations and computational material design
Protein-ligand docking (D. Zheltkov)

Density functional theory for large clusters (V. Khoromskaia)



## Motivation

Construction of reduced order models for multiparametric/stochastic systems in engineering


## Diffusion problem

$\nabla a(p) \Delta u=f(p)$,
$p=\left(p_{1}, p_{2}, p_{3}, p_{4}\right)$
Approximate $u$ using only few snapshots.

## Motivation

## Data mining and compression

Images


Computational data (temperature)

## Why tensors are important

The multivariate functions are related to the multivariate arrays, or tensors:

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The multivariate functions are related to the multivariate arrays, or tensors:

Take a function: $f\left(x_{1}, \ldots, x_{d}\right)$
Take tensor-product grid
Get a tensor:

$$
A\left(i_{1}, \ldots, i_{d}\right)=f\left(x_{1}\left(i_{1}\right), \ldots, x_{d}\left(i_{d}\right)\right)
$$

## Literature

- T. Kolda and B. Bader, Tensor decompositions and applications, SIREV (2009)
- W. Hackbusch, Tensor spaces and numerical tensor calculus, 2012
- L. Grasedyck, D. Kressner, C. Tobler, A literature survey of low-rank tensor approximation techniques, 2013


## Software

## Some software will be used:

- Tensor Toolbox 2.5 (T. Kolda)
- TT-Toolbox
(http://github.com/oseledets/TT-Toolbox)
There is also a Python version
(http://github.com/oseledets/ttpy) which has similar functionality now.


## Where tensors come from

- d-dimensional PDE: $\Delta u=f, u=u\left(x_{1}, \ldots, x_{d}\right)$
- PDE with $M$ parameters: $A(p) u(p)=f(p)$, $u=u\left(x, p_{1}, \ldots, p_{M}\right)$
- Data (images, video, hyperspectral images)
- Latent variable models, joint probability distributions
- Factor models
- Many others


## Definitions

A tensor is a d-dimensional array:

$$
A\left(i_{1}, \ldots, i_{d}\right), \quad 1 \leq i_{k} \leq n_{k}
$$

Mathematically more correct definition:
Tensor is a polylinear form.

## Definitions

Tensors form a linear vector space.
The natural norm is the Frobenius norm:

$$
\|A\|=\sqrt{\sum_{i_{1}, \ldots, i_{d}}\left|A\left(i_{1}, \ldots, i_{d}\right)\right|^{2}}
$$

## Curse of dimensionality

## Curse of dimensionality: Storage of a d-tensor with mode sizes $n$ requires $n^{d}$ elements.

## Basic questions

- How to break the curse of dimensionality?
- How to perform (multidimensional) sampling?
- How to do everything efficiently and in a robust way?


## Real-life problems

If you really need to compute something high-dimensional , there is usually a way:

- Monte Carlo
- Special basis sets (radial basis functions)
- Best N-term approximations (wavelets, sparse grids)

But we want algebraic techniques...

## Separation of variables

One of the few fruitful ideas is the idea of separation of variables

## What is separation of variables

Separation rank 1:
$f\left(x_{1}, \ldots, x_{d}\right)=u_{1}\left(x_{1}\right) u_{2}\left(x_{2}\right) \ldots u_{d}\left(x_{d}\right)$,
More general:
$f\left(x_{1}, \ldots, x_{d}\right) \approx \sum_{\alpha=1}^{r} u_{1}\left(x_{1}, \alpha\right) \ldots u_{d}\left(x_{d}, \alpha\right)$.

## Analytical examples

How to compute separated representations?
Analytical expressions (B. N. Khoromskij and many others):

$$
\begin{gathered}
f\left(x_{1}, \ldots, x_{d}\right)=\frac{1}{x_{1}+\ldots+x_{d}} \text { based on the identity } \\
\frac{1}{x}=\int_{0}^{\infty} \exp (-p x) d p \\
r=\log \varepsilon^{-1} \log \delta^{-1}
\end{gathered}
$$

## Numerical computation of separated representations

We can try to compute the separated decomposition numerically.
How do we do that?

## Canonical format

## Tensors:

Canonical format:

$$
\begin{gathered}
A\left(i_{1}, \ldots, i_{d}\right) \approx \sum_{\alpha=1}^{r} U_{1}\left(i_{1}, \alpha\right) \ldots U_{d}\left(i_{d}, \alpha\right) \\
\text { What happens in } d=2 ?
\end{gathered}
$$

## Two-dimensional case

## $A\left(i_{1}, i_{2}\right) \approx \sum_{\alpha=1}^{r} U_{1}\left(i_{1}, \alpha\right) U_{2}\left(i_{2}, \alpha\right)$

## Two-dimensional case

$A\left(i_{1}, i_{2}\right) \approx \sum_{\alpha=1}^{r} U_{1}\left(i_{1}, \alpha\right) U_{2}\left(i_{2}, \alpha\right)$
Matrix form: $A \approx U V^{\top}$,
Where $U$ is $n \times r, V$ is $m \times r$
Approximate rank-r approximation

## SVD: definition

The fabulous SVD (singular value decomposition):
Every matrix can be represented as a product $A=U S V^{*}$, where $U, V$ are orthonormal, $S$ is a diagonal matrix with singular values $\sigma_{i} \geq 0$ on the diagonal.

## SVD: complexity

## Complexity of the SVD is $\mathcal{O}\left(n^{3}\right)$ (too much to compute $\mathcal{O}(n r)$ decomposition)

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Are there faster algorithms?

## Skeleton decomposition

Yes: based on the skeleton decomposition

$$
A \approx C \widehat{A}^{-1} R
$$

$C-r$ columns of $A, R-r$ rows of $A, \widehat{A}-$ submatrix on the intersection.

- Ex.1: Prove it
- Ex.2: Have you met skeleton dec. before?


## Maximum volume principle

What happens if the matrix is of approximate low rank?

$$
A \approx R+E, \quad \operatorname{rank} R=r, \quad\|E\|_{C}=\varepsilon
$$

## Maximum volume principle

Select the submatrix $\widehat{A}$ such that volume is maximal (volume $=$ absolute value of the determinant)

$$
\left\|A-C \widehat{A}^{-1} R\right\| \leq(r+1)^{2} \varepsilon
$$

## Proof

E. E. Tyrtyshnikov, S.A. Goreinov, On quasioptimality of skeleton approximation of a matrix in the Chebyshev norm, doi: 10.1134/S1064562411030355

$$
A=\left(\begin{array}{ll}
A_{11} & A_{12} \\
A_{21} & A_{22}
\end{array}\right)
$$

$$
H=A-C \widehat{A}^{-1} R=A-\binom{A_{11}}{A_{21}} A_{11}^{-1}\left(A_{11} A_{21}\right)
$$

Need: $\left|h_{i j}\right| \leq(r+1)^{2} \delta_{r+1}(A)$

## Proof

$$
Z=\left(\begin{array}{cc}
A_{11} & v \\
u^{\top} & a_{i j}
\end{array}\right)
$$

Entry $h_{i j}$ can be found from:

$$
\begin{gathered}
\left(\begin{array}{cc}
l & 0 \\
-u^{\top} A_{11}^{-1} & 1
\end{array}\right) Z=\left(\begin{array}{cc}
A_{11} & v \\
0 & h_{i j}
\end{array}\right) \\
\operatorname{det} Z=h_{i j} \operatorname{det} A_{11}
\end{gathered}
$$

Therefore,

$$
\begin{gathered}
\left|h_{i j}^{-1}\right|=\left\|Z^{-1}\right\|_{C} \\
\left|h_{i j}\right| \leq(r+1) \sigma_{r+1}(Z)
\end{gathered}
$$

## Proof

Finally,

$$
\begin{gathered}
\sigma_{r+1}(Z)=\min _{U_{Z}, V_{Z}}\left\|Z-U_{Z} V_{Z}^{\top}\right\|_{2} \leq \\
(r+1)\left\|Z-U_{Z} V_{Z}^{\top}\right\|_{C} \leq(r+1) \delta_{r+1}(A)
\end{gathered}
$$

## Maxvol algorithm(1)

Ok, then, how to find a good submatrix?
Crucial algorithm: Maxvol submatrix in a $n \times r$ matrix.

Characteristic property: $A$ is $n \times r$,

$$
A \widehat{A}^{-1}=\binom{I}{Z}, \quad|Z|_{i j} \leq 1
$$

## Maxvol algorithm(2)

Problem: find maximal volume $r \times r$ submatrix in an $n \times r$ matrix.

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Maxvol algorithm:

- Take some rows, put them in the first $r$.
- Compute $B=A \widehat{A}^{-1}$
- $B=\binom{l}{Z}$
- Suppose maximal element in $Z$ is in position $(i, j)$.
- Swap $i$-th row with $j$-th row.
- Stop if maximal element is less than $(1+\delta)$.


## Maxvol algorithm(2)

Problem: find maximal volume $r \times r$ submatrix in an $n \times r$ matrix.

For an $n \times m$ matrix:
Find maximal volume in rows, then find maximal volume in columns

Ex. Implement an algorithm that searches for a maxvol submatrix.

## Maxvol algorithm (demo)

Let us see how maxvol works. . .

## Cross approximation

A typical scheme we use is the cross approximation approach, which uses minimal information from the matrix.

## Cross approximation

(1) $k=0$, Select $j_{0}, U_{0}=0, V_{0}=0$.
(2) Compute $j_{k}$-th column of the remainder $A_{k}=A-U_{k} V_{k}^{\top}$.
(3) Find maximal element $i_{k}$ in it, compute $i_{k}$-th row, compute maximal element $j_{k+1} \neq j_{k}$.
(9) Compute the next cross: $u_{k}=A_{k} e_{j_{k}}$,
$v_{k}=A_{k}^{\top} e_{i_{k}}, u_{k}=u_{k} / A_{k}\left(i_{k}, j_{k}\right)$,
$U_{k}=\left[U_{k-1}, u_{k}\right], V_{k}=\left[V_{k-1}, v_{k}\right]$.
(5) If $\| u_{k} v_{k}^{\top}$ is small, stop, else go to 1 .

## Randomized techniques

Randomized techniques for low-rank approximation became popular recently.
Sublinear randomized algorithms for skeleton decompositions, Jiawei Chiu and Laurent Demanet, http://arxiv.org/abs/1110.4193v2

## Theorem

Let $A=U S V^{\top}$ and $U$ and $V$ are $\mu$-coherent, i.e. $\|U\|_{C} \leq \sqrt{\frac{\mu}{n}}$
Then, with high probability, one has to sample $I=\mu r \log n$ columns and rows uniformly, to get a $\mathcal{O}\left(\sigma_{r+1}\right)$ bound.

## What is the best cross algorithm?

I strongly believe, that the "best" cross algorithm is still to be found

And it is very important in higher dimensions!

## Going to higher dimensions

How to generalize the idea of separation of variables to higher dimensions?

- SVD is good
- Best approximation exists
- Interpolation via skeleton


## Canonical format (2)

$$
\begin{aligned}
& A\left(i_{1}, \ldots, i_{d}\right) \approx \sum_{\alpha=1}^{r} U_{1}\left(i_{1}, \alpha\right) \ldots U_{d}\left(i_{d}, \alpha\right) \\
& r \text { is called (approximate) canonical rank, } U_{k}- \\
& \text { canonical factors. }
\end{aligned}
$$

## Canonical format(3)

Good things about the canonical format:

- Low number of parameters dnr
- Uniqueness results (Kruskal theorem)


## Canonical format(3)

Let $A$ be a 3-tensor with $(U, V, W)$ canonical decomposition of rank $R$,
, and $k(U)+k(V)+k(W) \geq 2 R+3$, then the decomposition is unique.
$k(X)$ Kruskal rank (spark in compressed sensing),
Def: $k(X)+1$ is the minimal number of linearly dependent columns in $X$.
Proof is highly nontrivial (Est time: ~1.5 lectures!)

## Canonical format (4)

Bad things about the canonical format:

- Best approximation may not exist
- Canonical rank is NP-complete (matrix rank is
...)
- No good algorithm


## Bad example (1)

$$
f\left(x_{1}, \ldots, x_{d}\right)=x_{1}+x_{2}+\ldots x_{d}
$$

Canonical rank $d$ (no proof is known), can be approximated with rank-2 with any accuracy!

## Bad example (2)

Canonical rank may depend on the field (matrix rank can not!)

$$
f\left(x_{1}, \ldots, x_{d}\right)=\sin \left(x_{1}+\ldots+x_{d}\right)
$$

- Complex field: 2
- Real field: $d$ (Ex.: prove it)


## Alternating least squares

The main algorithm for the computation of the canonical decomposition is the Alternating Least Squares (ALS) algorithm.

- Easy to implement
- Known for its very slow convergence (swaps)
- Local convergence proven only recently (Uschmajew, A.)


## Alternating least squares

Treat approximation as an optimization problem:

$$
\begin{gathered}
\|A-(U, V, W)\|_{F} \rightarrow \min \\
\text { Three steps: }
\end{gathered}
$$

(1) Fix $V, W$, update $U$ (linear least squares)
(2) Fix $U, W$, update $V$
(3) Fix $U, V$, update $W$.

Exercise: write down comp. formula and implement them.

## Example from the complexity theory

There are cases, where the canonical format comes from a model:

Matrix multiplication:

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Matrix multiplication:

$$
C=A B, c=f(a, b), c_{i}=\sum_{i j} E_{i j k} a_{j} b_{k}
$$

If the canonical rank of $E$ is $r$, computation of $C$ requires $r$ multiplications

- $2 \times 2$ : $4 \times 4 \times 4$ tensor, rank 7 (Strassen)
- $3 \times 3: 9 \times 9 \times 9$ tensor, rank is unknown $19 \leq r \leq 23$.


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It is fascinating.

## What about sampling?

## Can we generalize skeleton decomposition?

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## Can we generalize skeleton decomposition?

No
Try it yourself: a simple generalization of a "cross".

## Another attempt: Tucker

Another attempt to avoid was the Tucker format (Tucker 1966, Lathauwer, 2000+)
$A(i, j, k) \approx$
$\sum_{\alpha \beta \gamma} G(\alpha, \beta, \gamma) U_{1}(i, \alpha) V(j, \alpha) W(k, \alpha)$

## Tucker and SVD

You can compute Tucker by means of the SVD:

- Compute unfoldings: $A_{1}, A_{2}, A_{3}$
- Compute left SVD factors: $A_{i} \approx U_{i} \Phi_{i}$
- Compute the core: $G=A \times_{1} U_{1}^{\top} \times_{2} U_{2}^{\top} \times_{3} U_{3}^{\top}$.


## Tucker and the cross

You can generalize skeleton to Tucker (O., Savostyanov, Tyrtyshnikov, 2008)
Compute good columns in $A_{i}$, find core by interpolation.

## Problem with the Tucker format

Q: What is the main problem with the Tucker format?

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Q: What is the main problem with the Tucker format?

A: Curse of dimensionality
The core takes $r^{d}$ elements!

## Summary

## What we have?

- Canonical format: low number of parameters, no algorithms
- Tucker format: SVD-based algorithms, the curse


## Main algebraic question

## Can we find something inbetween?

## Lecture 2

- The Tree-Tucker, Tensor Train, Hierarchical Tucker formats
- Their difference
- Concept of Tensor Networks
- Stability and quasioptimality
- Basic arithmetic (with illustration)
- Cross approximation formula (with illustrations)
- QTT-format (part 1)


## Lecture 3

- QTT-format (part 2), application to numerical integration
- QTT-Fourier transform and its relation to tensor networks
- QTT-convolution, explicit representation of Laplace-like tensors
- DMRG/AMEN techniques
- Solution of linear systems in the TT-format
- Solution of eigenvalue problems in the TT-format


## Lecture 4

Advanced topics: New applications, recent results and open problems

- Solution of non-stationary problems
- Global optimization via the TT-cross
- Latent variable models (finance and natural language processing)
- Approximation results in quantum information theory (Hastings area law)
- Open problems

